

SPECTRUM

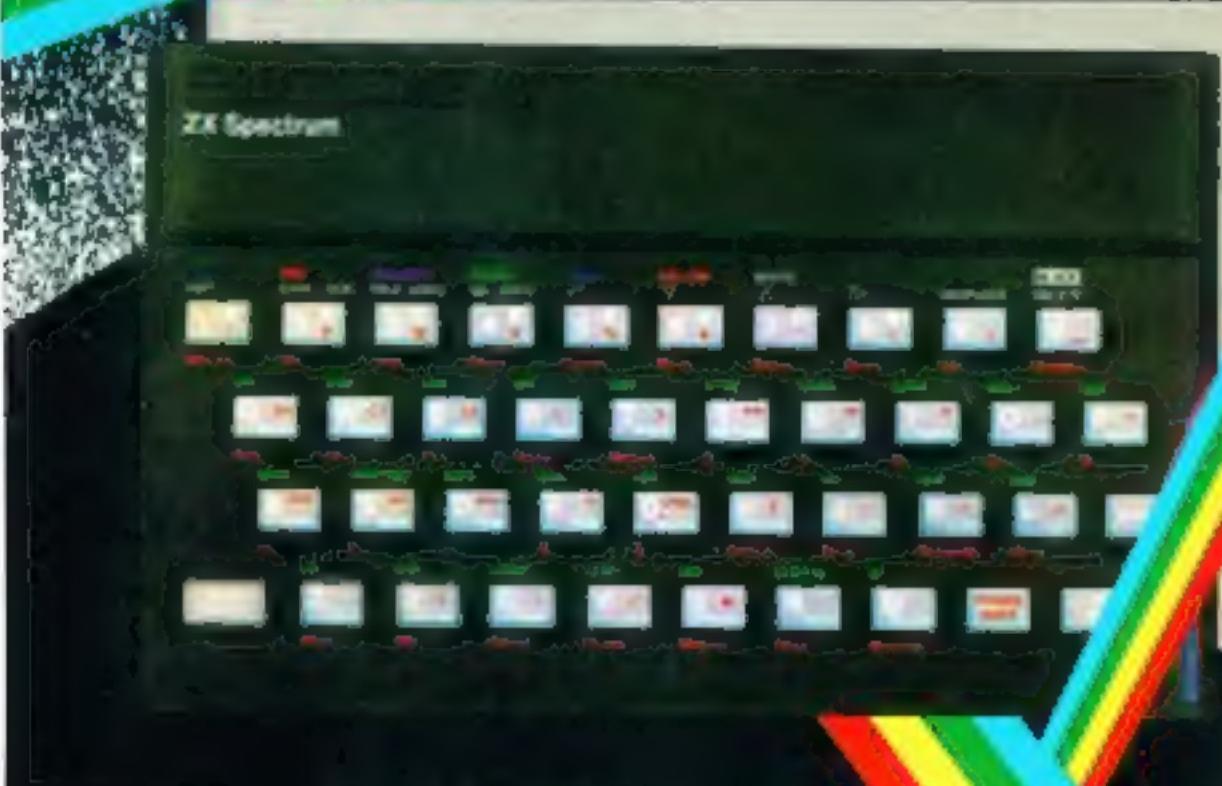


Melbourne
House

SPECTRUM MACHINE LANGUAGE FOR THE ABSOLUTE BEGINNER

Edited by William Tang

ZX Spectrum



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The Beginning

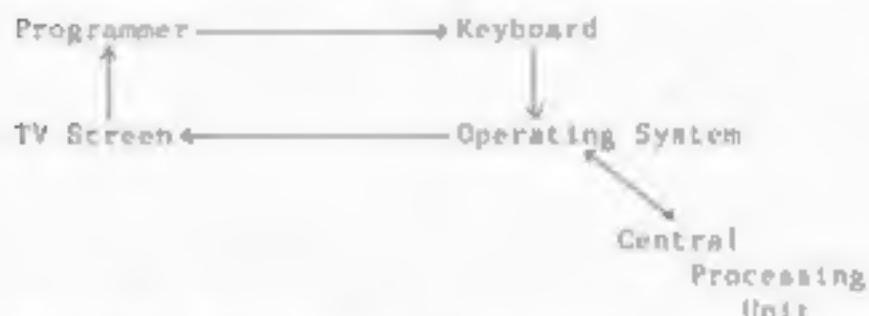
This book is designed as an introduction to the field of machine and assembly language programming for the "Sinclair ZX Spectrum."

It may be that you are coming to this book with no clear idea of what machine language programming is all about.

You may not even know what machine language is. You may not even be aware that there is a difference between machine language and assembly language, nor indeed how they differ from programming in BASIC.

Don't worry, and don't be frightened by the jargon - we will explain everything step by step.

First, let us look at the way a computer operates:



What this diagram aims to show is that there is a barrier between the programmer and the brain of the Spectrum, the Central Processing Unit. It is not possible under normal processing for the programmer to tell the Central Processing Unit - usually referred to as the CPU - what to do.

In the Sinclair machines, the CPU chosen is the Z80A chip, which is a faster version of the popular Z80 chip. There are 4 chips - Z80, 6502, 6809, and 8088 - which have become widely accepted as CPUs for microcomputers. The Z80 is by far the most popular chip.

I am sure it comes as no surprise to learn that the Z80A does not understand a word of 'BASIC'! Indeed no CPU has been designed so that we can communicate directly with the brain of the computer.

If you think about it for long enough, you will realise that it would be very difficult, if not impossible, in any case to give a chip in a computer an instruction that would make any sense to a human. Take the top off your Sinclair (if you dare!) and have a look at the chip nearest to the speaker - this is the Z80A CPU. Obviously this chip in your computer can only respond to electrical signals that are passed on to it by the rest of the circuitry!

What Is Machine Language?

The Z80 chip has been designed in such a way that it can accept signals simultaneously from eight of the pins connected to it.

The designers of the Z80 chip constructed it in such a way that different combinations of signals to the Z80 chip along these eight pins would 'instruct' the Z80 to perform different functions.

Keeping in mind that what is really happening are electrical signals, let's adopt a convention to represent these signals - for example showing a '1' if there is a signal to one of the pins, or a '0' if there is no signal.

A typical instruction might therefore look something like:

0 0 1 1 1 1 0 0

Quite a long way from something like

'Let A = A + 1',
for example, isn't it!

Nonetheless, this is what machine language is all about. The name says it all! It is a language for machines. Each manufacturer of the different chips has designed a different 'language' for its products!

At this stage you may be asking yourself - if this is what machine language is all about, why bother? Why not accept the benefits of someone else's work which allows me to program the computer in a language I can easily understand, such as BASIC or COBOL?

The reason is because of the main benefits of machine language which are:

- FASTER EXECUTION OF THE PROGRAM
- MORE EFFICIENT USE OF MEMORY
- SHORTER PROGRAMS (in memory)
- FREEDOM FROM THE OPERATING SYSTEM

All of the above benefits are a direct result of programming in a language that the CPU can understand without having to have it translated first. When you program in BASIC, the operating system is the machine language program that is really being run by the machine. The program is something like:

Next	Look at next instruction
	Translate it into a series of
	machine language instructions
	Perform each instruction
	Store the result if required
	Go to Next again

If you are wondering where the computer finds this program, the operating system, it is in the ROM. In other words, it is built into the Spectrum. (ROM is the abbreviation for Read Only Memory, memory locations whose content you cannot change, but can only be read/PEEKed.)

Programming in BASIC can be up to 60 times slower than a program written directly in machine language!

This is because translation takes time, and also the resulting machine language instructions generated usually are less efficient. Similarly, it is usually faster to drive yourself than to take public transport; you can take shortcuts you know, instead of following the public transport route which needs to cater for the GENERAL public CASE.

Nonetheless, we would have to be among the first to admit that programming in machine language does have drawbacks.

The main disadvantage of machine language are:

- * PROGRAMS ARE DIFFICULT TO READ AND DEBUG
- * IMPOSSIBLE TO ADAPT TO OTHER COMPUTERS
- * LONGER PROGRAMS (in instructions)
- * ARITHMETIC CALCULATIONS DIFFICULT

This means that you must make a very conscious decision of which programming method you should use for each particular application.

A very long program for financial applications should be written in a language designed to deal with numbers and one in which programs can be easily modified if required.

On the other hand there is nothing quite so bad as an arcade game written in BASIC - when you get down to it, it is just too slow.

Your own needs, the amount of memory in your computer, the response time required, the time available for development, and so on will determine your choice of programming language.

Thus, in summary, machine language is a series of commands which the CPU can understand and which can be represented by numbers.

What is Assembly Language?

Quite obviously if machine language could only be represented by numbers, very few people would be able to write programs in machine language.

After all, who could make sense of a program which looked like

```
0 0 1 0 0 0 0  
0 0 0 0 0 0 0  
0 1 0 0 0 0 0 0  
etc ...
```

Fortunately, we can invent a series of names for each of these numbers. Assembly Language is just such a representation of Machine Language so that it can be read by humans in a form that is easier to understand than

```
0 1 1 1 0 1 1 1.
```

There is only one difference between Assembly Language and Machine Languages: Assembly Language is one level higher than Machine Language. It is more easily read by humans than Machine Language, but on the other hand, computers can't read Assembly Language.

It is not an adaptation of Machine Language such as BASIC. For each Assembly Language instruction there is an identical (in function) Machine Language instruction, and vice versa, i.e. there is a ONE-TO-ONE relationship between them. We can therefore say that Assembly Language is EQUIVALENT to Machine Language.

Assembly Language makes use of mnemonics (or abbreviations) to enhance readability. For example at this stage, the instruction

```
INC HL
```

may not mean much to you, but at least you can read it. If you were told that 'INC' is a standard abbreviation (or mnemonic) for INCREASE and that HL is a 'variable', then by simply looking at that instruction you can get a feel for what is happening.

The same instruction in Machine Language is

```
0 0 1 0 0 0 1 1
```

Now obviously you can also "read" that instruction in the sense that you can read the number, but it isn't going to mean much to you unless you have a table to look up or when your brain is functioning almost like a computer.

Assembly language can be converted directly to machine code by a program or by you. Such a program is called an ASSEMBLER. You can see this as a program which performs the rather boring task of translating your assembly language program into a sequence of machine language instructions that the Spectrum will understand. And we understand that an ASSEMBLER for the ZX Spectrum is already available.

Nonetheless, such assemblers typically require 6K of memory, and are therefore of limited use on a 16K machine. The Spectrum display takes up 7K of memory, and after loading the Assembler you may have only 4K of memory left for your assembly language program. (This will mean about 1/2 K of machine language program).

The alternative to an Assembler program is for you to do the translation of the assembly language mnemonics into machine language by hand, using the tables provided in this book.

It's hard, it's frustrating at first, it's inconvenient, but it's wonderful practice and gives you a great insight into the way the Spectrum CPU works.

We would in fact recommend that you try writing short machine language programs in this way - ie writing them in assembly language and translating it into machine language by hand - before buying an Assembler program.



CPU

The central processing unit of the computer. This is the chip that does the controlling and calculating work in the computer.

Machine Language

The language understood by the CPU. For the Spectrum's CPU, this is the Z80 machine language which is made up of about 200 instructions¹.

BAS C language

A computer language designed to be intelligible to humans. When

but easier to write.

Assembly Language

The human shorthand representation of the machine language instructions so that each of the latter instructions can be understood more easily. For example, HALT is the assembly language equivalent of the machine language instruction 0 1 1 1 0 1 1 0

A program that translates assembly language instructions (easily read and understood by humans) into machine language (which can be understood by the computer eg the Spectrum).

Read Only Memory (ROM)

Program that has been FIRMLY built into the hardware of the computer; it will remain there even when the power is off. For the Spectrum, the ROM is in Z80 machine code, and was written specifically for it. The ROM of Spectrum occupies from memory locations 0 to 16183. You can only refer to the contents of these locations, unlike the rest of memory which you can refer to and change as desired.

BASIC Machine Language Concepts

WHAT IS THE CPU?

machine (the CPU) talks

unless we know what sort of information the CPU understands or

do things all the time.

Especially calculate

to keep track of what is happening. How does he do it?

The design of the CPU

only, but he is able to do those tasks very quickly.

We mentioned above that the CPU doesn't even have pencil and paper and that is part of the design of the CPU. Any number he can't remember or can't track of has to go in a box for safe keeping.

time in NEW YORK, knowing the time in LONDON.

it to do next, so it puts that information away in a box, say box #1.

result away, say in box #3.

$$10 - 5 = 5$$

The answer of course is 5 o'clock.

so it does exactly what you or I would do - it counts on its fingers and toes.

The CPU's hands and feet are called Registers.

hands and toes - but we will get to that later.

In the above exercise, let's call one of the CPU's hands "HAND A". Now

actually do given the above instructions

- * Count out the value of box #1 on the fingers of Hand A,
- * Subtract the contents of box #2 from what he has already on his finger.

Look at the value on the fingers of Hand A and store it in box #1.

conclusions to be drawn from this

1. The CPU would not be able to deal with a number like 11.50 - it could only deal in whole numbers.
2. The CPU would be limited in its calculations to whatever number it could count on its fingers.

This is indeed true.

The main consolation however is that the CPU has a lot of hands and feet and can count on each of them separately, and that it can count to 255 using only the 8 fingers of Hand A.

255 and each foot can be used to count to over 64,000!

Programmers have to understand what has been done is describe the processes.

The first two processes is of the example of machine language program shown as mnemonics. Abbreviations instruct the CPU at each step.

SETTING UP

LD	(BOX #1),	10	;Load box 1 with 10
LD	(BOX #2),	5	;Load box 2 with 5

WHAT AT ONS

LD	A,	(BOX #1)	;load A with box 1 contents
SLB	A,	(BOX #2)	;subtract contents of box 2

STORING THE RESULT

LD	(BOX #3),	A	;load box 3 with A value
----	-----------	---	--------------------------

These instructions may seem a little strange at first, but after all, mnemonics are mnemonic.

"LD" is an abbreviation for LOAD so that

LD A,1

for example, would mean load A with 1; that is count off 'one' on the fingers of hand A.

INTENTS OF WHATEVER IS INSIDE THE BRACKETS.

brackets do look like they are meant to indicate a container.

#1 and #2 with 10 and 5, ...etc... to get the final result of 5 in box #3.

All of this is fairly simple to follow and I am sure you can

on Hand "A" are used to represent the time in New York. A minute later they may be used to represent the number of employees in a company, and at some other time how much money you have.

If you are used to the concept of variables from your BASIC programming, you must leave that behind in machine language programming.

The fingers of Hand "A" are not a variable in the same sense as in a BASIC program. They are merely what the CPU uses to count with.

ONE OF THE BIG DIFFERENCES IN PROGRAMMING IN MACHINE LANGUAGE AND PROGRAMMING IN BASIC IS THIS LACK OF VARIABLES

You may realise that you can think of the BOXES we use to store

Yes, you are absolutely correct, but these are not variables either. They can be immensely useful and perform similar purposes to variables, but bear in mind that these boxes are no more than memory locations set aside for a specific purpose.

The way the CPU copes with negative numbers is different, and we will look into that later.

What if the CPU runs out of hands?

I should mention here that you would probably find the CPU a very strange looking fellow were you to meet him in the street.

His hands have eight fingers each, and he has eight hands! He only has two feet, but each foot has 16 toes, so he is extremely agile with his toes!

P D S C

and toes

Nonetheless, it is possible that in some cases the CPU will not

on or other the programmer will wish the CPU to
stop in the middle of calculation to do something else

P D S C

boxes it put the information in

The Z80 CPU gets away with using a stack, which is one of those

P D S C
books, spare notes, etc. I am sure you have seen those stacks where you spike one piece of paper on, and then the next one, and so on. It's a great filing system if you want the top piece of paper only, but very inconvenient if you want one in the middle because you have to riffle through all the pieces of paper on the stack.

As it happens, it's a very convenient system for the CPU because it ever only needs to look at the top piece of information

Whenever an interruption causes the CPU to stop doing its calculations, it PUSHes all the information it has on its hands onto the STACK, and as soon as the interruption is over, it POPS the top bits off, and continues with its work.

So now we have to take off the stack. To do this when we get it off, we "POP" it off.

Now to move information onto the stack we need to do the same thing, but in reverse. We need to move the information down the stack.

Information, there then needs to be many separate "POPs".

For example, if you wanted to move a stack of numbers up to keep the stack stuck to the ceiling. This means that the more numbers you want to move, the further you have to move them downwards.

information is 16 - it knows it is the last piece of information "PUSHed" on the stack. Naturally it needs to be a little bit

bit odd.

What can the CPU do?

I think it's worth considering at this stage the type of built-in functions built into the 280 chip.

Because the CPU has to be able to keep track of all its numbers the CPU can deal with

- * one handed numbers - ie numbers you can count on one hand
- * two handed numbers - ie numbers you can count off on two hands

You may find this difficult to believe, but the CPU cannot deal with numbers larger than those it can count on two hands!

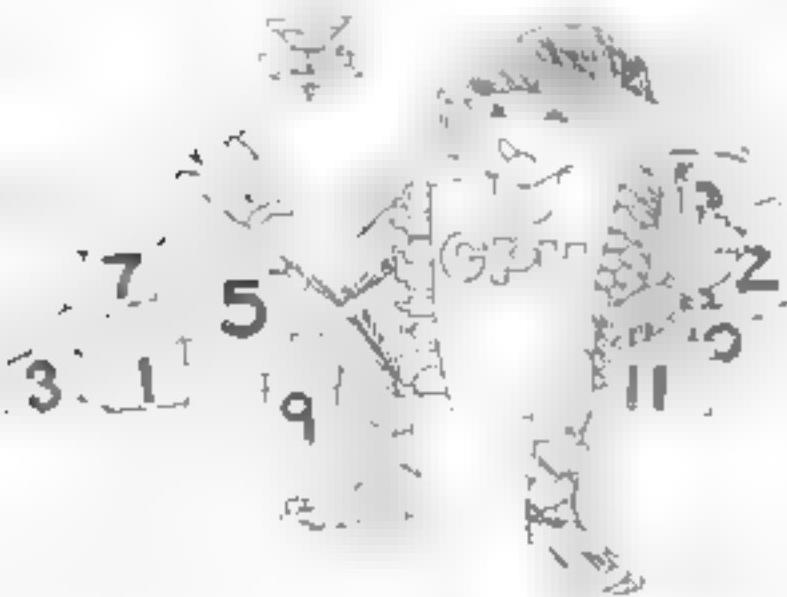
The types of instructions the CPU can perform are also very limited

- * counting off numbers on one hand
- * counting off numbers on two hands
- * adding, subtracting, increasing, decreasing, or comparing one handed numbers
- * adding, subtracting, increasing or decreasing two handed numbers
- * various manipulations on one handed numbers - eg making the number negative
- * making the CPU skip to another part of the program
- * trying to communicate one handed numbers to and from the outside world.

I am sure you will agree that this is a very limited set of
instructions to get the CPL to play chess, or to work out your wages!

have to write a program to do so

than writing programs in BAS C - you can only do things in tiny
little steps.



SUMMARY

Registers

these can be thought of as the CPU's feet. Each 'hand' has eight 'fingers', while each 'foot' has 16 'fingers'.

Memory locations

The CPU can transfer information from its hands into or from any other other hand, and into or from memory.

represent specific information.

The Stack

The CPU can use the stack to transfer information the programmer may wish to store temporarily. Information is retrieved by POPping the information off.

Four byte instructions

last type of information transfer and simple arithmetic calculations. All programs must be made up of series of these simple instruction

The Way Computers Count

As we have seen, the binary system has a major advantage over the decimal system in that it can be easily implemented in electronic components. However, it also has a major disadvantage in that it only manages to count to 10?

It is true that the binary system only counts to 10, but what does this mean? When we raise the thumb up as having your little finger raised?

With two fingers it is possible to represent different numbers in this way.

Number	0	1	2	3	4	5	6	7	8	9	10
Counting	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001	1010
Binary	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001	1010
Octal	000	001	010	011	100	101	110	111	200	201	210
Hexadecimal	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001	1010

With only two fingers it is possible to devise a way to count from 0 to 3, as follow:



00 + 0

We can indicate not having a finger raised as '0'.

1 + 1 + 1 = 3

representation 11 (or two fingers) have the value 3.

There is a direct relationship between this and binary

as we have seen, so if we can make to indicate an off (or '0' and '1') as convention dictates).

from 0 - 7

but all the possibilities for four fingers being raised

confusion in trying to write down the number eleven as opposed to indicating that two bits were set, a universal convention has been adopted

The numbers 10 - 15 are indicated by the letters A

Decimal	10	A
	11	B
	12	C
	13	D
	14	E
	15	F

This means we write the numbers from 0 to 15 decimal as
5 op's, isn't it?

This way of treating numbers is called the HEXADECIMAL FORMAT.

To prevent confusion, some people write "H" after a hexadecimal (0H). The "H" has no value, but serves to remind the hexadecimall convention.

In machine language programming, it is CONVENIENT to deal with numbers in hexadecimall format.

For your instruction in normal decimal format, it is convenient for us to use the hexadecimall format because

1. It is easy to convert from this form to binary, which tells us which bit (or finger) is doing what.
2. It gives us an easy means of seeing whether numbers are one handed or two handed - i.e. 8-bit or 16 bit
3. It standardises all numbers to sets of 2-digit numbers. (We will elaborate on this)
4. It is the common convention and familiarity with hexadecimall will allow you to read other books and manuals more easily
5. As the CPU is designed to process information represented by binary numbers which are cumbersome for humans to read, we need a representation which is more easily readable

But it is only a convention and not a sacred rule.

be represented by two hands of 5 fingers each.)

REGISTER IS THAT THIS IS THE STRUCTURE OF THE ZX SPECTRUM

on each hand.

Taking this one step at a time, let us come familiar with 4 fingers first:

left. If we number the fingers

3 2 1 0

then the value of each finger is 2 to the power N^t where N is the finger number. Let's call a 4-finger hand a "handle". (just small cigar is a cigarette)

Exercise

bits (or fingers) represent?
Decimal Hexadecimal

JO 0
1110
001
010
100

last few pages again before going on.

Say 10? We would use the next finger on the left, as

7 8 5 4 3 2 1 0

= 16 decimal = 10H (Hexadecimal)

The reason we write the number as 10H is that we divide the hand into two "4-bit handlets". We can therefore easily denote each (A-F).

In this way any 8-bit hand can be written as exactly two hexadecimal handlets



The "handlet" on the left indicates 16 times as much as the "handlet" on the right. This is much the same way as in decimal notation, the digit in the "tens" column is worth ten times as much as the digit in the "ones" column.

We convert numbers in decimal format such as 15 automatically to
15 = (1¹⁰) + 5

This is so automatic that we don't even think about it.

It is exactly the same thing in hexadecimal notation. To convert
the
hexadecimal number on the left "handlet" by 16. Using the example
above

$$10H * (1^16) + 0 \\ = 16 \text{ Decimal}$$

maximum is obtained when all fingers are held up

7	6	5	4	3	2	1	0
						$\sqrt{}$	F
						4	

$$\begin{aligned}FFH &= (F \cdot 16 \\&= (15 \cdot 16) + 15 \text{ (in decimal)} \\&= 255 \text{ (Decimal)}\end{aligned}$$

The smallest number is when no fingers are held up

Out = 0 Decimal

Note that all numbers, from the smallest to the largest, need only 2 digits to define the number.

Soon you will get the hang of it.

Also note that when you count in hexadecimal, you do the same as in decimal.

decimal: 26 27 28 29 30 etc.

Hexadecimal: 26 27 28 29 2A 2B 2C 2D
 2E 2F 30 etc

The values of the numbers in the decimal and hexadecimal series are the same, except for the first two, which are 2AH, not 30H.

The following BASIC program will enable you to input to your Spectrum a decimal number and convert it to a hexadecimal value.

```
100 REM decimal to hexadecimal conversion
110 PRINT "Please input decimal value."
120 LET T = INPUT
130 LET S$ = "" 135 LET n2 = INT (T/16)
140 LET n1 = INT (T - n2*16
150 LET S$ = CHR$ ((n) (= 9) * (n1 + 48) +
    (n1 * 16 + n1)) + $S
60 IF n2 = 0 THEN PRINT : PRINT "HEXADECIMAL =      "
    " = 1 TO 200  NEXT 1: R N
70 LET T = n2. GO TO 135
```

and use the BASIC program to test your answer.

- i. 16184 memory address of the start of Spectrum display file
- ii. 22578 memory address of the start of Spectrum attribute file
- iii. 15360 memory address of the start of Spectrum + 7 bit
- iv. 15616 address of the start of ASCII characters in Spectrum

SUMMARY

Decimal

The decimal notation is a convention of counting numbers in groups of ten units at a time. These are represented by 0, 1, 2, 4, 5, 6, 7, 8, and 9.

Hexadecimal

The hexadecimal notation is a convention of counting number in 4, 5, 6, 7, 8, 9, A, B, C, D, E, and F.

Sometimes an H is added at the end of a hexadecimal number to remind us it is written in this format. For example, 1800H

8 Bit Memory Locations

format as a two digit number.

How Information is Represented

Information is stored in a computer in groups of bits.

A bit stands for Binary digit ('0' or '1'); in the Z80A group of eight bits is called a BYTE.

The 8-bit byte is the standard unit of information in Z80 and most microcomputer CPUs talk.

Information is stored in memory in bytes. In RAM memory, each byte can store either binary data or alphanumeric text. We will thus discuss below these three types of memory.

Program Representation

A program consists of a sequence of instructions which are "sub tasks".

Instructions are represented by binary numbers consisting of two or more bytes.

With one byte at a time, and if it requires more than one, it

instructions where possible.

Instructions consist of a sequence of bytes.

Numeric Data Representation

* Integer Representation

number of fingers = number of bits = number of binary digits

numbers. Also, by using only 8 fingers (i.e. an 8-bit number), we could represent all the numbers in the range 0 to 255.

e.g. decimal 255 is represented by 0FF in
the binary

But what about negative numbers?

* Signed Integer Representation

Remember one byte is a HAND with eight fingers and a number is represented by holding different fingers up.

following convention (signed representation):

bit 7 bit 6 bit 5 bit 4 bit 3 bit 2 bit 1 bit 0
+-----+-----+-----+-----+-----+-----+-----+-----
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
highest bit - bit 7 is on.)

positive (depending on whether the thumb is up or not).

Now comes the crunch. When is a number with the thumb up a positive number and when is it a negative number?

positive if thumb is up
negative if thumb is down
particular time.

As the first word working with words was very
few days I had to be very patient in doing so.

or positive and negative.

* Choosing a Representation for negative numbers

number is negative, and not holding the thumb up means it is positive. Is this enough?

No. We need to decide which of the 127 possibilities of the remaining 3 fingers will denote -1, which one -2, and so on.

We need a representation of negative numbers, such that when a number is added to its negative we get zero. As an exercise, let's

$\mathbf{P} = \mathbf{M}^T \mathbf{W}^{-1} \mathbf{M}$ is a positive definite matrix.

will be up)

could it be?)

888888

get's tex 1.0.0.0.0.0.1. In other words, the name is a 1)

9, and converge it to seven decimal places.

number which will give us the right answer is

1111111 (FFH) in hexadecimal

To confirm this

With these results we can work out a general rule for the rank of the stable modes. The rule is as follows: If the number of the unstable modes is at the maximum, the number and arrangement of the stable modes will be such that

Let's try this rule on another number, such as 3, say.

100% 0% (SDH)

Let's add to this number to 3 and see what happens:

0 0 0 0 0 0 1 1

(carry) 0 0 0 0 0 0 0 It works!

We have found a way to represent negative numbers:
-0t => FF

The largest positive number is
0 1 1 1 1 1 1 = 7F => +27 Decimal
and the negative of this is
1 0 0 0 0 0 1 = 81 = -27 Decimal

The real test of this rule is to see if by applying the rule to a negative number we get back the positive again!

Let's try it out on -3 which we worked out above is FDH.

Number	1 1 1 1 1 0 1
Opposite	0 0 0 0 0 1 0
Add (+)	0 0 0 0 0 0 1 1 => 3

This is therefore a representation that works! We can apply it to get the negative of any number.

16-bit Negatives

Exactly the same reasoning applies to two hand numbers (16-bit numbers), except that the thumb of only one hand needs to be shown as 'ON' to indicate if the number is negative or not. (ie bit 7 of the high byte).

Convention

The computer terminology for this convention is called TWO'S
COMPLEMENT numbers at the Appendix of this book.

Remember that this is only a convention! You still have to decide at all times whether the numbers you are using are meant to designate numbers in the range 0 to 255 or numbers in the range -128 to +127.

Exercise

I If 1FF (0111 1111) is the highest positive

number which can be represented in this convention, how would you represent -128?

Alphanumeric Data Representation

perhaps

Our convention to represent characters is pretty straightforward: all characters are represented on a single hand (i.e. in an eight-bit code).

characters representant The ASCII Code and the EBCDIC Code

AS-11 stands for "American Standard Grade for Information".

EBCDIC is a variation of ASCII used by IBM.

You can find an ASCII conversion Table in the Appendix. Compare it with Fig. 14-18b.

Try this : PR NT CHR\$ 33
and you will get a "!", because "!" is
represented internally by 21H.

Variety of things

- a program instruction to the CPU
 - a number in the range 0 to 255
 - a number in the range -128 to +127
 - part of a two-handed number
 - an alphanumeric character

This is all true, and it is up to you, the programmer, to remember just what it is the CPU's hand is supposed to be holding.

SUMMARY

Memory Contents

we desire. There is no way of telling which is which just by examining the contents of a single memory location.

Programs

Program instructions are stored in memory as sequences of bytes. Some instructions require only one byte, while others require up to four bytes.

Numbers

Each memory location can be used to store either positive from 0 to 255 or -128 to +127.

Negative Numbers

A convention has been adopted that when we choose to have memory store a signed (+ or -) number, the following rule shall apply:

If bit 7 is on, the number is negative

If bit 7 is not on, the number is positive

To obtain the negative of any number, get the "2's complement" and add 1.

2's Complement

The 2's complement of any number is its opposite in binary form. Any bit that is on becomes off, and vice versa.

A Look Into the CPU

Introduction

under a licence from Zilog Inc.

0.000000266 of a second

that even if all instructions performed are the slowest one about 160,000 instructions can be performed.

A Physical View of The Brain

or in the Spectrum is a silicon chip with forty pins from 1 to 40. These pins are the communication lines

etc its clock signals from pin 6, etc

through pins 7 to 15 except pin 11. The rest of pins are control signals communication

find yourself totally lost at this stage. But no need to worry it's really to our advantage that we don't know the structure of the machine, and we don't need to know it to use its capabilities. It's just the same as with a calculator. The

other words we don't see it!). We are only interested in the

and how we can use it to our purpose.

Logical View of The Brain

Logically, the Z80 can be divided into five functional parts:

- They are
- i. the CONTROL UNIT
 - ii. the INSTRUCTION REGISTER
 - iii. the PROGRAM COUNTER
 - iv. the ARITHMETIC LOGIC UNIT
 - v. the 24 USER REGISTERS (the usable HANDS and FEET of the CPU)

* CONTROL UNIT

We can see the CONTROL UNIT as a supervisor for the CPU's processing. Its task is to time and coordinate the Input,

Output, and Memory. It also decides what task the CPU is asked to perform, whether the instructions come from the ROM program, or from your program.

* INSTRUCTION REGISTER

This is a HAND that the CPU uses to hold the current instruction that it is going to perform. The whole task which comprises a

instruction is stored somewhere in memory - either in the ROM or in

* PROGRAM COUNTER

The PROGRAM COUNTER is a HAND that holds the address of the location of the next instruction to fetch out.

* ARITHMETIC and LOGIC UNIT

This is the calculator inside the CPU. It can perform both

addition and subtraction, Incrementation (adding 1) and

fingers up or down, etc.

The ALU has a special bit called the CARRY FLAG register. This is discussed in more detail further on.

* SCR-R STER

These are the CPU's Hands and Feet, which you, the programmer, can control.

There are twenty four User registers within the Z80 microprocessor
- some are HANDS, and some are FEET

The Images we have been building up of hands, feet and boxes make the processes easy to visualise and are a good representation of what is going on, but computer bums tend to look askance if you say things like "...and then the computer shifted its information from its right hand to its left hand."

so that when faced with that situation, you will be able to say
"LD B,A

To start off with, computer bums refer to the hands and feet of the CPU as "Registers"

We mentioned earlier that the CPU has eight hands - these are called A, B, C, D, E, F,In our world, the definition of a hand is something with eight fingers.

The CPU has two feet - these are named IX and IY. The definition of a foot is anything with 16 toes!

register has only one letter in its name then it must be a hand (that is, contains 8 bits), while if it has two letters in its name then it must be a foot (that is, have 16 bits).

We will have you used to computer terminology in no time.

Actually the remaining two hands for the CPU after D, E, F,are not named "G" and "H" as one would expect but "H" and

The conventional way to represent all these registers is as follows



The reason the register pair "HL" was called "HL" instead of "LH" is that it's easier to remember than "LH". It's also easier to remember than "LH" because it's easier to remember your hands and toes. You can easily set up your fingers to

such as BC, DE, and HL.

The reason the register pair "HL" was called "HL" instead of "LH" is that it's easier to remember than "LH". It's also easier to remember than "LH" because it's easier to remember your hands and toes. You can easily set up your fingers to

such as BC, DE, and HL.

might think you meant to represent the number 73.

The "H" in "HL" stands for HI-H and the "L" stands for LOW, so there is no chance of confusion - right?

register in the other register pairs contains the high number:

B in BC
D in DE

because all the highs and lows are treated in the same order

The feet (IX and IY) also have a special name: they are called "index registers". This has a lot to do with the fact that they can be used to organise information in much the same way as a book index is organised. Alternatively, you can view them as table pointers.

OK, now that we understand the term noting the two main points:

THE ACCUMULATOR (A register)

The 8-bit single byte register is the most important register in the 8085. It is also sometimes known as the累加器 (Accumulator) when there was only a single register that could be used to 'accumulate' a result.

Now we have learned about the various registers in the 8085, it is time to focus on just the one register that is used for all calculations. This is what it means by the term 'Accumulator'.

Now, if you are a bit confused about what exactly the Accumulator is, don't worry. After all, it's a very complex topic. So, let's start with something that most people can perform some task with - addition! How would we add two numbers?

The Flags

Please note that "AF" is not usually treated as a register pair.

It is a separate register pair, which is why it is often referred to as a separate chapter.

So, what is the AF register? Well, it is a 16-bit register pair consisting of the AH and AL registers. The AH register is the high-order byte and the AL register is the low-order byte. The AH register is used for storing the high-order bytes of the results of division and multiplication operations. The AL register is used for storing the low-order bytes of the results of division and multiplication operations. The AH register is also used for storing the high-order bytes of the results of logical operations like AND, OR, XOR, etc. The AL register is also used for storing the low-order bytes of the results of logical operations like AND, OR, XOR, etc. The AH register is also used for storing the high-order bytes of the results of arithmetic operations like ADD, SUB, MUL, DIV, etc. The AL register is also used for storing the low-order bytes of the results of arithmetic operations like ADD, SUB, MUL, DIV, etc.

Maybe the RL register is the CPU's right foot?

An Alternate Register Set

also has a spare set of hands!

Not really so much a spare set of hands (all right, alternate
of work gloves).

of a hand with the number 3 counted off

on the other set of gloves!

retained there. Nor, naturally, can the glove perform any
calculations without a hand inside the glove!

information the gloves retain.

on a right hand.

The representation of all the ~~the~~ stems is now therefore

$$\begin{array}{ll} A = F & (==) \quad A' = F' \\ B = C & (xxx) \quad B' = C' \\ D = E & (==) \quad D' = E' \\ H = L & (xxx) \quad H' = L' \\ X \\ Y \end{array}$$

Note that the set of gloves you are wearing has the ~~the~~ the hand it is for, while the spare set is always indicated with the dash symbol).

only works on your HANDS, not your gloves.

The only instructions involving the alternate register set are of the "swap gloves now" type. For example

1. LD A, (Box #1) ; Load A with contents of x #1
2. EX AF, AF ; Port for exchange - i.e. swap gloves on AF with those of AF'
3. LD A,(Box #2) ;
4. EX AF,AF' ; Another exchange
5. LD A,(Box #3) ;

You will note that in the above 5 instructions there are no instructions which have specifically affected the alternate

register set. Try to work out what is happening.
Do you know what will be in register 'A' after each instruction?

For simplicity's sake, let's assume that the contents of the three boxes are as follows:

(Box #1) = 1
(Box #2) = 2
(Box #3) = 3

Then the following is what happens after each instruction

		Register A
1.	1	Not known
2.	Not known	1
3.	2	1
4.	1	2
5.	3	2

Really quite simple, isn't it?

Now we can see how the stack pointer is used to move data into memory or into memory. We will follow through this point

Even More Registers?

There are many more registers than we have shown you, but we will not be using these to any great extent.

The STACK POINTER

The STACK POINT R is another foot the CPU has (2-byte address register).

It always points to where the pile on the stack has got to. As the stack grows, it grows downward from high memory locations to low memory locations.

it every time you do a PUSH or POP.

you PUSHed on to the stack. You can be sure that this will cause your program to "CRASH".

The I Register

base address of a table of addresses for handling different responses to an interrupt, for example, Input/Output request.

However in the SPECTRUM this facility is not used. This register is involved in generating T V frame. It is unlikely you will ever have to use this register.

The R Register

The R register is the memory-refresh register. It is provided in

will disappear.

the R register thus cycles over and over from 0 to 255.

without ever worrying about refreshes, etc.

255. We will demonstrate this usage later.

user's Registers

There are eight main 8-bit registers in the (P, (A, F, B, C, D, L, I), and two 16-bit registers (IX and IY). Eight bit registers have only one letter in their name, while 16-bit registers have two letters.

Register Pairs

Six of the eight 8-bit registers can in some circumstances be used in pairs to operate on 16-bit numbers.

These are the BC, DE and HL register pairs. The name HL can serve to remind us which is the High order byte and which the Low order byte.

Preferred Registers

The Z80 CPU is designed in such a way that some 8-bit instructions can only be performed by the A reg set, while some 16-bit instructions can only be performed by the HL register pair.

Alternate Register Set

The eight main 8-bit registers can be swapped with another 'alternate' set of registers.

The values stored in the main registers are retained by the CPU while the alternate set is being used, but cannot be accessed.

Changing the reg set lets again allows us to operate on original values again.



This is all very well.

notation, and it all seems so irrelevant. It doesn't explain how you actually RUN a machine language program

the time? (When it's on). It's just that you are not aware of it. Even when you're not doing anything, just watching the screen, trying to think of what to enter as the first line of your

under the control of a machine language program.

This program is the one that is stored in the ROM chip and is referred to as 'the operating system'. For example, the part of the program that is running when you're sitting there looking at the screen does the following things

Scan the keyboard for entry

Note that no key has been pressed

Display the present screen (empty)

of the 'Interpreter' type as we have already explained. It looks at your next BASIC instruction, converts it to machine language, executes that part of the program, and then returns to interpret the next instruction.

All this stops being true when you run your own machine language program!

Total freedom from the operating system. The use of the VDU function hands over total control of the CPU to whatever commands you have placed at the VDU address. It will interpret whatever it finds there as valid machine language instructions.

This can be pretty terrifying as you could lose everything stored in memory should you lose control. One error, one wrong character, and you will have to turn the Spectrum off and start again from the beginning.

There are no error messages to catch what you have done wrong, no syntax checking for incorrect statements so if you make the slightest error, the hours of work you put in to enter your program could be lost!

At the end of this book we have included a BASIC program which will allow you to enter and edit machine language programs. Once you have entered this program on your Spectrum, save it on tape as it

language program at least once.

The worse that can happen is that you may have to turn your Spectrum off and on again.

"Editor" found at the back of this book and RUN it.

12000. Enter the number 32000 then press (ENTER).

The screen will now show

Command or Line (###):

new line of machine code,

screen should now show you all the lines you have ente

I c9

and at the bottom of the screen the prompt

Command or Line (###):

enter a command instead.

you have specified, namely 32000.

Congratulations: you have just entered a one instruction

listing address. Enter 32000 then (ENTER).

key "m" to return to the main command input stage.

What the instruction "C9" means is RETURN!

way you want to "return" to the safety of earth (or operating system as the case may be).

Now we can enter programs into memory from the keyboard followed by (ENTER).

We can now run the program that we just entered. The command used at the start.

Let's do this. The command is RUN. The value of the BC register pair will be the value of the BC register pair.

The command RUN (from the keyboard or from the same one) deals with the "USR" function.

We can now run the program that we just entered. The value of "USR", as in case 32000.

The value of "USR", as in
let A = USR 32000
naturally gave the answer 32000!

Let's do this. The command is RUN. The value of the BC register pair will be the value of the BC register pair.

Let us enter the following machine language program:

OR
19

We can now run the program that we just entered. The command "dump" and then the command "run".

This time the result will be 31999! This is because the value of the BC register pair was increased by 1).

instructions at the table at the back. Can you work out what the abbreviations mean?"

NEVER RETURN.

been damaged. Just turn off the power and reload everything.

Exercise

interesting.

How the CPU uses its Limbs

we have seen that the CPU has two main limbs - the left limb and the right limb. These limbs are used to control the memory and other components of the computer system. In fact, the CPU is like a person who has two hands - the left hand and the right hand. These hands are used to perform various tasks such as reading data from memory, writing data to memory, and performing arithmetic operations.

Imagine for a moment that you are the CPU

at with your other hand. There are also certain actions which left and right hands.

is the same in machine language - you can perform some tasks allowed is the key to success.

The equivalent hand on the CPU to your right hand is the

hand, foot and vice versa.

Computer boffins refer to this

one register to another.

Other examples would be LD A, B
LD R, E
and so on

Thus we would read

LD A, B

as "LOAD A IN TO B"

normal English sentence would be

There are also other combinations or ways other than register
another or from register to memory.

* as you can use the CPU's limb

and the possible combinations (addressing modes) that are

Let's look at the combinations offered by the 780

- * Immediate addressing
- * Register addressing
- * Register indirect addressing
- * Extended
- * Indexed addressing

What a list of names? Don't worry, just relax & confident and we will step through them one at a time.

The list above does not cover all the possible combinations possible - only those that apply to one-handed number

LD r, n

(or other instruction - we use LD as an example)
We use the abbreviation 'r' to mean any 8-bit register and 'n' any 8-bit number.

ing is a technique that involves only a single

CPU can execute the instruction IMMEDIATELY it receives the instruction. It doesn't need to look in memory to find more information in order to perform this instruction

for example, count off 215 on hand 'A'. I am sure you know enough about the mnemonics by now to be able to write this as

LD A, 215 or LD A, 017H

Once again you can do this with any of the registers, with any numbers whatsoever

The format for the immediate addressing type of instruction is shown below

byte 1 instruction (telling the computer what code is this instruction)

byte 2 n (the value of the actual
data for the instruction)

Since there is one byte allocated for the actual data, the limitation to the size of number you can specify is within the range 0 - 255. If you don't understand this, refer back to chapter on "The Way Computers Count".

We usually use immediate addressing to initialise counters and to define constants needed in calculations.

Immediate addressing is easy to use in machine language

would be

LET A = 5

entire programs this way

Immediate addressing is convenient but does not solve any major problems.

But at least we're starting to get somewhere! We can now specify which number gets loaded onto which register.

* Register dressing

We dealt with this mode briefly earlier. The general format is

LD r, r
(for other instructions)

This technique only involves two hands; in short, this is passing information from one hand to another.

the "P1" hand (which we should not think of as a hand at all. It is the 'FLAG' register and does not store numbers in the normal sense).

Register dress instructions only need one byte.

Instructions of this type are not only short (One byte), they are faster as well. The time needed to execute them is the time taken for 4 clock pulses, or less than 1 microsecond on the Spectrum.

used when possible to improve program efficiency in time and storage.

* Register Indirect addressing

LD (rr), A or LD A, (rr)
LD (H), B or LD B, (H)

This powerful type of instruction causes the transfer of data between the FPU and a memory location pointed to by the content one of the 16-bit register pairs (Feet).

Register indirect addressing is faster than ordinary indirect

however, we must load the register originally, and so register
same or neighboring address many times.

For example, LD HL,S-APE ;load HL with start of
;shape datab
LD A,(HL) ;Retrive a data
NC HL ;move pointer along
;continue until shape finished

* Extended Addressing

LD A, (nn) or LD (nn),A

Now we are looking at how to store and restore information from a
to your Hand and Feet from memory.

In Extended addressing, the instruction from the program supply the
CPU with address specified by two bytes.

If the transaction is to and from the accumulator, information
transfer will only affect the content of memory referred to by the
two-byte integer.

memory location will be affected

The format of this type of instruction is

byte 1	op code
byte 2	(possible additional op code)
byte 3	low order value of the 16-bits integer value
byte 4	high order value of that integer value

Now this is the way the program can read the memory into the us

relative address is used in the instruction
relative x when h is absolute address the instruction is referring to is relocatable

e.g. SHAPEx DB D,D,D,... ;shape data bytes

LD A,(SHAPEx) ;load first byte of shape
in accumulator

* Indexed addressing

LD r, (IX/IY + d) or LD (IX IY + d), r
(for other instructions)

This type of transaction involves a Port of the CPL, the IX or IY index register

The CPU adds the contents of the index register to the address supplied with the instruction in order to find the effective address

Another common 16-bit instruction type is the Block Load Instructions e.g. LDTR (Load increment and repeat).

One typical usage of this type of addressing technique is to perform Table operations.

of ea as. A displacement value is supplied in the instruction to point to refer to

e.g. LD IX, TABLESTART ;initialise pointer to start of table
LD A, (IX + 3) ;refer to the third byte from the start of the table

The format of instructions of this type is

byte 1	(op code)
byte 2	(op code)
byte 3	d ;displacement integer d

The number 'd' is an 8-bit number which has to be specified together with the instruction and can not be a variable.
e. the range of addressing is limited from -128 to 127 from the

address pointed to by the index register.

Indexed addressing is slower because the CPU must perform an addition in order to obtain the effective address. Yet indexed addressing is much more flexible since the same instruction can handle all the elements in an array or table.



SUMMARY

The program can be in 8-bit words and transfer it from 8-bit registers to memory.

immediate addressing

Defining in the program the number to be transferred to any register.

Register addressing

From any register to any other register.

Register indirect addressing

Either using BC or DE to specify the address, and A to hold the number to be transferred.

Or using HL to specify the address and defining the number in the program.

Extended addressing

Specifying the address in the program and using A to hold the 8-bit number.

Indexed addressing

Using IX or IY to specify the start of a table in memory, and any register to hold the 8-bit numbers. The displacement from the start of the table must be specified in the program.

The number to be transferred to memory can also be specified in the program if desired.

addressing modes are the only modes of transferring information to and from memory. No other combinations are

Instructions For One-Handed Loading Operations

Mnemonic	Bytes	Time Taken	Effect on		
			C	Z	Pv
LD Reg, Register	1	4	-	-	-
LD Register, Number	2	7	-	-	-
LD A, (Address)	3	13	-	-	-
LD (Address), A	3	13	-	-	-
LD Register, (HL)	1	7	-	#	#
LD A, (BC)	1	7	-	-	-
LD A, (DF)	1	7	-	-	-
LD (HL), Register	1	7	-	-	#
LD (BC), A	1	7	-	-	-
LD (DE), A	1	7	#	-	-
LD Register, (IX + d)	3	19	-	-	-
LD Register, (IY + d)	3	19	-	-	-
LD (IX + d), Register	3	19	-	-	-
LD (IY + d), Register	3	19	-	-	-
LD (HL), Number	2	10	-	-	-
LD (IX + d), number	4	19	-	-	-
LD (IY + d), number	4	19	#	#	-

Flags notation

indicates flag is altered by operation

0 indicates flag is set to 0

1 indicates flag is set to 1

- indicates flag is unaffected

Counting off Numbers on One Hand

Learn how to count off numbers on one's hands.

We discussed in the previous chapter some of the ways we can

register addressing.

from one register to another.

Examples are

LD A B

LD D, E

and so on.

Remember the terminology involved. 'LD' means
'load', and the mnemonic (abbreviation) instruction is
read in order as an English sentence.

We would thus read out loud something like

LD A

'Load A with B'. The next example would be read as

hand. Even the seemingly stupid instruction 'LD A, A' is
permitted.

A short shorthand of this is "LD r,r" where "r" represents any

register. We now know we can shuffle information between han-

about the mnemonics by now to be able to write this as

LD D, D7

(D7 is the hexadecimal representation of 215).

obvious, isn't it?).

Once again you can do this with any of the registers, with any number you can specify with 8 bits - D = 255.

A short shorthand of this is "LD r,n" where "r" indicates any implies 8-bits still applies.

Now we're starting to get somewhere - we can now specify which from hand to hand. But we still haven't learnt how to put any of many registers.

We showed you very briefly an example of "extern addressing" when we were doing the time difference exercise.
LD A, (Box #3)

The general mnemonic for this is
LD A, (nn)

etc.

Note two things about this:

1. You can only do it with Register A
2. You have to supply the number of the box as a two handed (8 bit) number

The REV - e instruction is also valid. This is one thing you will notice about the Z80 - there is symmetry about the instruction set:
LD (nn),A

Do notice that these instructions only apply to Register "A". There are of course other instructions for the other registers but none quite as clear as this one. It's the dominant hand concept again.

Let us pause here for a nanosecond and consider what these two instructions actually mean and do for us.

In the first place, the number range that can be defined by a two handed number (nn) is from 0 - 65,535. This is 64K, and means that 64K! This means that all the memory - ROM, program, display, and free memory - have to fit within 64K. On a "16K Spectrum" there is the "16K" refers to the RAM part only. On the "48K Spectrum", the

is available on a 48K Spectrum.

The instruction "LD A,(nn)" - which is read as "Load A with the contents of location nn" - is a very powerful instruction. It enables us to "read" the contents of any memory location, whether in ROM, or RAM.

You can use this instruction to explore to your heart's desire, even to a location where there is no memory - eg to try to see what

You will be surprised - it is not all zeros

The reverse instruction "LD (nn),A" - which is read as "Load the contents of memory location nn with A" - will attempt to write to

limits one

You can't write to a location that can't store that information, such as in non-existent memory beyond the size of your system

One of the limitations of this instruction is that we have to know EXACTLY OR WRITE INTO. The abbreviation "nn" means a definite number - eg. 17100 - and not a variable

You can't use this instruction in the machine language equivalent of a "For - Next" loop. The main use for this instruction is

eg
 ld define 32000 > speed
 32001 > height
 242 > fuel left
in a lunar lander type program

You could therefore plan a program where you get the fuel left, set it, and stored the new amount of fuel back into that location. You will know at the time of writing your program the address of that memory location which serves to act as a storehouse for that information

Let us be clear about this. Location 32002 is not a variable. It is only a memory location which you use to store information.

When writing your assembly language program you would therefore write something like
 LD A (Fuel)

Machine code for this instruction you would replace "fuel" by the hexadecinal address of the memory location you specified.

But what if we don't know the exact address of the memory location where the information we seek is? Suppose we can only calculate

where that information is going to be? Because we need 16-bits to specify the address of any memory location, we would need to store it in a 16-bit regster. This means one of the register pairs BC, DE, or HL, or one of the direct registers IX or IY.

One way we can do this is to have one of the register pair contain the address of the memory location. Because the register contains the information and because we don't have the address directly we call this form of addressing Register Indirect Addressing

The mnemonic abbreviations for these are

LD r,(HL)

LD A,(BC)

LD A,(DE)

The English reading of these instructions is

"Load the register with the contents of the memory location pointed to by HL"

"Load A with the contents of the memory location pointed to by BC"

"Load A with the contents of the memory location pointed to by D"

can load to any register - even K or L, as strange as that may sound - but that using BC or DE we can only load into the A register.

In the same way that the A register is the favoured single register.

store information into memory locations in a similar way

LD (HL),r

LD (BC),A

LD (DE),A

Now the direction of flow is in the opposite direction the information flows in.

to the memory location.

The short shorthand of these instructions is

LD r,(IX + d)

LD r,(IY + d)

"r" is again any register, and "d" is the "displacement"

(confused - we don't mean register 'D' but d = displacement)

The number "d" is a one handed number (8-bit number) which has to
This is the weakness of this particular instruction and means that
data

The symmetrical instruction is also available

LD (IX + d),r
LD (IY + d),r

don't worry you are unlikely to need it in your first few
programs

The Z80 chip used in the Sinclair computers is nothing if not
we described above.

the number you want loaded) with external addressing (ie
specifying the address to be loaded by using a register base).

This is called a surprise, surprise - 'Immediate External
Addressing' .

short hand is therefore

.LD (d),n

This is useful as you can directly load a memory location without
first having to load that value in a reg ster.

"Immediate Indexed Addressing

This is of more limited use, and the abbreviated form for the
instructions are

LD (iX + d),n
LD (IY + d),n

Using These Instructions in a Machine Language Program

Let's try to put some of these 'LD' instructions into practice.

'USR' machine language program the value of the 'USR' is the contents of BC. Let's run the following progr

the Loading address to 32000,

1	0e	00
2	c9	

Now use the DUMP command to place this code into memory.

From now on, we will no longer be giving you such explicit

understanding into the point of the program.

be showing all of our programs as follow

3E	00	LD C,0
C9		RET

& Turn) and which instructions require 2 bytes, etc. (you will

The other point is that we shall try to make all our programs

It does not matter what you specify as your loading address.

or any other loading program you may design yourself

code into memory and then use the 'run' command in the EZ Code program) what would you expect the result to be?

The program sets the "c" register in the register pair BC to zero, which is 32,000.

will be answer be A. 000D
 B. 2000C
 C. 31896

and reread the chapter on "The Way Computers Count"

Now try running the following program

```
0B 00 LD B,0
JF 00 LD C,0
C9 RET
```

Registers B and C have been set to 0).

x

number, transferring to L, setting H to 0, and so on.

Exercise

To the attribute file by the following program

```
26 58 LD H,58
2E 00 LD L,0
```

using the LD (AL),n command.

The structure of the attribute file is described in the Spectrum manual. Let us set the first character to ink red, paper white. Flash on. This is

10111010 = BAH

so the next line of the program will read

```
36 BA LD (AL),BAH
```

program, so the last line must be

```
C9 RET
```

RUN this machine language program. Did it work?

Flags and Their Uses

Flags are those nice bunting you can wave on state occasions..... - wrong!

that a certain condition exists.

indicate distress, country, piracy or whatever.

Just performed.

at the start of the last chapter a lot of instructions to be discussed in that chapter, and that chapter affected any of the flags.)

at least to understand the Flag.

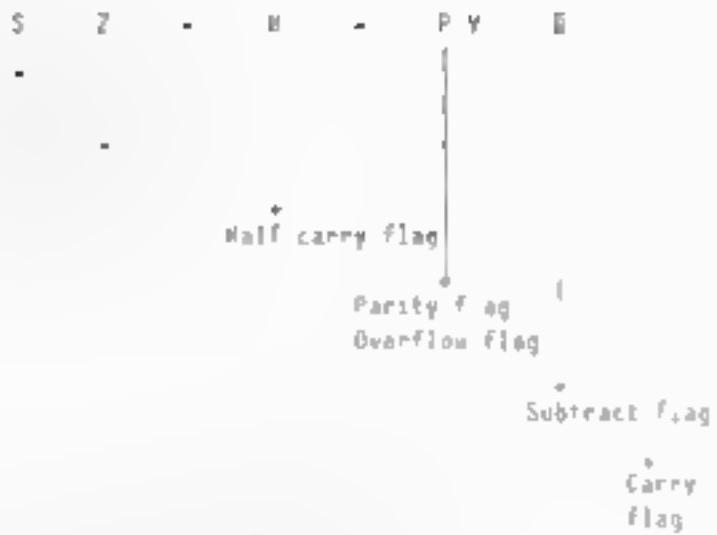
register is zero.

you would only need one bit to define the zero flag.

These flags have been computed enough so that on or off and require only one bit.

The Different Kinds Of Flags

could only think of 6 flags!



Actually the designers thought of seven flags, but decided that one register could serve as both flags - the parity/overflow flag.

Let us now look at each of these flags in detail.

Zero Flag

This is the flag we have already discussed above. Its application is obvious, and the flag is usually set after an arithmetic operation as it serves to indicate the contents of the 'A' register.

Note carefully however that it is possible to have the 'A' register contain 0 and for the zero flag not be set. This could easily happen by us as the

$D \leftarrow A, 0$

instruction. We mentioned above that none of the one-handed 65-bit load instructions have any effect on any of the flags. The zero flag would NOT be set yet A would contain zero.

The zero flag is also set if the result of the "rotate and shift" series of instructions results in a zero.

As well, the zero flag is the only visible result of some testing instructions, such as the "bit testing" group of instructions. In those cases the zero flag is put on if the bit tested is zero.

Sign Flag

The sign flag is very similar to the zero flag and operates on very much the same set of instructions (with the major point of departure being the "bit testing" group where the concept of a

negative bit is somewhat meaningless in any case.

Carry Flag

This is one of the more important flags available to language, for without it the results of assembly language arithmetic would be totally meaningless.

The point to remember is that assembly language instructions always refer to each one-handed 8-bit, or two-handed (16-bit) numbers.

This means that the numbers we are dealing with can be either:

$$\begin{array}{ll} 0 = 255 \\ + 0 = 65535 \end{array}$$

include carry,

$$\begin{array}{ll} \rightarrow 0 = 256 \\ \leftarrow 0 = 65536 \end{array}$$

$$\begin{array}{l} 255 \\ + 255 \end{array}$$

$$= 255 + 1$$

This is a direct consequence of only having a limited number range available, and the same thing can obviously happen with 16-bit number.

We've already discussed that you can only count to 255 on one hand what happens if a register is already showing 255 and you add 1? You might like to think of the register as operating the same way as your car. Once you have reached the maximum, it 'clocks' over and begins counting from zero again.

In the same way, if the register or car meter shows all zeros, and you turn it backwards, you will get the highest value showing, or 255 on an 8-bit register.

This is why the result of $200 - 201$ gives 255. If we were car drivers we would obviously like an indication that the meter has clocked over, whether in a forward direction - in which case the car has travelled further than it seems - or a backwards direction - in which case the meter has been tampered with.

is called the carry flag. Fortunately we do not need to worry about registers being tampered with.

We have seen that the carry flag can be set by subtraction operations if there should be an 'overflow'.

It is therefore convenient to think of the carry bit as the 9th bit of the 'A' register.

Number	Carry bit	Number in bit form
112	-	
* 113	-	
263	1	

But as we do not have 9 bits, the 'A' register would contain the number C8H (Decimal 204) and the carry would be on (i.e. = 1).

Leave a '1' there as well.

using Flags in the

Machine Language Equivalents

such as

```
if A>0 then  
    what follows can be 'Let...',  
        or 'Goto...',  
        or 'Gosub...'
```

exactly the same kind of decision can be programmed in machine language (except for the 'Let...'). Instead of saying "If A>0", we only look at the zero flag - if it is on, then we know A=0.

the only ones which allow us to make a choice in the next instruction to be executed

The format of such instruction is as follows:
For example,

JP cc, End

where 'JP' is the mnemonic for 'jump' and 'end' is a convenient label.

The instruction is read in English as "jump on condition cc to End".

The condition "cc" could be any of

- > Zero
- Not zero
- Positive
- Minus
- (Carry set)
- No carry

The other three flags tend not to be of so much use in every day programming. They are

Part or Overflow Flag

This flag acts as the parity flag for some instructions, and as the overflow flag on others, but there is rarely any confusion as the two types of operations do not commonly occur together.

The parity side of it comes into effect during logical operations and is set if there is an even number of set bits in the result we deal with this in greater detail in the chapter on logical operations.

The overflow is a warning device that tells you that the arithmetic operation you have just performed may not fit into the 8-bits. Rather than actually telling you that the result needed a 9th bit, this tells you that the 8th bit changed as a result of the operation.

In the example above, adding 132 and 175, the 8th bit was '1' prior to the addition and '0' afterwards, so that the overflow would have been set. But the overflow would also be set by adding

Subtraction Flag

This flag is set if the last operation was a subtraction

Half-Carry Flag

This flag is set in a manner similar to the carry flag but only in the case of an overflow or borrow from the 5th bit instead of from the 9th bit.

Both the subtract flag and the half-carry flag are of use only in "binary coded decimal" arithmetic, and we deal with these flags in the chapter on "BCD Arithmetic".

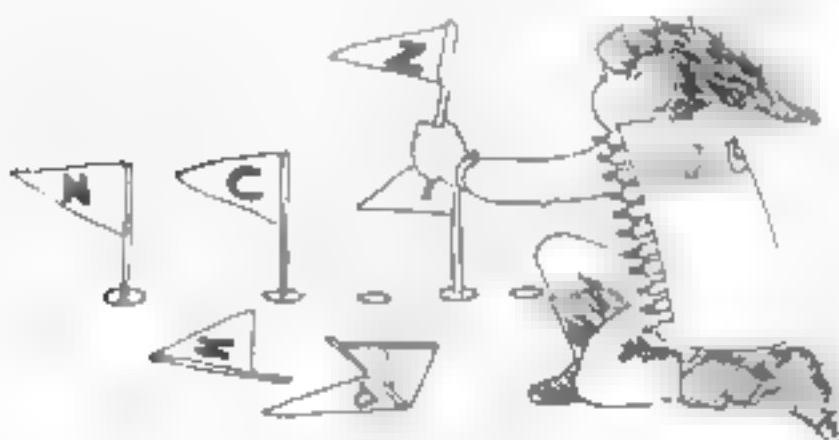
Flags are used by the CPU to indicate certain conditions after

There are six such flags, each of which can be set to b
FF. The six bits representing these flags
bits in the P register. The other two bits

The conditions indicated by the various flags are

- arity or Overflow
- Sign
- Negative
- All Carry

Not all instructions affect each flag. Some affect all flags,
only specific flags, while others have no effect on the flags.



Counting Up and Down

numbers onto its fingers and toes

fingers or we can decrease the number represented.

essentially: whatever number you have on your fingers, increase it by one.

monitoring the traffic at a particular intersection

Counting Up

mnemonic

INC E

"INC" is read in English as "increasing" and is therefore fairly self-explanatory.

feet, as we saw).

This increasing of the count on our toes is written

INC RT

NC IX

NC IY

using 8-bit numbers and which are 16-bit numbers.

The 8-bit numbers are denoted by a single letter, while

The 16-bit numbers are denoted by two letters.

index registers or the 'favoured register pair', HL

INC (IX + d)
INC (IY + d)
INC (HL)

(where d is the displacement - not the register D!)

Important note

Remember carefully our convention of reading brackets

brackets → mean → 'contents of'

This is very important - there is a lot of similarity between the instructions

INC HL
INC (HL)

but a world of difference in their execution.

first would be read as "increase the contents of HL" - one would be read as "increase the contents of the location whose address is HL" (the second reading is often abbreviated to contents of HL.)

As long as you remember the rules of the mnemonic - you will be saved from this kind of confusion. Let us examine how each works, and let's assume that HL = 5800H.

Look at HL. Increase the count on its fingers by one. Result:
HL = 5801H

INC (HL): Look at HL. Find the memory location referred to by this number. Increase the count in that location by one. Result:

HL = 5801H
(5800H) = (5800H) + 1

These are significantly different operations. (You might like to P/N both versions - 5FH is the start of the attribute file). Note also that while 'INC aL' is an instruction acting on a 16-bit number, 'INC (IY)' is an instruction which acts on an 8-bit number only - the number stored in location 5810H.

Decreasing the Count

It is equally important to

certainly ensure that everything you can increase you can also decrease, and this is indeed the case.

DEC t
DEC rr
DEC IX
DEC IY
DEC (HL)
DEC (IX + d)
DEC (IY + d)

Effect on Flags

Since the effects of these instructions are very similar to those of the add and subtract instructions, it would be a good place to review the operation of the flags.

Most of the time, the carry and zero flags will be set or cleared by the same operations that set or clear the sign flag. However, some operations, such as divide, do not affect the carry or zero flags.

Sign: This flag will be set (=1) if bit 7 of the 8-bit result is 1.

This means it will be on if the thumb is up using our previous analogy. Note that this will happen whichever convention you are using for the numbers.

Zero: This flag will be set (=1) if the 8-bit result is zero.

Overflow: This flag will be set (=1) if the contents of bit 7 of the 8-bit number is changed by the operation.

Negative: This flag will be set (=1) if there is a carry into or a borrow from bit 4 of the 8-bit number.

Subtract: This flag is set if the last instruction was a subtraction. Thus it is not set (=0) for "INC" and set (=1) for "DEC".

Use the "LD", "INC" and "DEC" group of instructions to return numbers you want as a result of the 'USR' operation.

This will give you familiarity with these instructions.

increase or decrease the contents in any of the 8-bit registers or in any of the 16-bit register pairs or in either of the 16-bit indexing registers

We can also increase or decrease the contents of memory locations whose address is specified by the Hu register pair or by the indexing registers

or in memory, affect all the flags except the carry flag.



Instructions for One-Handed Arithmetical Operations

Mnemonic	Bytes	Time Taken	Effect on FL		
			C	Z	PV
ADD A, register	#	H	#	*	0
ADD A, number	#	H	#	*	0
ADD A, (HL)	#	C	#	#	0
ADD A, (IX + d)	#	Z	#	#	0
ADD A, (IY + d)	#	E	#	#	0
ADC A, register	#	#	#	#	0
ADC A, number	#	#	#	#	0
ADC A, (HL)	#	X	#	#	0
ADC A, (IX + d)	#	#	#	#	0
ADC A, (IY + d)	#	H	#	#	0
SUB register	#	H	#	#	Y
SUB number	#	H	#	#	Y
SUB (HL)	#	X	#	#	Y
SUB (IX + d)	#	H	#	#	Y
SUB (IY + d)	#	H	#	#	Y
SBC A, register	#	#			
SBC A, number	#	#			
SBC A, (HL)	#	#			
SBC A, (IX + d)	#	#			
SBC A, (IY + d)	#	#			
CP register	#	#			
CP number	#	#			
CP (FL)	#	#			
CP (IX + d)	#	#			
CP (IY + d)	#	#			

Flags notation

- Indicates flag is altered by operation
- Indicates flag is set to 0
- Indicates flag is set to 1
- Indicates flag is unaffected

One Handed Arithmetic

It should be noted that our microprocessor has two hands or fingers for performing arithmetic. It is the right hand that must be carried out through our dominant hand, register A.

Even though we are dominant and knows how to add up to his fingers, it would be foolish to assume that is the case that the abbreviation 'A' is even omitted in some mnemonics. For example, in a PDP 11, we would normally expect to see
SLB A,B
but in fact the mnemonic is
SLB P

So what does this mean? Well, if we have a 16-bit number in the A register, we can add it to whatever number we have in the B register.

A + A	Add any single register to A
A + n	Add any 8-bit number to A
A + A _x + d	A + 8-bit number in the box whose address is given by H _x
ADD A, (IX + d)	A + 8-bit number in the box whose address is given by IX + d
ADD A, (IY + d)	A + 8-bit number in the box whose address is given by IY + d

So, if we want to multiply several registers in one instruction, we will need to define each of them in memory and then add them all together and we will get any sum we want in the final location.

The one that is missing is

ADD A, (nn)

where we define the address in the course of the program.

As far as the only way to set such an instruction would be to write

LD HL,nn
ADD A, (HL)

We also note the fact that the LD HL,nn command is almost always followed immediately by the BC or DE reg's set to zero.

This is because the BC and DE registers are 8-bit registers which can only contain values from 0 to 255 as we have already seen.

For example, LD A,80H

ADD A, 81H

We expect the result to be in A' but the bit ~~overflow~~ is set to indicate the result did not fit in.

It makes sense to store the overflow bit in the carry register. Convert the numbers to decimal and check the addition:

Hexadecimal addition and subtraction is the same as ordinary arithmetic:

1 = 2
2 = 3
3 = 4
4 = 5
5 = 6
6 = 7
7 = 8
8 = 9
9 = A
A = B
B = C
C = D
D = E
E = F
F = G

So we get a carry value of hex F, which is the next number after a number bigger than 'F' instead of '9' as in decimal arithmetic.

The addition will be carried out in the following manner, to be as follows:

$$\begin{array}{r} \text{A} \\ + \text{B} \\ \hline \end{array}$$

$$1011\text{H} \quad \text{as } 8 + 8 = 16 \Rightarrow 100$$

What can you do about this carry error?

There is a special instruction "ADC" (ADD WITH CARRY), which adds the number in A to the number in A' plus the bit in the carry register into the Carry.

This is a very useful instruction: "ADC", which we read as "ADD WITH CARRY".

It adds the number in A to the number in A', plus the bit in the carry register. It is similar to ADD, except that the carry is added on (if it is set).

It is a chaining operation:

result in B:

```
LD A,E8H ;Lower part of 1st no.  
ADD A,DOH ;Lower part of 2nd no.  
LD C,A ,Store result in C
```

LD A,07H higher part of 1st no.
ADC A,07H higher part of 2nd no
LD B,A Store result in B

After the first addition (ES + DO), we will have the carry set
(check this for yourselves!)

carry

farther than in a register pair.

8-BIT SUBTRACTION

This is exactly the same 8-bit addition. Two sets of commands exist, one for ordinary subtraction, and one for subtraction with CARRY

SUB B,A Subtract B
SBC B,A Subtract B with carry

The notation '1st' is meant to denote the same range of possible operands as for the add instruction.

COMPARING TWO 8-BIT NUMBERS

exactly what it is we mean when we compare two numbers

same - they are "equal". One way to denote this in an arithmetical

subtraction, if the result is zero, then the new number would be negative.

Similarly if the new number is smaller, then the difference would be positive.

LD A,5	Number we have
5 IN H	Number being compared

Then we will have the following results -

If $H = S$ Zero flag set, carry flag not set

If $H < S$ Zero flag not set, carry not set

If $H > S$ Zero flag not set, carry flag set

It is therefore clear that the test for equality will be the zero flag, and the test for "greater than" will be the carry flag. (The test for "less than" is the absence of set flags).

Flags which were set or reset by the operation may have been altered by the operation.

are the same as for addition

"Compare" is exactly the same as "subtract" except that the contents of A^+ are unchanged. The only effect is therefore on the flags.

Eight bit arithmetic on the Z80 is limited to

- addition
- subtraction
- comparison

and can only be performed through the A register.

Given this limitation however, a wide range of addressing modes exist

i A R M S W C

affected by arithmetical operations. We can use this as a warning of overflow.

A R M S W C
i d d d d d

Instructions for Logical Operators

Mnemonic	Bytes	Time Taken	Effect on Flags					
			C	Z	PV	S	N	H
AND Register	1	4	0	#	#	#	0	1
							0	0
AND (HL)	1	7	0	#	#	#	0	1
AND (rA + d)	2	19	0	#	#	#	0	1
AND (rY + d)	3	19	0	#	#	#	0	1
OR Register	1	4	0	#	#	#	0	0
							0	0
OR (HL)	1	7	0	#	#	#	0	0
OR (rX + d)	3	19	0	#	#	#	0	0
OR (rY + d)	3	19	0	#	#	#	0	0
XOR Register	1	4	0	#	#	#	0	0
XOR Number	2	7	0	#	#	#	0	0
XOR (HL)	1	7	0	#	#	#	0	0
XOR (rX + d)	3	19	0	#	#	#	0	0
XOR (rY + d)	3	19	0	#	#	#	0	0

Flags Notation

altered by operation
 set to 0
 1 to 1
 F unaffected

Logical Operators

There are 3 logical operators which are used to test if one bit or many bits are set. In C/C++ programming, these may be used in conjunction with ordinary arithmetic.

These are represented by ASCII characters and are:

AND
OR
XOR

We will start with AND, except for one small difference. In order to do the AND operation, the bits must be ANDed together, i.e. on the individual bits of the number (or fingers of the CPU's hand).

Let us look at one of these operations, 'AND':

Btt A	Bit B	Result of Bit A 'AND' Bit B
0	0	0
1	0	0
0	1	0
1	1	1

So, what does this mean? It means that the result is 1 only if A and B both contained a '1'.

Now, how many bits would we need to store the result of ANDing two numbers?

Is it possible to do this with just 1 bit? If so, what is the operation?

The answer is no. To see why, consider the following example. Suppose we have the number 5. We want to invert the 4th bit from the right. In other words, we want to change the 4th bit from 0 to 1. In order to do this, we must know the value of the 4th bit. If we don't know this information, the number would be at least 8.

e.g. 0 0 0 0 0 1 0 1 = 5
(_____)
These bits must be '0'.

We therefore say a number whose value we do not know and apply

lies in the range 0 - 7
 e.g. $0110\ 1001 = 105$
 $0000\ 0111 = 7 = \text{Mask}$
 result of AND $0000\ 0001 = 1 \Rightarrow$ in
 range 0 - 7

Note that the result of the AND operation is placed in the 'A' register. 'A' can be 'AND'ed with an 8-bit constant or with memory locations (HL) or with (IX+d).

AND 7 Note that as only the 'A' register
 AND E can be acted on, it need not be
 AND (HL) mentioned in the instruction.

(is true for the other Boolean operations, 'OR' and 'XOR').

The 'OR' operation is very similar in concept to the 'AND' operation:

Bit A	Bit B	Bit A 'OR' Bit B
0	0	0
0	1	1
1	0	1
1	1	1

It is obvious that the result of an 'OR' operation is to give us a '1' if either A or B contained a '1'.

An example of using AND and OR in assembly language:
 LD A,Number
 OR 1 make number = 1

The above two lines would be a typical assembly listing

The result of 'XOR' is a '1' only if one of A or B contains a '1'.

all cases except when both A and B contain a '1'.

Bit A	Bit B	Bit A 'XOR' Bit B
0	0	0
1	0	1
0	1	1
1	1	0

have on the flags.

Zero Flag	This flag will be on (=1) if the result is zero.
Sign Flag	This flag will be on (=1) if bit 7 of result is set.
Carry Flag	Flag will be off (=0) after 'AND', 'OR', 'XOR' i.e. carry will be reset.
Parity Flag Note that this flag also doubles as overflow flag)	This flag will be on (=1) if there is even no. of bits in the result $0\ 1\ 1\ 0\ \ 1\ 1\ 1\ 0 \Rightarrow \text{OFF}$ $0\ 1\ 1\ 0\ \ 1\ 0\ 1\ 0 \Rightarrow \text{ON}.$
Half-Carry Flag)	Both flags turn off (=0) after 'AND', 'OR', 'XOR'.
Subtract Flag	These flags are useful if 'BCD' arithmetic is being used.

Use of Boolean Operations on R

There is a special case of the Boolean operators which is very useful - that of the register A operating on itself.

A is unchanged, carry flag cleared

A is unchanged, carry flag cleared

A is set to 0, carry flag cleared.

byte to do what might otherwise require two, such as LD A,0.

The carry flag often needs to be cleared - eg. as a matter of routine before using any of the arithmetic operations such as

ADC Add with carry

SBC Subtract with carry,

and this can easily be done by the instruction AND A without affecting the contents of any of the registers.

SUMMARY

There are three logical operators which are useful in machine language

AN

R

XOR

be stored in the A register. The result of the operation is returned in the A register.

different to its meaning as a BASIC instruction.

Individual bits.

Coping with Two Handed Numbers

So far we have been dealing only with one handed 8-bit numbers, but we have raised about the fact that the 8086 can quite hand two-handed (16-bit) numbers in some cases.

One case we have already mentioned is the index registers. These "feet" have 16 "toes" (16 bits), and can only handle 16-bit numbers.

As we work with two-handed register pairs, we need to know how to use 16-bit numbers and these hands having together "register pairs". They are BC, DE, and HL.

The 8086 can handle numbers in many ways, but it is not always appropriate or convenient. We need to understand what is available, what is appropriate, what is good, what is bad, and the way we handle them is slow and limited.

Let's look at the instruction set (or assembly language, or mnemonics) available for dealing with 16-bit numbers.

Immediate Extended addressing

 LD tr, nn
(or other instruction)

This is the equivalent of 8-bit immediate addressing. It is merely immediate addressing extended to 16-bit immediate values for transfers.

A problem with this is that pointers to 16-bit numbers are typically more than twice as big as 8-bit pointers. For example, while 8-bit memory addressing requires only 2 bytes, 16-bit addressing requires 4 bytes. A 16-bit number, being expressed in 4 bytes, is 16-bit = 4 bytes = 32 bits.

The most common form of Extended Addressing is as follows:

Byte 1	Instruction
Byte 2	n1
Byte 3	n2

Low order byte of the number
High order byte of the number.

We use this type of addressing to no avail today on the 8086s of today's parts. For example, a pointer to memory is 32 bits.

Register addressing

The Z80 has 8-bit registers. These are used to hold data, addresses and control words.

the registers

are limited in the register combinations allowed.

e.g. ADD Hu, BC

is not allowed because the BC register pair is not a valid source register.

Later chapters

Register indirect addressing

The Z80 has 16-bit memory locations. The address of a memory location is held by a register pair.

In the Z80, this type of addressing is again mainly applied using the register pair HL.

JP (HL)

Register indirect addressing is similar in concept to register indirect addressing, except that the value you want is not held in a register pair, but in a pair of memory locations.

e.g. LD HL, (nn)

where nn must be specified at the program stage.

Exercises.

5.8. Write a program to add two numbers in memory.

1. Immediate extended addressing

```
010FOO LD BC,15 ;load BC with value 15
C9 RET ;return
```

When this program is run, you will see that the value of USR on the screen is 15. This is because the immediate value 15 is added to the initial value of the number in the program.

2. Register addressing

We will now add a line to the program above:

```
210040 LD HL,4000H ;load HL with 16384
010FOO LD BC,15 ;load BC with 15
9 ADD HL,BC ;add the two numbers
C9 RET ;return
```

If you run this program, you will still get the same output as above, namely 15! Why? Didn't we add 16384?

The reason is that the first part of the program was never run, so we didn't see any of it! To see what happened, we add a few lines to the code:

3. Extended addressing

```
010FOO LD BC,15
22647D LD (7D64H),HL ;put HL in 32100 and 32101
E04B647D LD BC,(7D64H) ;get value of BC from
;32100 and 32101
C9 RET
```

This part of the program, reading information from HL into BC, would normally be run at the beginning of the P1FH and P1F programs, and more often than not. But here we illustrate what happens if we run this over time - the limited addressing mode of the Z80.

command to check on this program as well.



Manipulating Numbers with Two Hands

In the last chapter we have seen just how big or the numbers the Z80 can handle in one hand and we have seen just how small it can handle two-handed numbers.

The Z80's math is not very fast, so it is important to note why you might want to use a large number with more than one hand. Why then bother with two-handed numbers?

This is where we find out just how impossible it is to do anything useful with such a small number of bits. If we were limited to 8 bits, how could we add? The 8-bit number is our computer would indeed be a very limited machine.

To make up the example of adding two numbers by specifying the address of a memory location, we implied that such a memory location contained two words, each 8 bits long, such as LD A,(HL).

The reason for this is that we want to add two 8-bit numbers. Each of them is a byte, or part of a larger number.

For example, if we had some bytes in memory, say 1000H, 1001H, 1002H, 1003H, we would have to manipulate those numbers in memory to add them together. This is what the next chapter will deal with: 16-bit arithmetic.

Specifying Addresses with 16-Bit Numbers

Please note that all addresses must be specified by a 16-bit number.

You just can't specify an address with only 8-bits, even if it is only addresses from 0 to 255. The way the CPU works, it's not an address unless it is 2 bytes of 8 bits each.

We implied this when we used the short shorthand of

LD A, (nn)

As we have said, 8-bit numbers are stored in register pairs. The number first, the high nibble, our chapter on "A look into the CPU" shows - "HL" stands for H = "high"; L = "low".

Storing 16-Bit Numbers in Memory

= xxxx xx xx

There is one facet of Z80 design which is very difficult to explain or justify:

Instructions For Two-Handed Loading Operations

Flask Notation

- # Indicates flag is altered by operation
0 Indicates flag is set to 0
1 Indicates flag is set to 1
- Indicates flag is unaffected

from that of register pairs is used.

The low bit is always stored first in memory[†]

memory v r w f > m r s < t

Before:	Location	Contents
H L	32000	00
01 02	32001	00
	32002	00

memory locations are all empty.

After:	Location	Contents
H L	32000	02
01 02	32001	01
	32002	00

(listing) is that the low bit is always stored first.

Give with it.

Read carefully and make sure that you are
using most important source of errors in programs!

In registers: High bit stored first
In memory and programs: Low bit stored first

we have to pay.

contents of memory using the "mem" command.

loading 16-Bit Numbers

LD rr, nn

using the notation of 2 letters to indicate a
"rr" means any register pair, "nn" any 16-bit

tables at the back of the book - then the discussion we had on the
order of the 16-bit numbers in memory becomes crucial.

Even if you do have an assembler, you should be aware of these
reversals of order to enable you to "read" the code when ~~copy~~
into memory.

Let us look at a specific example.

Load HL with 258
the mnemonic for this is
LD HL,D102H

XX XX

To see that, the number D102H needs to be inserted in place of
XX XX. But because of the reversal rule, we do not enter
this as 1002H.

The proper instruction is therefore:

21 02 01
In our examples we will show you this as

21 02 01 LD HL, D102H { = 258 }

your own programs.

Other 16-Bit Load Instructions

LD IX, nn
LD IY, nn

location).

The general instructions are

LD (nn), RR

LD (nn), IX

LD (nn), IY

Put the value in memory location nn into the register pair RR.

into the register pair. It is not necessary to specify both

16-bit operations.

For example, if we want to put whatever is in memory locations

21653, (nn)

21654, (nn)

21655, (nn)

and 21656,

In BASIC we can determine this by using the line

PRINT PEEK 21653 + 256 * PEEK 21654

(21653 = SC65H)

ED 4B 65 3C LD BC, (21653
C9 RET

THE OTHER WAY AROUND,

this once in each program.

recall, the value of USR is the contents of the BC register pair when the machine language program has finished.

Note that "LD BC, (NN)" is a four-byte instruction

two byte variables listed in the Spectrum manual on pages 173 - 176.



SUMMARY

memory location where the 16-bit number is to be found.

the register pairs of from the index registers.

in program instructions involving 16-bit numbers)
The low byte is always stored first !!!

Instructions for Stack Operations

Mnemonic	Bytes	Time Taken	Effect on flags					
			C	Z	PV	S	N	H
PUSH Reg pair	1	11	-	-	-	-	-	-
P SH IX or IY	2	15	-	-	-	-	-	-
POP Reg pair	1	10	-	-	-	-	-	-
POP IX or IY	2	14	-	-	-	-	-	-
LD SP, Address	3	10	-	-	-	-	-	-
LD SP, (Address)	3	20	-	-	-	+	+	+
LD SP, HL	1	6	-	-	-	-	-	-
LD SP, IX or IY	2	10	-	-	-	-	-	-

Flags Notation

- # Indicates flag is altered by operation
- 0 Indicates flag is set to 0
- 1 Indicates flag is set to 1
- Indicates flag is unaffected



Manipulating the Stack

You may recall the image we developed in the beginning of the book of the stack as being where the CPU was able to keep information without having to remember the address of that particular information.

One of the advantages provided by the architecture of the stack operations is that we can only PUSH and POP values on to two-handed stack items. This is because the stack is primarily designed to remember addresses and we need to specify addresses as 16-bit numbers.

The general instructions for pushing or removing from the stack are:

- PUSH rr
- PUSH IX
- PUSH LY

The general instructions for popping or removing back from the stack are:

- POP rr
- POP IX
- POP LY

This is a very simple example to try, and you will notice the lack of need to specify an address.

For the ordinary register pairs = i.e. not the index registers - these instructions are only a single byte long and therefore very economical in terms of programming space.

However, there are some side effects. In the first place, if you do a POP, it means the same 16-bit value is after the POP as before.

So, for example, if we do PUSH any register pair and then do a POP afterwards, the register value POPed won't be the same as the one you PUSHed.

For example:

```
PUSH BC  
POP HL
```

The effect of the two commands is to leave the contents of the B register unchanged but set the H register to whatever the previous value of the B register was at the time of the PUSH instruction.

This effectively adds an instruction of the type

```
LD rr, rr'
```

to the 68000 code group which was ensured as you might

A 16-bit P-P instruction for register pairs is only one byte long, the cost in terms of memory is not expensive.

The other extra benefit is that we are able to PUSH or POP the
current value of AF. This is one of the few instructions where AF is
not considered a memory location but it is oh so easily sensible because
the processor will always preserve the contents
the flags.

Flags unchanged

Moving the Stack Around

are PUSHed to or POPped from.

have 16K of memory or whether you have 48K

but not in any of the LOAD, etc. Instructions because it

and will add with this help to the
people's understanding and the support of
the cause that is available.

LD SP, nn
LD SP, {nn}
LD SP, IX
LD SP, IV

Looking to the last 30 - 60 bytes before RAMD0

3. **AC** **DC** **AMPS** **HZ** **WATTS** **BTU** **T**
3. **AC** **DC** **AMPS** **HZ** **WATTS** **BTU** **T**

requires on the stack and changes will cause it to bomb out.

Stack pointer manipulation is dangerous.
stack pointer unless you are sure of what you are doing.

Note

Arithmatic calculations involving pointers must be done carefully.
miscalculation may lead to strange results

Example

When a user subroutine is called from by 'PUSH'ing the value out of the

```
C1      POP BC    ;Get address in BC
C5      PUSH BC   ;Put it back on the stack
9       RET
```

Instructions for Two-Handed Arithmetic

Flags Notation

- Indicates flag is altered by operation
- Indicates flag is set to 0
- Indicates flag is set to 1
- Indicates flag is unaffected
- Indica effect is not known

Two Fisted Arithmetic

The 16 bit bus is not being able to have 6 bit capacity so what we do is we take the 8 bit processor as the we can use 8 bit to do addition (problem is that the 8 bit integer part is stored 2's complement so the negative numbers are to be permitted).

Now the question is why in some processor it is not the case that the 8 bit part is the same as the 16 bit part and the range is -32768 to +32767.

But the thing is that if you want to add with two 8 bit numbers then the result will be 16 bit number so it is very complicated to handle 16 bit. The range of options is just not there!

Favoured Register : L

In the same way that the 'A' register is the favoured register in 8 bit arithmetic, so there is a favoured register in 16 bit arithmetic, and it is the HL register pair.

For example if we want to add two 16 bit numbers we do not omit the name of the register pair.

Addition

The additions are quite straightforward

```
ADD HL,BH  
ADD HL,DP  
ADD HL,HL  
ADD HL,SP
```

But that is it?

No, there is one more possibility to add in the 16 bit number which is ADD HL, another pair of 8 bit registers. To profit in this kind of addition we need to

```
LD DE,nm  
ADD HL,DE
```

What is the problem? That has now us up to 3 of the 8 bit registers. You will see this is not something you want to do too often.

No, also by there is no addition between HL and other 8 bit registers. You will also remember that there is no LOAD instruction which permits you to transfer the contents of TX to Y

P SH IX
POP DF
ADD HL,DE

The one point of note is the 'SP' register - the stack pointer.
It is a proper register, but obviously you can't use it as a variable!
The contents of 'SP' at will

Effect on Flags

that the 'add' instruction is not a subtraction!

by the calculation.

Add With Carry:

Because of the limited nature of 16-bits, we are able to chain

'ADC' - operates in a similar manner to 'ADD'
range of register pairs.

HL,BC

HL,DF

HL,SP

6 Bit Subtraction

to clear the carry flag before any subtraction operation.

SBC HL,BL

Here SP and the result negative will be the amount free memory available (it's up to you to write a simple program to do that - see the end of the chapter to confirm your solution).

Effect Of Carry Arithmetic On Flags:

Any time you add two numbers together the carry flag is set by the simple 16-bit addition instructions.

Carry flag is set if the sum of the two numbers is greater than 65535. It is reset if the sum is less than or equal to 65535. These flags are set according to the result of the operation.

Index Register Arithmetic

The index register arithmetic instructions are:

HL = HL + DE
HL = HL - DE
Registers is extremely limited:

- Adding the 'BC' or 'DE' register pair
- Adding the index register to itself
- Adding the stack pointer.

Solution to Memory Limit Exercise

...
and 21654 in the Spectrum manual.

Obviously if we do `LD A,(STKPNR)` with the contents of that location we are halfway there.

`D HL,(STKPNR)`
then subtract the 'stack pointer' (`SBC HL,SP`)

Because of the 'carry' we need to clear the carry flag. This is most easily achieved by the 'AND A' instruction, which is covered earlier in the book (p77).

`AND A`
`SBC HL,SP`

I don't know how to do it. One-quarter marks if you forgot all about the carry.

negative.

thus). We first want to shift HL to BC, but there is no 'load' command. So we can do this:

```
pop
```

```
PUSH HL
```

```
POP BC
```

HL still has the same information as before, so H =

again

```
AND A
```

```
SBC Hw,BC
```

```
SBC HL,BC
```

i.e., the positive number of bytes left.

to get a result from the 'JSR' function. To get HL back into BC

```
push HL
```

```
POP BC.
```

and finally a return from the USR function:

```
KFI
```

Did you get this right?
Notice how handy the stack is



Loops and Jumps

Type and run a short BASIC program at your computer. Notice how it reads the code and executes it line by line. It's really getting places.

It's all well and good for programs to exist in this state, but it's difficult to follow, and almost impossible to debug.

We have to make some changes to our program to make it more readable. One way is to use labels. Labels are just names we can use to identify parts of our program. This is away from good program design.

Machine Language Equivalent of 'GOTO'

In BASIC, you are familiar with the instruction 'GOTO', which jumps to another part of the program. In machine language, you 'GOTO'.

It's just a simple jump instruction. You can see that the first byte is the address of the next instruction and you are half-way there.

The most simple instruction is "Jump To":

```
JP XX XX  
JP (H),  
JP (IX)  
JP (IY)
```

It's just a simple jump instruction. It depends on the status of one of the flags, such as the carry flag. This conditional jump instruction is:

```
JP cc, nn
```

where cc is the condition to be met. If we had

```
JP Z,000C
```

and start again with 'K',

'JUMP', it will jump. Because almost any code can be construed as

read the byte at the address it finds and presumes that is the

start of the next instruction.

The way the CPU works out the jump instructions is really quite
simple. It takes the address of the instruction it is currently executing
and adds the value of the jump instruction to it. This value is usually
relative to the instruction to the program counter.

Thus if it meets a 2-byte instruction, it adds 2, while a 4-byte
instruction will make it add 4 to the program counter.

It is important to remember that you must never change the
contents of the program counter with whatever value you have
specified. That is why you cannot allow any errors to creep in



Long Jumps and Short Jump

jump to anywhere the Z80 chip can possibly go.

The disadvantage of the long jump is that:

- A. Often we don't want to jump that far but still have to use a 3 byte instruction.
- B. We cannot easily relocate the program to another part of memory because we are specifying the absolute address.

can be specified in one byte!

Relative Jump Instruction

IR d
where d is the relative displacement.

example. These conditional jumps are written as
IR cc, d
where cc is the condition to be met.

The value of the displacement 'd' is added to the program counter.

This means it takes the present value of the program counter + d
as the starting point for the jump. This is why the range of the displacement is -128 to +127.

which would be executed if the condition was not met.

This is what happens to the program counter after the relative jump.

Eg In a program such as

Location	Code
32000	ADD A,B
32001	JR Z,02H
32003	LD B,0
32005	Next LD HL,4000H

Load byte at 32000

Because the byte is only a 1-byte instruction so set program counter to 32001.

Execute instruction.

Load byte specified by Program Counter (32001)

Byte is part of 2-byte instruction so add 2 to Program Counter to make it 32003

Get next byte to complete instruction

Execute instruction

specified by Program Counter (32001)

2-byte instruction so add 2 to (now equal to 32005)

Get next byte to complete instruction

Execute Instruction

At location 32001 the program encounters the Relative Jump

the CPU does nothing.

In general, the CPU executes jump instructions as follows.

If the zero flag is set, add 2 more to the Program Counter (this would make it = 32005)

If the zero flag is not set, do nothing
(Program Counter remains = 32001)

In other words, the relative jump allows us to jump over the instruction "LD B,0" in certain cases

for this instruction. It takes less time to do nothing than to calculate the new program counter.

The CPU will therefore execute either the instruction at 32003 or the instruction at 32005 depending on the zero fl

already mentioned

Exercise

JR -2

Machine Language "For . . . Next" Loops

You are, I am sure, familiar with the BASIC form of the "For . . . Next" loops

Loop using the arithmetic functions and the relative jump

```
,Set counter to 1
,Max. of counter + 1
,C = C + 1
,Increment counter
; Is B = A'
JR NZ,LOOP ; If not loop again
```

This will work, but note the following:

We are tying up 2 registers

not set any flags on completion.

A much better way would be if we counted down

We know that we have to do the loop 6 times so why not set 'B' to 6 and count down?

This will give us

```
LD B,6      ;Set counter
LOOP INC C    ;C = C + 1
DEC B      ;Decrease counter
JR NZ,LOOP  ;Loop is not finished
```

The Z80 chip has a special instruction which combines the last two lines above.

This instruction is written

DJNZ d
This results in a b
and therefore saves one byte on the above coding.

Because of the existence of this special instruction, the B register is usually used as a counting register.

```
LD B,10H      ; B=16
BIGLOOP PUSH BC    ;Save value of B
      LD B,0    ;Set B 256
LITLOOP          ;Whatever calculation
                  ;Done 256 times?
DJNZ LITLOOP    ;Get back value of B
                  ;Do bigloop 16 times
```

register after each instruction in the above program.

Waiting Loops:

There are times in machine language programs when things happen fast it is necessary to just wait a little while.

(such as when you are printing thousands of characters a second).

instruction:

```
LD B, Count
WAIT   DJNZ WAIT
```

before proceeding again.

when you write

```
WAIT   JR WAIT
```

Instructions for Call and Return Group

Mnemonic	Bytes	Time Taken	Effect on Flags					
			C	Z	PV	S	N	H
Call address	3	17	-	-	-	-	-	-
Call cc, address	3	10,17	-	-	-	-	-	-
RET	1	10	-	-	-	-	-	-
RET cc	1	5/11	-	-	-	-	-	-

Note: cc is condition to be met for instruction to be executed.
The following are the conditions which can be used:

Flag	Abbreviation	Meaning
Carry	C	Carry Set (=1)
	NC	Carry Clear (=0)
Zero	Z	Zero Set (=1)
	NZ	Zero Clear (=0)
Parity	PE	Parity Even (=1)
	PO	Parity Odd (=0)
Sign	S	Sign Minus (=1)
	P	Sign Pos. (=0)

Flags Effected

No flags are affected by these instructions.

Timing

When the condition is met, the instruction takes 17 ns.
In case of the condition not being met,


```
CP      3  
CALL    2,Three
```

the value stored in the location 'number', but note that it assumes that the subroutines do not change the value in Register A !!!

A shorter routine is possible if you know that there are only the above three possibilities for the value stored in 'number':

```
LD      A, Number  
CP      2  
CALL    2,Two      ; A = 2  
CALL    C,One      ; A ( 2 = ) A = 1  
CALL    Three      ; A ( 2 = ) A = ?
```

Flags and the call instructions do not affect any flags.

Similarly the use of the conditional return from a subroutine is

Instructions for Block Compare and Move Group

Mnemonic	Bytes	Time Taken	Effect on Flags					
			C	Z	P/V	S	N	H
LDI	2	16	-	-	#	-	0	0
LDD	2	16	-	-	#	-	0	0
LDIR	2	21/16	-	-	0	-	0	0
LDAR	2	21/16	-	-	0	-	0	0
CPI	2	16	-	#	#	#	1	#
CPD	2	16	-	#	#	#	1	#
CPIR	2	21/16	-	#	#	#	1	#
CPDR	2	21/16	-	#	#	#	1	#

Flags Notation

Indicates flag is altered by operation

0 Indicates flag is set to 0

1 Indicates flag is set to 1

- Indicates flag is unaffected

Timing

The shorter time indicated is for the case of the instruction terminating - e.g. for CPIR, either BC = 0 or A = (HL).

Block Operations

You should by now be very familiar with the flagging you can do as you carry out the learning & task. Now we will look at how to assign to know what to do next.

This chapter covers some of very useful instructions. It might seem a bit daunting but don't worry, you will see how general things must still think we can make them work with what you already know.

We start with the easiest of these: **PLA** (PICK UP ACCUMULATOR).

The instructions covered in this chapter are by their very nature designed to move data between memory locations which consist of a block of memory rather than just single 8-bit bytes.

Let's start with the simplest of these:
CPI

With the following assembly code:

```
        LDX #1000      ; Load X with address of first byte of block
        LDY #0          ; Load Y with address of first byte of block
        CPI #100        ; Compare Y with the value at X
        BNE loop        ; If not equal, branch to loop
```

It is in fact an extended compare.

It is read in English as "compare and increase". (You will remember that compare and decrease was the name given to the **DCR** instruction.)

With the following assembly code:

```
        LDX #1000      ; Load X with address of first byte of block
        LDY #0          ; Load Y with address of first byte of block
        CPI #100        ; Compare Y with the value at X
        BNE loop        ; If not equal, branch to loop
        INC X           ; Increase X by one
        STA X           ; Store X back to memory
        BNE loop        ; If not equal, branch to loop
```

next location ready for a repeat.

With the following assembly code:

```
        LDX #1000      ; Load X with address of first byte of block
        LDY #0          ; Load Y with address of first byte of block
        CPI #100        ; Compare Y with the value at X
        BNE loop        ; If not equal, branch to loop
        INC X           ; Increase X by one
        STA X           ; Store X back to memory
        BNE loop        ; If not equal, branch to loop
```

Search CPI
LR NZ, Search

With the following assembly code:

```
        LDX #1000      ; Load X with address of first byte of block
        LDY #0          ; Load Y with address of first byte of block
        CPI #100        ; Compare Y with the value at X
        BNE loop        ; If not equal, branch to loop
        INC X           ; Increase X by one
        STA X           ; Store X back to memory
        BNE loop        ; If not equal, branch to loop
```

Unfortunately this is not such a good idea because once it finds a mismatch the program will never end! Fortunately the designer of the processor thought this: the **CP** instruction automatically decreases BC!

We will be able to use the length of the block we wish to

search through and thus specify an end to the search.

is less than 255 bytes long, so that the BC count would only be stored in the C register, we could write

Search	CPI
	JR Z, Found
	INC C
	EC C
	JR NZ, Search
Notfound	.
	.
Found	.

register

instruction

(CD)

which is read in English as "compare and decrease". The decrease refers to HL of course and the effect on BC is still the same.

Even more powerful than these two instructions are the real supermen

These are read as "compare, increase and repeat" and "compare, decrease and repeat".

memory until either a match is found or the end of block is reached. (Naturally we have to specify A, HL and BC before

match found in middle of block or no match at all) we have to ensure we use some code at the end to differentiate between the two possibilities.

language, CPIR and other similar instructions can be very time consuming instruction

by
second.

that the screen is displayed every 1/50th of a second or so you realise that it can be significant.

The remaining block operations are along the lines of "Move it, Mac".

These are

LDI	LDI R
LDD	LDD R

Obviously part of the "load" family these are read as

following set of actions:

- Load (DE) with (HL)
- Increment DE, HL
- Decrement B

Note that this is the only instruction that will load from one register first.

Clever - this way you never forget which register holds the destination address!

one where the information is and the one where the information is going) overlap

Suppose we are using this instruction in a word processing application, and we want to delete a word from a sentence.

The big brown dog jumped over the fox.
7 9 1 3 5 7 9 1 3 5 7 9 1 3 5 7 9

If we want to delete the word 'brown' all we need to do is to move

the rest of the sentence to the left by 6 characters.

DE = destination = character 9
HL = source = character 15
BC = count = 24 characters

Let us start with LD_I after one instruction we have

original = The big brown dog jumped over the fox.

move one char! d (---d

-----> -----> -----> -----> -----> ----->

and HL = 10, DE = 16, BC = 23,

After 2 more instructions

The big dogum dog jumped over the fox.

And After all the instructions have been completed:

The big dog jumped over the fox.e fox.

original sentence and increasing BC to say 30.

overwrite the information we want to shift:

e.g. HL = Source * Character 9
 DE = Destination * Character 15
 BC = Count * 24 Characters

After one instruction we would have

original = The big dog jumped over the fox.e fox.

move char d---) d

new = The big dog judged over the fox.e fox.

After 6 instructions we would have:

The big dog judog juver the fox.e fox.

ood. But another three gives
g dog judog jud og the fox e fox.

time yourself by hand.

It is therefore better to use the 'LD_I' instruction, with the DE
information will not be overwritten in the move.

shift thousands of bytes around very quickly.

Exercise

Write a program that copies memory to the screen.

Note how the 32 first bytes in the screen are arranged.

Now try 256 bytes, then 2048 bytes.

Instructions that are less frequently used

Register Exchanges

We have discussed earlier this chapter the fact that the 6502 language is word processor friendly, and this is done by permitting direct access to memory locations.

You must remember that the 6502 is a single precision processor. This means that it can only work with bytes at a time. It is not able to do arithmetic or counting by themselves.

The first instruction is

EX AF,AF'

B	X	A	C	H	S	N	D	P	M	R	K	S	T
RL	AF	AF'	RL	AF									
B	X	A	C	H	S	N	D	P	M	R	K	S	T
AF	AF	AF	AF	AF									

The next general swap gloves instruction is
EXX

This is a more general swap gloves instruction and follows:

B	C	B'	C'
D	E	D'	E'
H	L	H'	L'

This instruction is very useful, but it is very slow, and makes it limited in use. This is because it acts on all the registers and cannot possibly be byte accurate.

For example, if A and B is modified by EXX

The swap operation proceeds from the short routine along the lines of

PUSH	HL
EXX	
POP	HL

This means that it will save the values of B, DE and HL in the stack, then swap registers by saving H. It is up to you to work with it.

The last instruction in this group does not mean that we have no swap gloves type

EX DE,HL

of DP.

This instruction is indeed very useful, because as we saw HL is a
when the value we want to manipulate is in DE,



Bit Set and Reset

The first two instructions in the set/reset group are used for the manipulation of 8-bit or 16-bit numbers.

Op Code	Op Name	Op Desc	Op Type	Op Size	Op Cycles	Op Flags	Op Side Effects
SET	Set bit	Set bit n to 1	Instruction	8-bit	1	NZP	None
RES	Reset bit	Set bit n to 0	Instruction	8-bit	1	NZP	None

Set and Reset Instructions

For example, if you wanted to set the entire byte at memory location (HL) to 1, you could do this:

$SET \#1, (HL)$ would set the entire 8 bits of that memory location or register.

Op Code	Op Name	Op Desc	Op Type	Op Size	Op Cycles	Op Flags	Op Side Effects
SET	Set bit	Set bit n to 1	Instruction	8-bit	1	NZP	None
RES	Reset bit	Set bit n to 0	Instruction	8-bit	1	NZP	None

Logical Operators

The next two instructions are logical operators. They allow you to turn any bit "on" or "off" at will, or even just look at a specific bit to see what its status is.

Let us look at the first set of instructions.

```
SET n, t  
SET n, (HL)  
SET n, (IX + d)  
SET n, (IY + d)
```

The "SET" instruction turns "on" ($\neq 1$) the bit numbered "n" (using the notation $0 \dots 7$) in register "t" or in the specified memory location.

No changes are made to any of the flags.

It turns the bits "off" (i.e. = 0),

in register "t" or in the indicated memory location.

If Bit 0 then zero flag is set on ($= 1$)
If Bit 1 then zero flag is set off ($= 0$)

~~Take me back to the main menu~~
if the bit is zero, then the zero flag is raised; if the bit is on,
then naturally the zero flag would not be raised.



Rotates and Shifts

You can see where I'm heading now; you can move them to the right, but you can shift those registers any way you like.

The trick is to differentiate between the various shifts and rotations. You must know what kind of shift when a direction is given, because bit 4 of the B register is to be a nibble of the 8-bit A register. The carry is bit number 8 if the bits are numbered 0 - 7).

Now, since we are talking about shifting the carry, it is the 9th bit of the 8-bit register that goes in which direction? It is

left, right or 'rotate' (that's what the meaning of each notation will be made clear later in this chapter):



It is important to note that 'rotate' is not the same as 'shift'. If you rotate a register, you are 'turning it round' (going the long way round). An example of this is the 'RGA' instruction:



Most microprocessors have a rotate as well as shift instruction set. In fact, most of the time, they can be used interchangeably. The difference is that the carry bit is not necessarily preserved. We know RGA just means whatever bit length of transfers are necessary, so it will ignore the carry bit.

Left Rotations

There are basically two types of left rotations:

- * ROTATE LEFT REGISTERS this is a 9-bit cycle rotation as illustrated above for 'RGA'

RGA = "Rotate Left Accumulator"
RL r = "Rotate Left Register r"



D E FF E L R C L A C R M B T R A S T A

R A instruction on a stack of 8 bits

RLCA	- Rotate left circular "A"
RLC r	- Rotate left circular "r"
RLC (HL)	- Rotate left circular (HL)
RLC (IX + d)	- Rotate left circular (IX + d)
RLC (IY + d)	- Rotate left circular (IY + d)



D E FF E L R C L A C R M B T R A S T A

SLA - Shift Left Accumulator



D E FF E L R C L A C R M B T R A S T A

as nothing is transferred to the accumulator. (Think of $A = 0^H$).

RIGHT ROTATIONS

D E FF E L R C L A C R M B T R A S T A

and rotations can be spanned to the right as to the left.

RRA - Rotate Right Accumulator
R R - Rotate Right Register



RRA Rotate Right Circular 'A'
 RRC r Rotate Right Circular 'r'
 RRC (HL) - Rotate Right Circular (HL)
 RRC (IX+d)- Rotate Right Circular (IX+d)
 RRC (IY+d)- Rotate Right Circular (IY+d)



A similar shift right is available as for shift left:

SRL r - Shift Right Logical Register 'r'



In this case this is pure division by 2 as long as we are using
255).

negative numbers by setting bit 7 to 1 (ie. giving us a range of
SRA r - Shift Right Arithmetic 'r'



As you can see this is also a division by 2 but it preserves the sign bit.

In and Out

In or Out - what's it? Simple - In & Out are terms of machine language programming.

There are times when the CPU needs to get information from the outside world. It can be keys and such as control keys or from the cassette player.

As far as the CPU is concerned that's totally irrelevant. It just needs to receive bytes from the outside. The monitor is prepared to receive them in one form or another. The CPU doesn't know and doesn't care to know how a cassette player works.

A byte is an 8-bit number which is the basic unit of memory. It's not a single digit, it's three or four digits. It's not a popular notion to number the bytes of a program. Programs are often described as made by the hardware manufacturer, but that's not true. In fact there is only the keyboard, the program is the cassette player.

By the way, if you want to know about how the data is transmitted, it's ASCII. So, if you're going in or going out, it's an 8-bit byte.

The next question the cassette player has to ask is "which side of the keyboard do I do this?" so that's going to depend on what you do and you use the instruction

IN A,(FE)

Now, when looking your keyboard up the keys of the keyboard are arranged so as to be represented by 8-bit bytes.

The question is now what is a word? Except for the first few letters, all my other keys are letters in the西方的字母表. If you specified which letter is which, then it's reasonable to be examined!

The keyboard is divided into two sections by two blocks of 5 keys:

	2	3	4	5	6	7	8	9	0	.	,
*	Q	W	E	R	T	S	D	F	G	H	I
A	S	D	F	G	Z	C	V	B	N	M	L
FT	Z	X	C	V	B	N	M	+	SPC	{=}	

You can see that there are 8 blocks of letters and we should therefore be able to correlate this with the 8 bits of 'A'.

This is in fact the case.

All of the bits of 'A' are set to 'ON' except for one bit which specifies the block to be read.

For example, if the keyboard has four blocks of information, then the value of 'A' determines which piece of information it gets.

Thus to read the keys in the block "1 2 3 4", it is bit 3 of 'A' which should be off:

$$A = 1111\ 0111 = E7$$

The contents of the keyboard are returned in 'A' with the information coming into the lower bits of 'A'.

- (e.g. Key '1' -> Bit 0 of 'A')
- Key '2' -> Bit 1 of 'A'

If block 4 was chosen instead (i.e. $A = \text{EFH}$) then the information would come in as

- Key '0' -> Bit 0 of 'A'
- Key '9' -> Bit 1 of 'A'

Since block 4 does not contain any keys, A from the keyboard would just be zero, so that both '0' and '1' would both go to bit '0' of register 'A'.

For example, you may want to read all four blocks at once. You can do this by doing it all in one instruction (rather than the two instructions which would be required to read one block at a time).

Information at once
e.g. $A = 1110\ 0111 = E7$

Note that both bits '3' and '4' are 'OFF'.

For example, if both keys '1' and '2' were pressed simultaneously, both would set bit 0 of 'A'.

- (e.g. '1' or '0' -> Bit 0 of 'A')
- '2' or '9' -> Bit 1 of 'A'

As you can see, this is a very useful technique for reading blocks which they belong to different blocks in the keyboard.

Note that if you use the instruction

IN R, (C)

selected

cassette input/output doors.

This is still door FF, as mentioned above. The major problem

instruction path.

The OUT instruction is also used to generate sound on the Spectrum and to set the border colour.

Page 160 of the Spectrum manual discusses the BASIC OUT

bit 4 sends a pulse to the internal loudspeaker.

To set the border colour, load A with the appropriate colour then execute the OUT (FF),A instruction. Note that this

manual).

(BL instruction).

Creating your own sound

16K of RAM

This is done by having the RAM T
line low.

The effect of this is that any program that requires exact or regular timing is impossible as it is not possible to predict the

such interruptions occur if the program and data the Z80 is accessing is in the ROM or in the upper 32K of memory.

To summarise this in layman's terms, you can produce sounds and

BEEP routine ~ see the chapter on the Spectrum's features).

To create sound, you need to send a pulse to turn on the
little while later, you need to send another pulse to turn it off.
Then a little while later, on again,

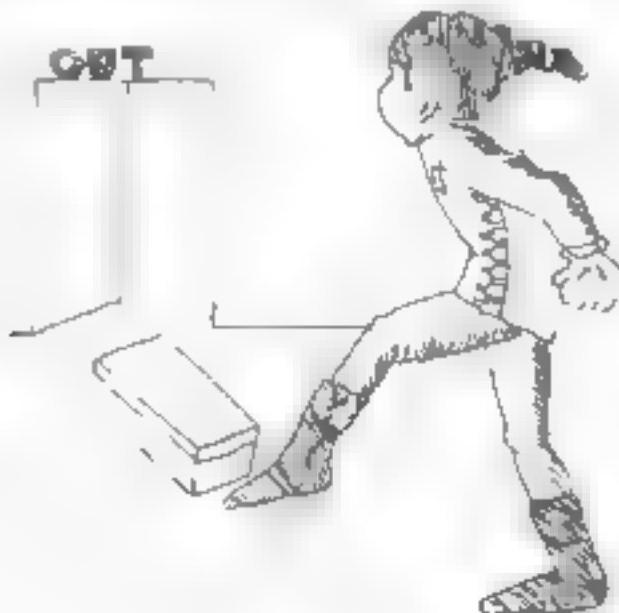
In this way sound is created. The total length of time between

you & minimal degree of control over volume.

Note that you must use a value of A for on and off such that the pattern similar to the LOADING pattern.

Exercise

Frequency.



BCD Representation

B stands for binary-coded decimal. This is a way of representing information in decimal format.

Each digit requires four binary digits from 0000 to 1001. There are 16 possible combinations and this representation.

4 bits are needed to encode a decimal digit. Two bits make up every byte. The first nibble contains the tens.

e.g. 00000000 is BCD representation for decimal 00.
00101000 is BCD representation for decimal 50? "10" is the BCD representation for "58"? "10"?

Is "10100000" a valid BCD representation?

BCD ARITHMETIC

The byte convention in representing numbers can lead to problems in addition and subtraction.

Considering the following:

BCD	08	0000 1000
BCD	03	0000 0011
BCD	11	0000 1011

You notice that the result of the operation is wrong. To correct the mistake, an instruction called "decimal adjust arithmetic", called "DAA", must be used. A half carry is generated if the sum is greater than 9.

The next problem is still started by the same example. The result will be 100000000000000000000000000000000. The rightmost digit is the carry. This extra carry must be cleared (zeroed) and added to the second BCD digit.

The "half carry flag", H is used to detect this carry.

```
LD A, 12H    ;load literal B = 2
ADD A, 24H    ;add literal BCD "24"
DAA           ;dec and adjust result
LD (addr), A  ;store result
```

You will have to use BCD representation in your programming. It is suggested to know that the Z80 chip set supports this representation and the DAA instruction which makes the life of a small group of BCD users simpler.

Interrupts

An interrupt is a signal sent to the microprocessor, which may occur at any time and will generally suspend the execution of the current program (without the program knowing it).

The 64 pre-programmed ones are provided in the Z80. We have already seen the non-maskable interrupt NM and the Soft Interrupt (INT).

In most programming projects, however, we will only look at the usual maskable interrupt (INT).

The DI (disable interrupt) instruction is used to set (mask), while the EI (enable interrupt) instruction is used to set (unmask).

When an interrupt occurs, the current program is suspended and control is passed to the interrupt service routine. This routine is usually a short one-line A RET instruction. It must always return to the point from which the interrupt.

Another important interrupt is the keyboard interrupt, which handles key presses and releases. This interrupt allows the keyboard to be controlled by the ROM BIOS or by your own routine.

The keyboard interrupt is generated when a key is pressed or released. You can write your own interrupt routine to do so.

It is not necessary to write your own interrupt routine, though. The keyboard interrupt is handled by the ROM BIOS.

Restarts

This is another reason why you should not implement interrupt handling. That's why you will be unlikely to use RST instructions in your programs.

The RST performs the same actions as a call, but all the arguments are one of eight addresses in the first 256 memory locations: 0H, 08H, 10H, 18H, 20H, 28H, 30H or 38H.

The disadvantage of the RST instruction is that it frequently needs several bytes of assembly language code. The RST instruction also takes less time than a CALL instruction.

A disadvantage of RST over function calls is that it can only be used to call one of the above eight possible locations.

As a final note, if you are working with the ROM, you should know that it is possible to put your program in the ROM. This is done by marking the ROM as being read-only. If you know what address you want to have the RST instructions

Y	M	B	C	L	H	W	M	R	S	T	U	V	W	X	Y	Z
R		HK	AK	Y	H	HK	RIM	SY	CL	CR	SR	TR	VR	XR	YR	ZR

Programming Your Spectrum

Planning Your Program

Machine language is a very powerful language. It allows you to do anything at all.

Fortran or Cobol or any other language can be done in machine language.

Machine language programs can be run faster on the Z80 than Fortran or Cobol.

This total flexibility can however also be a trap to the unwary programmer. In machine language there are no checks on whether the statement is a legal one.

Or another, the Z80 chip will process everything

totally illegal to any higher level language.

without doing any coding for a long time.

Suppose you wanted to write a lunar lander program

INSTR	Display Instructions
INSTR	Jump back to INSTR till ENTFR pressed
DRAW	Draw Lander
LAND	Move Lander
	If fuel finished go to CRASH
	Jump back to LAND if not ground
	Jump back to INSTR for next GO
	Jump back to INSTR for next GO

Notice how this 'program' is written totally in English. At this stage, no decision has been made whether the program is to be that decision - the concept of the Lunar Lander program is not dependent on the coding.

Now comes the part of logic testing,

you wish to see included in the program are covered

things be put into subroutine

way to finish the program

arcade machine, but in your program you may decide you would like to be able to turn the program off!

We now change the last part of the program as follows.

GROUND Print Congratulations
 Jump to FINISH

FINISH Ask player if finished
 If not, jump to INSTR
 If yes, STOP

Note that we have used labels to describe certain lines in program. These are very valuable, the more so if you choose labels which are descriptive to their meaning.

thing to one of the lines or modules above.
This is why this approach is called the top down approach.

For example we can expand the 'finish' module above:

FINISH Clear screen
 Print "Would you like to stop now?"
 Scan keyboard for input
 If input = yes then stop
 Jump to INSTR

run a particular module on its own, so that it is ready for the final program.

Let us go down one level further again, and look at the Clear screen

see in more detail.

If you were writing in BASIC, all you would have to say is
PRINT "SOS"

but in machine lang e that simple sentence, 'Clear screen' can be
decept(ve.)

We might therefore do something like

Fill next 6144 positions with blanks

a clear screen routine is meant to do and what it will actually do.

and 40-50

routine above will clearly be inadequate.

- to work on the attribute file well. (Note how much more complex certain tasks can be in machine language than in BASIC.)

where used to expand the acronym to read

Ficed Access Points

REST 8144 bytes with R=55%

feature file `background`

next 768 bytes with no F/Fok desired

The next level down is the one where you must finally do the coding, so let us look at filling the screen with blanks.

CLEAR	LD HL,SCREEN	;Screen start
	LD BC,6144	;Bytes to clear
	LD D,0	;D=blank
LOOP	LD (HL),D	;Fill blank
	INC HL	;Next position
	DEC BC	;Reduce count
	LD A,B	
	OR C	;Test if BC = 0
	JR NZ,LOOP	;Again if not end

this way build up very complex programs indeed.
By the way, you no doubt understand now why machine language
language programs'

x

There are more ways than one to write any particular routines.
Let us look at the simple clear screen routine written above.

This could be handled by several different approaches.

Exercise 1

only so that we may make use of the 'DJNZ' instruction?

Exercise 2

blanked using the more powerful 'LDIR' instruction?

Think carefully of what 'LDIR' does: it is not always necessary to have 6144 blank positions elsewhere.

Answers:

More than one possible answer can be "right" - the only test is "does it work?" In other words does it do what YOU want?

Using DJNZ

CLEAR	LD HL,SCREEN	
	LD A,0	
BIGLOOP	LD B,24	,Set B=24
	PUSH BC	Save value
LITLOOP	LD B,A	,Set B=256
	LD (HL),A	;
	INC HL	,Fill in 256 blanks
	DJNZ LITLOOP	
	POP BC	(Get back value of B)
	DJNZ BIGLOOP	(Do it until end)

We have been able to use 24 times 256 (=6,44) to clear the screen

Points of note are

We can set B = 0 to go through the DJNZ loop

256 times. (Why?)

This procedure would not normally be used in a program unless we were also using register C for other purposes.

Using LDHR

CLEAR	LD HL,SCREEN	;Source
	P SH HL	
	POP DE	
	NC DE	
	LD BC,6144	
	LD HL,0	
	LD R	;Move it

Note that we have found $DE = HL + I$ by getting $DE = HL$ and increasing DE . This can be achieved more easily by loading the value of SCREEN + I into DE directly but this requires 1 more byte!

The reason this LDHR works is because we are using the fact that the data is overwriting the block to be written as we proceed. This Move Chapter.

If you add up the memory required, the first method requires 14 bytes, the second 16 bytes, and the last 13 bytes.



Features of the Spectrum

Some hardware features of the Spectrum are specially designed to match machine language programs.

Input - keyboard

As far as input to the Spectrum is concerned, we will ignore cassette input and concentrate on the keyboard.

The keyboard is the only input which provides real-time

eight rows and five columns as in appendix A.

In their normal state (when they are not pressed) the keys in a high mood i.e. the interrupt level is zero.

When a particular key is pressed, the interrupt level corresponding to that key will be raised to one.

which can be used in machine language programming.

(p 160) of the Spectrum manual.

For example, the interrupt levels are listed in the leftmost column of the table in appendix A.

e.g. For the "H - ENTER" half-row we load A with value BFH

LD A, BFH

the INPUT instruction is issued.
e.g. The port used is the FFH port

IN A, (FEH)

Since there are 8 keys per half row, we are only interested in the five low order bits of the returned byte in A.

If no key is pressed in that half row, the value of the 5 lower five bits will be $(2^{4*4} + 2^{4*3} + 2^{4*2} + 2^{4*1} + 2^{4*0}) = 16 + 8 + 4 + 2 + 1 = 31$.

register A = xxxx1111 when no key is pressed.

If we want to test whether the right most bit is pressed we have to see whether that bit is low.

There are two ways to test this:

- i. Use Bit Test Instructions, eg. BIT 0, A
If the bit is low (not set) then the Zero flag will be set.
- ii. Use Logical AND instructions AND 1
If the bit is low (not set) then the result will be zero and the Zero flag will be set.

The first method needs to be done separately for each row. To test a specific column you can use the BIT command to test one bit at a time or you can use the AND instruction to test all keys in that half row which would be more efficient. Both test instructions are followed directly by relative jumps.

e.g. To test bit 0 and bit 1 using the first method

```
BIT 0, A      ;test bit 0 of A set or not
JR Z, NPRESS  ;jump if not pressed
BIT 1, A      ;test bit 1 of A set or not
JR Z, NPRESS  ;jump if not pressed
```

do whatever if both are pressed

Method 2

The second method of testing using logical AND requires a little more code. To test bit 0 we use "AND 1"; to test bit 1 we use "AND 2"; to test bit 2 we use "AND 4" and so on.

To test two keys we use "AND x" where x is the sum of the values we will use when testing each one key individually.

e.g. To test both bit 0 and bit 1 of A are set

```
AND 3      ;test both bit 0 and bit 1
            ;is set
CP      3      ;test if both set
```

```
JR      NZ,NBOTH    ;jump if not both pressed
```

```
To test if either bit 0 or bit 1 of A are set  
AND      3          ;test either bit 0 and bit  
                  .1 is set  
JR      Z,NOTONE    ;jmp if not one is pressed
```

Ex:

for your Spectrum.

You will need to

- a. check the row address that needs to be loaded into A.
- b. send it to the input port FEH.
- c. test the bit that is set by the (ENTER) key.

Output Video Screen display

The Video screen display is the main source of output for the computer to communicate to the user.

screen memory of the Spectrum is organised

210040	LD HL,4000H	,load HL with start of +d splay file
16FF	LD (HL), FFH	fill that screen location
110140	LD DE,4001H	,load DE with next byte
		+in display
010100	LD BC,I	,BC contains number of +bytes to be transferred
FEBO	LDIR	,move a block length BC +from (Hl) to (DF)
C9	RET	end of progr

transferred from (Hl) to (DF).

Now change the fourth line to read LD BC, 31 (011FOU). You may be

Note how a very thin line has been drawn across the top of the screen. The first J2 bytes of the screen memory relate to the first byte of each of the first 32 characters.

surprised. The next byte after the 32nd one is not on the second row of dots on the screen. It is the first byte of the 32nd character! And so on up to the 256th character.

that line to LD BC, 2047 (0 FF07) and run the program. You will find that the top third of the screen only has been filled.

You can experiment with this, using different values for BC up to LD BC, 6143 (0.FF17). In this way you can watch the way Spectrum organ sees the screen.

The screen memory is actually divided into three lots.

1. Memory 4000H - 47FFH (=*) first eight lines.
2. Memory 4800H - 4FFFH (=***) second eight lines.
3. Memory 5000H - 57FFH (==) third eight lines.

Not only that, but you will retain that each character of the Spectrum is composed of eight 8-bit bytes which makes up 64 dot

Eg. For the character "A", its character representation is

0	00000000	0H
16	00000000	00H
16	00000000	00H
16	00000000	00H
6	00000000	0A+
0	00000000	0H
6	00000000	00H
0	00000000	0H

the first 256 bytes from 4000H to 40FFH correspond to the first byte of each of the 256 8-byte character of the first eight lines.

Then the next 256 bytes from memory location 4100H to 41FFH correspond to the second byte of each of the 256 8-byte character of the first eight lines and so on

thus

First character of the screen is

1st byte	4000H
2nd byte	4100H
3rd byte	4200H
4th byte	4300H
5th byte	4400H
6th byte	4500H
7th byte	4600H
8th byte	4700H

Isn't it? But we have to accept the Spectrum the way it is.
It.

Can you write down the eight bytes that correspond to the 31st character of the third line of the screen? You can refer to Appendix B, the screen memory map.

4D0H, 415EH, 425EH, ..., 475EH.

To follow on the concept we have developed about the screen display, the memory locations that correspond to the first character of the second eight lines lot is
4800H, 4900H, 4A00H, 4B00H, 4C00H, 4D00H, 4E00H, 4F00H.

Similarly, the first character of the third eight lines lot has its eight bytes in memory locations

5000H, 5100H, 5200H, 5300H, 5400H, 5500H, 5600H, 5700H.

There are some advantages, however, in using machine language. The apparent complexities are worth overcoming. As a trivial example, in BASIC, if you try to PRINT into the input section of the screen (the bottom two lines), the BASIC system will object most violently. But in machine language you have full access to the

whole screen.

For address the first 4 bytes of memory are more or less you will see at the right part by **First Byte** & **RF8x** at each character of address which locates the three memory positions the character is in.

For example, if $40H \leq HOBFB \leq 41H$ char is in first eight lines lot
if $48H \leq HOBFB \leq 49H$ char is in second eight lines lot
if $50H \leq HOBFB \leq 51H$ char is in third eight lines lot

Note: In the above diagram the bits of the 4-Bit High Order Address are shown which are the first four bytes of the character it belongs.

To take some example consider the address $4A36H$ & we want to know which portion of the memory contains the displayed character (if any??).

Try the following example,

Suppose we are given an address $4A36H$. The High Order Byte of the address is $4AH$ so,

i. we know that it is within the screen display memory since its value is in between $40H$ and $58H$.

ii. Its binary representation is $010010,0$

iii. from the lower three bits we know that it belongs to the third byte of a character position on the screen.

iv. If we made the lower three bits zero, then the value would be $40H$. Thus we know this belongs to the second eight lines lot, i.e. the middle portion of the screen display.

The conclusion we can reach is that the byte given refers to the third byte of a character in the middle portion of the display memory.

What character of the middle portion is the byte belongs to? To answer this question, we'll need to know the value of the Low Order Byte of the address.

We know that $3B$ & the address is $36H$ & the address refers to character at $48H$ i.e. the 54 th position away from the first character of the middle portion.

Considering this 5 characters, the position of address is

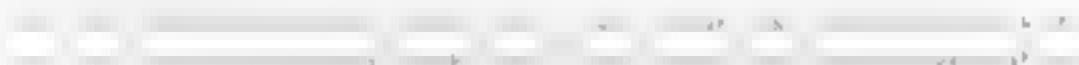
the second line of the middle screen display portion and is the 154-128th character of that line.

of the 23rd character of the 10-th line from the start of the

Exercise

Which byte of which character does the address 364FH refer to?

x



Input Video display attribute

The display attribute memory is easier to understand than the display memory because it has a one-to-one relationship with the screen display characters.

The attribute file is located in memory from 5800H to 5AFFFH. It is 768 bytes, which correspond to 24 lines of 32 character each. In other words, there is one attribute byte for each character position.

Thus, 5800H corresponds to the attribute of the first character of the first line, 5801H the second character, 5802H the third,...,581FH the thirty second character of the first line.

Similarly, 5820H holds the attribute of the first character of the second line, 5840H of the third line, ... and 5A20H the attribute of the first character of the last line of the screen.

We know that for each character position on the screen, there is a corresponding attribute byte in the attribute memory, made up as follows:

attribute byte b b bbb bbb

bit 0 - 2 represents the ink colour of the character 0 to 7.

bit 3 - 5 represents the paper colour of the character 0 to 7.
t = 1, normal if 0.
b = 1 if 0

What is the address of the attribute byte that contains the first byte of the middle screen section? What is the address of the first byte of the third section? Answers are given on the next page, but try to work it out for yourself.

Exercise

Can you write a subroutine that converts a given address on the screen to its corresponding attribute address
e.g. 4529H

You must in effect determine which character of this screen this belongs to, and then add this to 5800H

The following program shows a short method of achieving this:

LH AL, 4324H
LD A

oad the given address to HL
oad the high order byte to A
Drop bits 3 and 4 to
determine which portion of the
screen the address below
left right accumulator
three times - ie divide by 8
result can either be 0,1 or 2
depending whether H was
48H, 50H or 52H
transform to attribute memory
+ contain attribute address
58H, 59H or 60H
using the same

You may be to think about this for a while!

The way the program works is related to the answer of the first exercise

1st char. of 1st screen section = 4000H	A	58H
1st char. of 2nd screen section = 4800H	A	59H
1st char. of 3rd screen section = 5000H	A	60H

2nd char. of 1st screen section = 4001H	Attribute add	59H
etc.		

This should make things a little clearer!

Output - Sound

offers is sound. It would be a waste if we didn't make full use of this facility.

generating sound.

- I. Sending signals to the cassette output port 254 for certain duration of time using the OUT instruction 254.
eg OUT (254), A
- ii. Set HL, DE to certain values and call the ROM sound routine used to generate sound.

The input parameters are

DE = duration in sec * frequency
HL = (437,500 / frequency) - 30,125

Then

CALL 0185H.

The first way of sound generation has the advantage of being free from any ROM call. It is shorter in terms of time to execute. But there is a limitation.

For example, if you want to generate a sound for 1 second, you have to

call the ROM routine 437,500 times. This is not possible on a 48K machine.

So, you can't generate a sound for more than 1 second.

Using this method, but it will not be 'clean sound'. You have to use the second method of sound generation (of calling the ROM routine) to get that result.

For example, if you want to generate a sound for 1 second, you have to

call the ROM routine 437,500 times.

program. You can think of the DE register pair as holding a value for the duration of the sound, and HL a value for the frequency.

sound you want.

The limitation of this method of course is that you are restricted to whatever sounds you can create with the BEPP command.

Monitor Programs

EZ-Code Machine Language Editor

This is a machine code monitor program that allows you to

- I. INPUT your machine language program module in either a fully assembled format
OR a semi-assembled format with all relative jump and absolute jumps expressed in the form of line number.
- II. LIST the source input program module.
- III. DUMP the input program module into the specified memory address.
- IV. EXAMINE a range of memory locations.
- V. SAVE EITHER the "source module"
OR the dumped program in fully machine code format
- VI. LOAD a saved "source program" from the cassette.
- VII. RUN the dumped machine program module

not need to calculate relative or absolute jumps
200 Instructions
to wipe off the EZ-code program.)
CONCEPT behind the EZ-code

of a BASIC program.

code) has a line number and up to 4 bytes of machine code.

A major benefit is therefore the ability to "edit" any line. The "source program" can also be SAVED separately to tape, allowing work in progress to be saved.

A major innovation in this program is the ability to insert relative jumps or absolute jumps without having to calculate the number you wish to jump to!

the scope of a relative jump.

to memory.

EZ-code Instruction Summary

Note that the first question the program will ask you is
"Loading address".

This cannot be below 31500.

*** Entering LINES ***

1. To ENTER lines of "source program"

(line no)(blank)(maximum of 4 bytes in Hexadecimal)
ENTER;

e.g. 1 210040 will insert the machine code instruction
D HL,4000H into line number 1.

2. To EDIT a line

(line no)(blank)(retype new bytes)(ENTER)

e.g. 1 210140 will change line number 1 to the instruction
10 HL,4000H.

3. To DE-ERASE an instruction line

line-no)(ENTER)

e.g. 1 (ENTER) will delete line number 1.

iv. To specify RELATIVE or ABSOLUTE jump

(line-no)(blank)(jump instruction) ("lower case "L)
(line no ENTER)

e.g. L c312 represents the instruction JP to line 2.
2 18.1 represents the instruction JR to line

other COMMANDS ***

v. dump ENTER

* dump the source listing into the memory starting
from the specified LOADING addre
* this must be done before running the machine code
program.

abbreviations: du

vi. exit ENTER

* exit from the EZ-code and re-enter BASIC system.

abbreviations: ex

vii. list ENTER

* list the first twenty-two instruction lines of the
source listing
* press any key except 'm' and 'BREAK' to continue
listing

abbreviations: l,

list # ENTER

* list twenty-two lines of the source listing
starting from line number #, a number between 1 and
210 inclusively.

abbreviations: NO ABBRV

viii. load ENTER

* load a source listing module from the cassette
replacing the existing module.

abbreviations: lo

v. mem ENTER

prompt: Starting address:

* enter memory address you want to start displayin
from.
* can be from 0 to 32767 for 16K Spectrum or 0 to
65535 for 48K Spectrum
* press "m" to exit memory examine mode.

abbreviations :sc

sc - newENTER

- * clear the current module and re-run the EZ-code.
- * this is useful when you want to start coding in another program module

abbreviation :re

vtl:run(ENTER)

- * run the dumped program module from LOADING address you specified when you start running the EZ-code program or when you LOAD a new source listing

abbreviation :ru

- * save either the source listing or dumped machine code onto cassette.

prompt: Enter name

enter the name you want to do

Source or Machine code (s or m)

enter s for source listing saving

enter m for machine code saving

Start tape, then press any key,

make sure that the cassette is properly

press any key when the cassette is ready

abbreviation :sa

NOTE

1. If you don't want the result of BC register returned after running, change line 3090 to :

"

restart the EZ-code:
either use R'N and resulting with all variables
reinitialised.
Or use GOTO 2020 which returns the prompt
"Command or Line(##) "

2. All numeric entry except machine instruction
code has to be in decimal format.

4. To enable you to insert additional lines in the current
listing, it is good to space out the listing.
i.e. instead of entering instruction lines as 1, 2, 3
enter as 1, 3, 10 etc
This will makes the input of the module more flexible.

TYPE SE or EZ-code

Enter the following code:

210040	LD HL,4000H	full screen
110140	LD BC,4001H	
01FF17	LD BC,6143	
FFF	LD A, 0FFH	
77	LD (HL), A	
FEBO	LD1R	
F7F	LOOP: LD A, 7EH	trap BREAK key
07E	IN A, (FFH)	
F6C1	AND I	
20F8	JR NZ, LOOP	
9	RST	

To enter the above code using EZ-code

CMDN

address : 0000H
Line 1 : 210040 ENTER
at Line 000 : S 110140 ENT R
Line 001 : D 01FF17 ENT R
 S 3e7f ENTER
 D 077f ENTER
Line 002 : 25 ed><ENTER
at Line 003 : D 0e7f ENTER
Line 004 : 15 dffe ENT R
Line 005 : 40 eb 1 ENT R
Command or Line 006 : 45 2e< ENTER

(This is 24 then lower case 'a', then 10, in other words:
F NZ, Line 10 >

Command or Line 007 : 0 c9 ENTER
Command or : 1st ENTER
Command or Line 008 : dimp ENTER
Command or Line 009 : 0000 ENTER
Starting address : 31500 ENTER

this is the key to exit the memory

Command or : 1st <

Note how there must be a space after the line numbers.

EZCODE

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```

100 REM Machine
110 REM Machine_code_monitor
120 TO TO 900
130 DEF FN d($$) = ($$ > "9")*( CODE ss 55)
+($$ <= "9")*( CODE ss-48) +($$ > "A")*72
140 DEF FN o(D$) = ((0$ = "ca")+(0$ = "da"))
+(0$ = "ea")+(0$ = "fa")+(0$ = "e2")
+(0$ = "d2")+(0$ = "e2")+(0$ = "f2")
+(0$ = "c3")+(0$ = "2B")+(0$ = "30")
+(0$ = "2D")+(0$ = "20")+(0$ = "1B")
+(0$ = "10")
150 REM
160 REM Line PRINTING routine IBA
170 LET F = ze + PRINT AT z - ze
180 FOR J = p11 TO p12
190 IF C$(J, on) = " " THEN GO TO 1110
200 PRINT TAB tr - LEN STR$ J1 J1 TAB fr$ "."
210 IF C$(J, tw, on TD on) = "1"
    THEN PRINT C$(J, on)+ " "+C$(J, tw)+C$(J, tr)
    + GO TO 1490
220 PRINT C$(J, on)+ " "+C$(J, tw)+ " "
    + C$(J, tr)+ " "+C$(J, fr)
230 LET F = Fton
240 IF F = 22 THEN TO TO 112
250 NEXT J
260 PRINT AT z - 251 "-----"
270 LET FN
280 REM
290 REM main_routine IBA
300 INPUT "Command or Line(LEN) $ " A$
310 IF ASL TO fr$ = " " THEN GO TO m
320 IF ASL on > "9" THEN GO TO 2000
330 LET K$ = "" + FOR K = on TO fr
340 IF ASL TO K$ = " " THEN GO TO 2000
350 LET K$ = K$+ASL(K TO K)
360 NEXT K
370 IF K = 5 OR VAL K$ = ze OR VAL K$ > 1n
    THEN GO TO m
380 LET J = VAL K$ + LET n = J
    + REM Line number must be 3 bytes
390 LET AS$ = AS$+on TO 1
400 LET K$ = ""
410 FOR K = on TO LEN AS
420 IF AS$K TO K$ <> " "
    THEN LET K$ = K$+AS$K TO K
430 NEXT K
440 LET AS$ = 1$
450 IF AS$on) =      THEN GO TO m

```

```

*CLS = FOR I = on TO 7 STEP tw
LET K = INT (I/16+on)
LET CS(i, K) = AS(I TO I+on)
NEXT I
IF CS(i, on) = "++" THEN GO TO 2250
IF n < TP THEN LET TP = n
IF n > BP THEN LET BP = n
GO TO 272
*IF n < BP THEN GO TO 278
*IF RF = on OR CS(BP, on) <> "++"
THEN GO TO 2720
LET BP = BP+on + GO TO 226
IF n < TP THEN GO TO 2720
IF CS(TP, on) <> "++" THEN GO TO 232
IF TP < BP AND TP < 1n THEN LET TP = TP+on
GO TO 224
LET TP = on
LET po = n
IF n < TP THEN LET pp = TP + 80 TO 2380
LET numpp = 20
IF po = TP OR numpp = 11 THEN GO TO 2380
IF CS(po, on) <> "++"
THEN LET numpp = numpp+on
LET pp = po+on + GO TO 2350
LET p11 = pp + LET p12 = BP
GO SUB 100
REM print a block of lines
GO TO 28
REM
REM ***** Commands***** IBW
LET VS = ABS TO tw
VS = "d0;" THEN GO TO 5000
VS = "e0;" THEN STOP
VS = "f1;" THEN GO TO 4000
VS = "f0;" THEN GO TO 7000
VS = " " THEN GO TO 6000
VS = " " THEN R IN
VS = " " THEN PRINT USR R
THEN GO TO 8000
**  

*REM ***** routine***** IBW
LET p13 = TP + LET p12 = RF
LET n1 = CDR AS(6 TO 6)
IF LEN AS > 16 AND n1 > 47 AND n1 < 53
THEN LET p13 = VAL AS(5 TO 8)
GO SUB 1000
GO TO 28
REM
REM ***** DUMP routine***** IBW
CLS = PRINT AT 28, 253 INK on; INVERSE on
+ FLASH on; "DUMPING" + LET G = R

```

```

PRINT AT on, ze
FOR J = TP TO BF
  IF (S(J, on) = "+") THEN GO TO 5470
  IF C$(J, tw, on TO on) <> "3" THEN GO TO 5780
  POKE B, ze + POKE G+on, ze + POKE G+tw, ze
  * POKE G+tr, ze
  LET zt = VAL (C$ J, tw, tw TO tw)+C$(J, tr)
  PRINT TAB tr- LEN STR$ J1 INVERSE on+ J
    * TAB frt INVERSE zet "+"
    * C$ J, on)+"+"+(%,J, tw)+C$(J, tr)
    * "+>"+
  IF j1 < ze OR j1 > ln THEN GO TO 5460
  LET CJ = FN O(C$ J, on))
  PRINT TAB 17- LEN STR$ j1t INVERSE on+ ,1
    * TAB 10t INVERSE zet "+1 C$(1), on)
    * "+1 C$(j), tw)+ "+1 C$(1), tr)+ "+"
    * C$(j), frt
  IF ABS CJ <> on THEN GO TO 5460
  LET dd = j1 > j1-1+j1 < j1
  LET ja = G + LET dd > ze
  IF j1 = J THEN GO TO 5270
  LET cl = J+dd
  LET nl = ze + IF (%,cl, on) = "+"
    THEN GO TO 5240
  IF C$(cl, tw, on TO on) <> "1"
    THEN LET nl = on+(C$(cl, tw) <> "+")
      + (C$(cl, tr) <> "+")
      + (C$(cl, fr) <> "+")
  * GO TO 5 2
  LET TJ = FN o(C$cl, on))
  LET nl = (TJ + on)+tr+(TJ - on) $tw
  IF cl = j1 AND dd > ze THEN GO TO 527
  LET dn = dp+rt
  IF cl = j1 THEN GO TO 5 70
  LET cl = cl+dd
  GO TO 5 80
  IF CJ = on THEN LET ja = ja+dd$dp+(dd > ze) $tr
  * GO TO 5 10
  IF dd > ze THEN LET dp = dp+2
  IF dp > 126 AND dd < ze THEN GO TO 5460
  IF dp > 129 AND dd > ze THEN GO TO 5460
  LET V = 168 FN d(C$(J, on, on TO on))
  * FN d(C$(J, on, tw TO tw))
  POKE B, V + LET B = G+on
  IF CJ = on THEN POKE B, ja- INT (ja/qk)+qk
  * LET B = G+on + POKE B, INT (ja/qk)
  * LET B = G+on + GO TO 5360
  IF dd < ze THEN LET dp = -dp
  LET dp = dp-tw = POKE B, dp + LET B = G+on
  PRINT "ok"
  GO TO 5470
  FOR I = on TO 7 STEP tw

```

```

519 LET K = INT (I/twton)
520 LET V = 160 FN d(CS(J, K, on TO on))
      + FN d(CS(J, K, tw TO tw))
521 IF V < ze THEN GO TO 544
522 POKE G, V
523 LET G = G+on
524 NEXT I
525 GO TO 547
526 PRINT "33"
527 NEXT J
528 PRINT AT ze, 25; "*****"
      + GO TO sr
529 REM
530 REM IBM Memory display XXXXXXXXXX IBM
531 INPUT "Starting address ?"; dm
532 CLS + PRINT AT 20, ze
533 LET G = dm + LET F = ze
534 LET F = F+on
      + PRINT TAB 5+ LEN STRG G1 G + TAB 7
535 FOR I = on TO fr
536 LET V = PEEK G
537 LET H = INT (V/16)
538 LET L = V-16*H
539 PRINT DS(H)+DS(L); DS(L+on) + "A"
540 LET G = G+on
541 NEXT I
542 PRINT "A"
543 IF F <> 22 THEN GO TO 616
544 LET $ = INKEY$ + IF $ = "" THEN GO TO 615
545 IF $ = "A" AND $ = "M" THEN LEFT $ = " "
      + POKE 23692, q+on + GO TO 615
      + POKE 23692, on + PAUSE 20 + GO TO sr
      + M
546 REM IBM LOAD XXXXXXXXXXXXXX IBM
547 S
548 INPUT
      Load program address any key when ready...
      $ + $
549 PRINT AT ze, 25; INVERSE on + FLASH on + "LOADING"
550 LOAD "source" DATA C$,.
551 FOR I = on TO ln
552 LET TP = I
553 IF CS(I, on) <> " " THEN GO TO 7100
554 NEXT I
555 FOR I = ln TO on STEP -1
556 LET RP = I
557 IF CS(I, on) <> " " THEN GO TO 7140
558 NEXT I
559 PRINT AT ze, 25; "*****"
560 GO TO 9150
561 REM
562 REM IBM SAVE XXXXXXXXXXXXXX IBM

```

```

8020 INPUT "Enter name <_>" n$  

8030 IF n$ = "" THEN GO TO 8020  

8040 INPUT  

    "Source or Machine code <_&ts.or.am>"  

    t$ k$  

8050 IF k$ > "S" AND k$ < "M" THEN GO TO 8140  

8160 IF k$ = "S" THEN SAVE n$ DATA (t$) : GO TO mr  

8070 INPUT "Starting address <_>" ss  

8080 INPUT "Finishing address <_>" sf  

8090 LET sb = sf-ss+1  

8100 SAVE n$ CODE ss, sb  

8110 GO TO mr  

9000 REM  

9010 REM initialisation  

9020 LET ze = PI - PI * LET on = PI / PI  

    * LET tw = on+on * LET tr = on+tw  

    * LET fr = tw+tw * LET qk = 256  

    * LET mr = 2020 * LET ln = 200  

9125 BORDER 7 + PAPER 7 + INK on + INVERSE ze  

    + OVER ze + FLASH ze + BRIGHT ze  

    + BEEP .25, 24 + BEEP .25, 12  

9030 DIM A$(13) + DIM D$(tw)  

9040 LET TP = ln + LET BP = on  

    + REM line,number,buffer  

9050 DIM C$(ln, fr, tw) + REM holds,code  

9160 PRINT AT ze, 201 INVERSE on; FLASH on  

    + "INITIALISING"  

9070 FOR I = on TO ln  

9080 FOR J = on TO fr  

9090 LET C$(I, J) = ".  

9100 NEXT J  

9110 BEEP .01, 20  

9120 NEXT I  

9130 PRINT AT ze, 201 "*****"  

9140 LET D$ = "0123456789ABCDEF"  

9150 CLS + PRINT "Lowest address <_>" T15H  

9160 INPUT "Loading address <_>, R : PAUSE 1  

9170 IF R < 31500 THEN GO TO 9160  

9180 CLS + GO TO mr

```

Hexload Machine Code Monitor

This section will explain how to program BASIC programs into memory and save them to cassette tape. Hexload is a cassette monitor that can read and write programs from memory or cassette and also has a built-in cassette to memory.

It is designed to run on the Spectrum 16K and 48K. To use it, you must have a cassette tape and less than 200 instructions.

When we use Hexload, we use basic operations to save each module as machine code on cassette. These modules have been saved in RAM, program or load them into memory again by moving them about their appropriate memory locations.

We will actually apply this technique as we develop the FREEWAY PROG program.

Concept behind Hexload

The concept behind Hexload is extremely simple.

It reads programs from memory RAM or BASIC system 26999.

It then writes them to cassette tape. It is designed for 48K Spectrum.

It offers basic monitoring functions like:

- WRITE onto memory in Hex format
- SAVE from memory to cassette
- LOAD from cassette to memory
- LIST memory contents from a starting address
- MOVE memory contents from one locations to another.

Hexload Instructions Summary

1. WRITE

Write code in HEX format onto the memory.

Format:

Input start of memory where you want to write to in decimal format in response to the prompt.

The address is limited to 27000 - 32758 for 16K
37000 - 65536 for 48K

e.g. Write to address 27000(ENTER)

- b. Enter codes in hex format.
- c. Press "m" to return to main menu.

2. SAVE

Save memory to cassette.

Procedure.

- a. Input memory from which saving starts. Can be any address 0 - 32767 for 16K
0 - 65535 for 48K
- Input number of bytes to be saved
- Input name of the module to be saved
e.g. ROM, etc is T
- Option of verifying the module saved on to the cassette
- It is good to verify so as to ensure that there is no corruption of the module during the saving procedure.

Load machine code module from cassette.

Input memory address to which the module is start loading. The address is limited to same range as in write command

- b. Enter the name used when the module is saved if you are not sure of the name, just press ENTER.,

4. LIST

Display memory contents starting from an address

Procedure

- a. Input address start listing from.
Can be any address as in SAVE command above
Type any key to continue the display.
- c. Type "m" to return to main menu

Move memory contents from start address to finish address into new memory address.

Procedure

- Input move from memory, any address as in the range of SAVE command.
- Input move until memory, any address as in the range of SAVE command.
- c. Input move to memory, address range as in WRITE command
You can even copy the ROM into RAM by using this command

e.g. Move from memory: 0 ENTER
Move until memory: 1000 ENTER
Move to memory: 32000 ENTER

this will move ROM 0 to 1000 to RAM address 32000.

NOTE: Any of the input in above commands which breaches the address range will result in the input being prompted.

<+RUSP>

Try using this monitor to input the module we have developed with F2-coded.

THE YOUNG CAR

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```

100 REM
110 REM monitor program
120 CLEAR 26999 : LET ze = PI - PI
  : LET on = PI / PI * LET tw = ONTOP
  : LET qk = 256 * LET lm = 27000
  : LET ar = 140 * LET wl = 340
130 GO S R 2000
140 CLS
  : PRINT "Start_of_machine_code_area = "
    : lm
150 PRINT "menu" : PRINT
  : PRINT
    : "Write_machine_code.."
160 PRINT
  : PRINT
    : "Save_machine_code"
170 PRINT
  : PRINT
    : "Load_machine_code....."
180 PRINT
  : PRINT
    : "Last_machine_code.."
190 PRINT
  : PRINT
    : "Move_machine_code...  <-->
      : "
    : "Appropriate_hex."
  : " "
  : " "
220 IF q$ = "m" OR q$ = "M" THEN STOP
230 IF q$ = "" OR q$ < "1" OR q$ > "F"
  THEN GO TO 210
240 CLS
  : PRINT "Start_of_machine_code_area = "
    : lm
250 GO TO 2000 VAL q$
260 REM Line Write to Address Line
270 INPUT "Write to address #": d
280 IF d > mm OR d < lm THEN GT TO 210
290 PRINT : PRINT "Write_Address #:": d
  : PRINT "To return to menu enter #"
300 LET as = ""
  : " " : " " THEN INPUT Enter_hex_code :
    : as
360 IF as onl = "m" OR as onl = "M"
  THEN GO TO ar
770 IF LEN as tw > INT 1 LEN as tw
  THEN PRINT "Incorrect entry."
  : GO TO ml

```


The Freeway Frog Program

Program Design

The frog has to cross the highway from one side to the other.

police cars frequently patrolling the highway.

other traffic.

This can be logically subdivided into several parts by the programmer.

This is merely the problem definition stage.

the design and development of the whole project.

FREWAY FROG program structure

Now we have to divide the program into well-defined logical modules.

They are as follows.

1. INITIALISATION

perform all initial tasks.

2. TRAFFIC FLOW

control of traffic on the highway.

This can again be logically subdivided into

- i. regular traffic flow eg. trucks, cars and motorcycles.
- ii. Irregular flow traffic eg. police car.

3. FROG

control the movement of the FROG, crash testing as well as home testing.

4. GENERAL PROGRAM CONTROL

this part of the program takes care of the score calculation and display, testing for termination of the game.

5. TERM NATION

perform the house keeping job before returning from the program.

Developing the FREEWAY FROG program

the logical breaks shown above.

stage is working before proceeding to the next stage.

The six stages will be

1. Data Base design

involving the design of objects shape, the creation of

2. Initialisation

involves the setting up of the screen, and the initialisation of various variables.

2. Traffic flow

different logic.

6. Police car

we develop and fight the police car movement.

3. EFOR

crashing, calculating scores ..etc.

6. Program control

game, return from the program.

block of memory and generate the sum as a "checksum".

9000 REM
9010 REM checksum

```
9020  INPLT "From address"
9030  INPUT "To address"
4      L= S
9050  FOR I=F TO t
9060  LET s=s+PEEK I
9070  NEXT I
9080  PRINT "Checksum = "
9090  GO TO 9020
```

Enter the starting block, number, and at the end of the memory block which you want to do the checksum in decimal value. The BASIC program will generate the checksum value.

Stage1-DataBase

**** Design of object shape ****

As this is a two way traffic game, we need to design two truck shapes a left truck shape and a right truck shape etc...

For the truck there will be four possible shapes and so there will be four shapes, one for each direction.

for drawing each object:

If the shape is composed of four characters

A B

the position pointer will be pointing to character A.

Character A drawn first, then character B ...until the whole row is finished.

... draw C D ...
above to character C

Thus, we will organize the shape data

Shape AB: D

Shape AB: D
A1, A2, A3, A4, A5, A6, A7, A8
B1, B2, B3, B4, B5, B6, B7, B8
C1, C2, C3, C4, C5, C6, C7, C8
D1, D2, D3, D4, D5, D6, D7, D8

Look like this

Shape a1, a2, a3, a4, a5, a6, a7, a8
b1 b2 b3, b4, b5, b6, b7, b8
c1, c2, c3, c4, c5, c6, c7, c8
d1, d2, d3, d4, d5, d6, d7, d8

If a character is copied by when we want to move it, then we change the attribute file.

We can change the attribute file after its screen memory data

Unlike the shape, for each character there is only one corresponding attribute data byte.

So, to cater for the attributes data we have four attribute data bytes after the above thirty two shape data bytes. (for a four character shape)

**** Input of object shape ****

Label	Line#	from(H) to(H)	from(D) to(D)
FROGSHP	120		27055 27190
LBIKE	340	6A17H 6A76H	27191 27254
LBATT	430	6A77H 6A7EH	27255 27262
RBIKE	460	6A7FH 6A81H	27263 27326
RBATT	560	6A8FH 6AC6H	27327 27334
LCAR	610	6AC7H 6B26H	27335 27430
LCATT	730	6B27H 6B32H	27431 27442
RCAR	770	6B31H 6B92H	27443 27538
RCATT	900	6B93H 6B9EH	27539 27550
TRUCK	940	6B9FH 6C76H	27551 27766
LTATT	1230	6C77H 6C91H	27767 27793
RTRUCK	1280	6C92H 6F49H	27794 28009
RTATT	1570	6D6AH 6D84H	28010 28016
RANK	1620	6F85H 6F88H	28037 28040

Module from 27055 to 28040, 986 bytes, checksum is 79197.

Suggested name "shapdb",(shape database).

data bytes followed by attribute data bytes.

one time, either GREEN when it is alive, or RED when it is dying,
or YELLOW when it reaches home.

(7) as paper colour.

O and the ink colour will be that given in its database.

cassette, it is assumed that you understand character representation in memory.

shape FROG1, starting at line 160.

In line 160, you will see

69B7 6F 160 FROG1 DB 111,15,31,159,220,2 6,120,48
OF 1F 9F DC DB 78 30

6987 is the memory address in hexadecimal format
 6F is the start of the eight bytes of the current DB instruction in Hexadecimal value.
 The hexadecimal value of the next seven bytes are in the next line between Line 160 and line 170, i.e DFB, 1FH, 9FH, DCH, D8H, 78H, 30H.
 160 is the line number of the assembler listing.
 FROG1 is the label. This is for our benefit only.
 DB is a mnemonic. It means that what follows is a sequence of bytes. (Similar to DATA in BASIC).
 +11,15,31,159,220,216,120,48
 are the bytes to be loaded into the memory.

Now let's build the FRM11 file.

00	00000000	00000000	"
01	00000001	10000000	A
23	00100011	11000000	
25	00.00101	10 00 00	+
6F	01101111	11.10110	F
4F	01001111	11110010	
7F	11011111	111.11.1	F
FF	11111111	111.11 1	03
6F	01101111	1111 110	F6
CF	01.101111	1111 0000	F
4F	00011111	1.11.000	F8
9F	10011111	1.111001	F9
DC	11011100	00111011	3B
D8	11011100	00011011	1B
78	01111000	00011110	1E
30	00110000	00001.00	0C

Remember,

- I. we draw the bottom row first from left to right.
- II. Then we draw the next row up.
- III. For each character, we draw the eight bytes from top to bottom.
- IV. Then at the very last, we fill in the attributes.

For example, the above binary representation of the characters in the string FROG1 is:

to find the correct shape given the direction of the frog.

DEFW is a mnemonic that means we want to define a 2-byte "nn". The least significant byte is first while the most significant byte is next.

**** Input of shape database ****

listing. Enter only the hex bytes as shown in column 2.

part of this stage)

**** Design of the objects database ****

in the two lane of the highway
These are randomly distributed between the two lanes

Object database will store information about the current status of the traffic

For example, for each object we need to know:

It's partly on the screen or not, Position pointer,
Shape database pointer, Attribute database pointer,
Number of Rows the shape occupies,
Number of column the shape occupies.

The dat carries this information about each object in each game cycle.

generated randomly

One simple way is to prepare the initial information for each possible vehicle and store this in memory.

memory locations and restore the data

program

* temporary database memory map

		x	x	x	v			
Existence		DEFB		1	byte			
Cycle count		DEFB		1	byte			
Direction		DEFB		1	byte			
Real/abstract		DEFB		1	byte			
Position		DEFW		2	bytes			
Shape pointer		DEFW		2	bytes			
Attribute		DEFW		2	bytes			
Row		DEFB		1	byte			
Column		DEFB		1	byte			
				TOTAL	12 bytes			
Label	line#	from(H)	to H	from(D)	to D			
M1EXT	1710	6E25H	6E30H	28197	28208			
OB2EXT	1800	6E31H	6E3CH	28109	28220			
OB3EXT	1850	6E3DH	6E48H	28221	28232			
OB4EXT	1900	6E49H	6F54H	28233	28244			
OB5EXT	1950	6F55H	6F60H	28245	28256			
OB6EXT	2000	6E61H	6F6CH	28257	28268			
PCAREXT	2070	6E6DH	6E78H	28269	28281			
FRYEXT	2180	6E79H	6F80H	28281	28288			

As mentioned above, these are only temporary working storage information that they contain changes as the game proceed.

There are two other major temporary working storage areas car respectively.

label	line#	from H) to(H)	From D) to D
FRGSTR	1650	6D89H	6DACH
PCSTR	1660	6DADH	6F24H

them. We only need to build up the following database.

The object database is organised in the following way:

"	right bycycle database
"	left bycycle database
"	right car database
"	left car database
"	right truck database
LTOB	left truck database

LPCDB	left police car database
TOPATT	top straight database
RHDB	right police car database
RPCATT	right police car database

label	Line#	from(H)	to(H)	from(D)	to(D)
FRGDB	2260	6E81H	6F80H	28289	28296
DBINDEX	2320	6E84H	6E94H	28297	28308
RHDB	2400	6E95H	6EA0H	28309	28320
LBDB	2470	6FA1H	6FACH	28321	28332
RCDB	2540	6EA1H	6FB8H	28333	28344
LCDB	2610	6FB9H	6FC4H	28345	28356
RTDB	2680	6EC5H	6ED0H	28357	28368
LTDB	2750	6FD1H	6FDCH	28369	28380
LPCDB	282D	6EDDH	6EF8H	28381	28392
LPCATT	2840	6EE9H	6EF4H	28393	28404
RPCDB	2930	6EF5H	6FO0H	28405	28416
RPCATT	3000	6FO1H	6FORH	28417	28428

Module from 28289 to 28428, 140 bytes, checksum 7697.

Suggested name is "objdb". (object database).

It is a
database.

The meaning and contents of each byte is:

* Existence (1 byte)

- set to zero when the object is nonexistent.
- set to value n where (n - 1) is the number of cycles that the object will wait before it is allowed to move
- n value for left and right cycle is 2
 left and right car is 3
 left and right truck is 6
 police car is 1
 frog is 8

In other words, the police car moves every cycle,
the motorcycle move every alternate cycle etc.

* Cycle count (1 byte)

- Initially set as 1 so that it is ready to move straight away and decrement by one every cycle.
- When it reaches zero, the object will be allowed to move and the count will be reinitialised to the value held in the existence byte

* Direction (1 byte)

- all left to right traffic (ie. top lane traffic) will have direction value zero.

- all right to left traffic (ie bottom lane traffic) will have direction value one.
- * Abstract/Real Flag (1 byte)
 - this defines whether objects is partly off the screen
 - all left to right traffic will start off with value zero (abstract).
 - left to right traffic will change this to one when their position points to the real screen 4820H.
 - all right to left traffic will have flag start off with value one (real); the object has a position pointing to the screen, ie 480FH
 - as the right to left traffic moves off the screen, i.e., when the position pointer moves from 480CH to 48BFH, this will be changed from real to abstract.
- * Position pointer (2 bytes)
 - 2 bytes pointer storing the current position of the object.
- * Shape pointer (2 bytes)
 - 2 bytes pointer pointing to the shape database of the object.
- * Attribute pointer (2 bytes)
 - 2 bytes pointer pointing to the attribute database of the object.
- * Row (1 byte)
 - store how many rows the object shape occupies.
- * Column (1 byte)
 - store how many columns the object shape takes
 - this value includes two columns of blanks, one at each end of the object.
 - The purpose of these two extra columns of blanks is to avoid the traffic getting too close to each other.

Now you can key in the object initialise data' e from listing 2270 to 30 0.

You can use EZ-code or Mex-load to enter this module.

If you use EZ-code, remember to save the source listing . . . > the dumped listing.

**** General database ****

We have covered so far the database from 69AFH to 6F0FH (21055 to 28428).

Now we are going to build up the rest of the database and we

classify this as "general database".

This is organised as below:

line	500 to 630	SEND
	660 to 690	SCORE MESSAGE
	720 to 1210	GENERAL

label	line#	from R) to, H	from D) to D)	checksum
PCTON1	500	6F0DH	6F10H	28429
PCINH2	510	6F11H	6F14H	28433
BOMTON	540	6F15H	6F30H	28437
SCRMS1	600	6F3DH	6F42H	28477
SCRRE	670	6F43H	6F48H	28483
SCRMS2	680	6F49H	6F53H	28489
BISCR	690	6F54H	6F58H	28500

Module from 28429 to 28504, 76 bytes, checksum 4813.

Suggested name : general (. , e).

You only need to input from line 500 to line 690.

From line 720 to line 1210, memory 6F59H to 6F82H
(28505 to 28546), these are all variables used by the program.

Line 1100 to 1150 are instructions with mnemonic EQU. This is a value to the corresponding label and is used by the program. You do not have to enter anything.

Conclusion

Now we have covered the whole database area from memory 69AFH TO 6F82H (27055 to 28546).

FREEWAY FROG program.

You should have now developed three modules.

name	from mem	to mem	length	checksum
shpdb	27055	28040	~	~
objdb	28289	28428	~	~
gendb	28429	28504	76	~ ~

Note that the database occupies nearly 1400 bytes!!

Stage 2- Initialisation

*** Screen Setup ***

In this module, we set up the highway, the score display, the frog as well as initialise all control variables.

We will do it in three parts.

First, clear the screen and put in the highway.

Secondly, put in all the frog

Thirdly, display the score.

This module includes the following routine

routine line#	from(H)	to(H)	from(D)	to(D)	
DRAWHwy	1820	2000H	7040H	28683	28736
FITHWY	2070	7041H	7054H	28737	28756

Module spread from 28347 to 31749, 2201 bytes.

Suggested name "init". (initialisation).

MEMORY locations:

Then enter INIT routine. Enter three bytes of zero for the called haven't been developed yet.

line#	address(H)	address(D)
1410	6FAFH	
1470	6FB8H	
490	6FC H	
1510	6FCFM	
1570	6FD6H	
1590	6FD8H	
1610	6FE7H	

code in memory 32000.

```

E3      D1      ;Disable interrupt
D9      EXX      ;Preserve HL
E5      PUSH    HL
D9      EXX
CDB8>6F  CALL    INIT
IE7F  KEY   LD     A,7FH      ,TRAP SPACE KEY
DBFF  IN     A,(FEH)
F601  AND    1
                  ;HL=1
C0FE?7  CALL    FINAL   ;finalisation
D9      EXX      ;restore HL
E1      POP    R
B9      EXX
FB      EI      ;enable interrupt
C9      RET

```

You should see the screen blacken and four white lines

w

N.T.

```

set border colour to black
initialise frog-crash flag, (frog existence, gametlag
number of frog
t random ROM pointer
t frog station (also initial position of frog) to
" A "
call clear screen
call draw highway
call line-up tr (five of them)
load score message
print score
load high score message
print high score
initialise all objects as nonexistent
initialise chase flag, siren sound flag and score

```

DRAW-WY

fill top highway line (32 characters of 40A0H)

fill middle highway line (32 characters of 4860H)

fill bottom highway line (32 characters of 5020H)

*remember that highway is white paper black ink

unfill top two-character bytes of top highway

(therefore, they are white)

unfill bottom two-character bytes of bottom highway

(they are also white now)

redraw middle two-byt of the middle highway

FINAL

```
    initialise fill character ()FFH
    set loop count to 32      {one line 32 characters}
    draw one character (8 bytes)
    move pointer to next character each time

FINAL
    set white border
    blank screen
    set screen attribute file to white paper and black ink
```

If everything is fine, save the module first from memory
location 28500 to 30800, 2300 bytes.

Now enter the RWF module. Type the break and save the
whole module again under the same name. Same addresses
from the hex dump from 6FAFH : 89 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
corresponds to line 1430 of the assembly listing
i.e. CD \$5 70.

Finally you will see five frogs come up at the bottom of
the screen.

The following are descriptions of what the two routines do:

LINFLP

```
set frog direction to 1 (facing right)
set frogs shape to FROG2
set attribute number to 2 (green)
if frog left
    then RETURN
else
    for number of frog
        push BC, DE, HL onto stack
        draw the frog by calling DRWFRC routine
        pop HL, DE, BC from stack
        update draw position
```

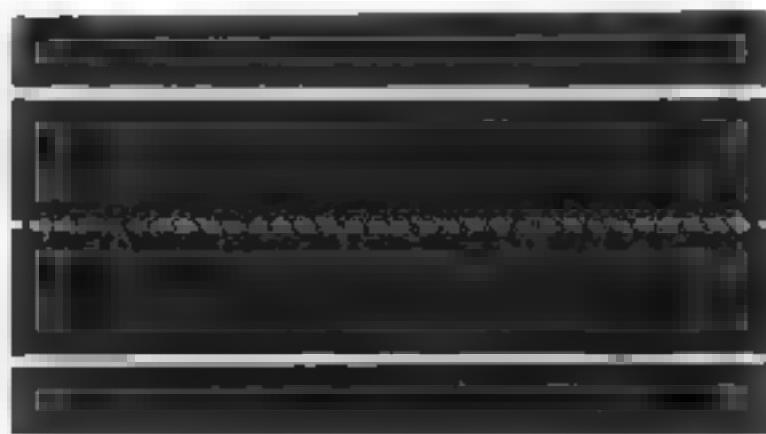
DRWFRC

```
draw shape using convention discussed earlier
calculate attribute pointer
fill attribute of frog
```

Y = ρ \hat{q} M τ n_{sc} b κ T k Rd c E
 μ β γ σ δ λ α

www.hinduismtoday.com | The EC by R. S. Agarwal | 2022-07-22

• **Figure 1** shows the results of the simulation of the effect of the number of nodes on the performance of the proposed scheme. The proposed scheme is compared with the scheme proposed by [1] and the scheme proposed by [2]. The proposed scheme shows better performance than the other two schemes.



REVIEW ARTICLE | **REVIEW ARTICLE** | **REVIEW ARTICLE** | **REVIEW ARTICLE**

ANSWER *Answers will vary.*

Digitized by srujanika@gmail.com



Stage 3 - Regular Traffic

In this stage, we develop the regular flow of traffic.
i.e. all traffic except police car

```
Traffic control (including regeneration of traffic)
    regenerate traffic
Moving traffic
    moving control
        swing traffic
        determine drawn shape
```

Below is a table of all routines in this module.

name	line#	from(H)	to(H)	from(D)	to(D)	func	param
TFCTRL	3090	70B0H	70D8H	28861	7888H	1-8	
RGEN	3320	70D9H	710EH	28889	4000H	5A-7	
MVTRF	3700	710FH	71AFH	28943	4000H	1-8	
MVCTRL	4720	71AFH	7208H	29103	4000H	*	
DRAW	5560	7209H	7295H	29193	4000H	1-7	
NSHAPE	6630	7296H	72D6H	29334	7000H	6H	
RANDOM	15050	77CCH	77D1H	30668	1-8H	2-7	

Module from 28861 to 30685, 1824 bytes.

Suggested name "regtrf". (regular traffic).

It is useless to generate a total checksum of the whole module.
use the memory range covers some undeveloped memory
and it is better to check the checksum after each
routine after entering it.

We develop this module in two parts.

Firstly, the draw routine for traffic.
Secondly, the traffic control and draw control.

Now we can write the two modules and save them.

Load the INIT module. Enter test routine program starting at memory 32000.

F3	D1
D9	EXX
E5	PUSH HL
D9	EXX
C0836F	CALL INIT
3E03	LD A, 3 ;ROW COUNT A.0H A, 41H <--> H.W A, 00H <--> H.W
325F6F	LD (COLUMN), A ;STORE IN COLUMN
11926C	LD DE, RTBLCK ;right truck shape
216A6D	LD RL, RTATT ;right truck attrl
226A6F	LD (ATTPTR), HL ;STORE IN ATTPTR
3E01	LD A, 1 ;set to real pos
212248	LD HL, 4822H ;
C0D972	CALL DRAW ;DRAW
3E78 KEY	LD A, 7FH ;KEY TRAP
DBFF	IN A,(OPEN)
E601	AND 1
20F8	JR NZ, KEY
FFE77	CALL FINAL
D9	EXX
E3	POP HL
D9	EXX
FB	EI
C9	RST

Load the INIT module in the order they are created.
Load the INIT module.

Load the routines you developed in this stage.

Load the routines you developed in previous stages. This includes:

Includes all routines you have developed so far.

Enter and save the above test routine in memory 32000.

Run the program. It will print a right truck in the top lane as well.

Save the current memory dump. Copy it in and run it in a monitor. Call INIT and CALL DRAW to test all other object shapes.

Below is a brief description of the two routines.

INIT

Similar logic to DRWFRG.

RESHAPE

trap lower 5 bits of low order byte of position parameter
subtract from 1FH and add 1

trap lower 5 bits again
determine SKIP and FILL depending on real or abstract
calculate attribute position and store in ATIPOS

save the whole module.

Edit the testing routine as follows

```
    N
    FXX
    PUSH   HI
    FXX
    CALL   INIT
CDBD70  MOVE   CALL   TFLTR,
CDDF71  CALL   MOVTRP
3E7F    LD     A, 7H
38FE    IN     A,(OPEN)
FB01    AND   1
20F2    JR    NZ, MOVE
        ALL   F NAJ
        XX
        FNP   H.
        XX
        F
        KIT
```

we are developing that stage

Once a module is fully developed and tested, it will be merged

12000

test run the new "frag" module.

plus all traffic moving at a very fast speed in the two lanes. This is because there is no delay between each program cycle.

A short description of the routine

TF TR,

load generation flag
 If not regeneration
 decrement flag count

else

regenerate the first nonexistence object by calling

the REGEN routine

REGEN

REGEN

```
    save existence database pointer
    generate random number 0 to 5
    test the first two character of the screen position
        where the object is created
    if the sum of the attributes of those two position is
        not equal to zero
        then return (traffic jam)
    else
        determine the initialise database
        load into temporary working database
        set regeneration cycle count to 2
        return
```

MVTRI

```
for all existing objects
    decrement cycle count
    if count reaches zero
        reload count from existence
        move one character left or right
        store new position in NFWPOS
        test attribute correspondence of the front of
            objects
        if any nonzero ink
            if not green
                set jam flag
            else
                set crash flag
        if jam flag set
            load cycle count with 2(move one cycle
                later)
            return
        else
            store new position
            call MVCTRL (move control)
```

MVCTRL

```
if edge reached
    change real/abstract flag
if left moving
    if on edge (position low order byte = 1FH)
        if abstract flag
            set non exist, return
        else
            goto L1
    else
```

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100

+ - * / = % ^ < > <= >= << >> <<= >>= <<< >>>

+

Enter an ATTPTR
to the column of shop

Help

what random pointer pointing to in ROM

F N D G H M

Score @ HIGH SCORE @



Score @ HIGH SCORE @



Stage 4-Police Car

In this stage we will introduce the Police car into the program.

The Police car will be generated randomly and enter with a RIN sounding "FROG" every time the car is on the road in its course. It will overtake any regular traffic before it.

This is what the program needs to save what is and rewrite the F1 car and put it back when the police car moves on.

Below are all routines in this module.

name	line#	from(H)	to(H)	from(D)	to(D)	checksum
RSPC	9560	7450H	74C1H	29726	29889	11011
POLICE	7930	734AH	73DEH	29514	29662	15769
STRPC	8810	73D9H	744FH	29663	29775	10615

For any code you may need to use, see the [code developed in previous modules](#).

Module from 29514 to 29889, 376 bytes

Suggested name "police".

Load the "frog" module first so it can hook modules like them onto the "police" module.

Then edit the testing program to the following:

```
        FXX
        PUSH    H1
        FXX
        CALL    INIT
        MOVE   CALL    TFCTR
        CALL    MOVTR
        \D4A73  MOVEI  CALL    POLICE
        .D     A,7FH
        IS     A,(OFFH)
        AND    I
        20E5  JR     NZ,MOVE.
        CALL    F_NA.
        EXX
        POP    H2
        FXX
        F1
        RFT
```

Load the "frog" module, then the "police" module.

fast on the highway.

If you want to put in other traffic as well, change the relative jump to JR NZ, MOVE. Remember to recalculate the displacement offset, (EFH).

shape as it overtakes them. It may be so fast that you wouldn't notice.

In Front of them.

Let's look at these routines

POLICE

```
if police car non-exist
    get a random number
    if not a multiple of 3,
        return
    else
        set chase flag
        determine top or bottom lane randomly
        load corresponding initial data
        t direction
        store position pointer
        retrieve position
        move and store NEWPOS
        set ROW, COLUMN, BEAL/ABSTRACT flag, POS before
            call RSHPAE
        get resulted ATTPOS and test head of shape for green
        if green attribute
            set crash flag
            blank front of police car
        call STRPC {for storing of what's underneath policecar}
        update position database
        call MVCTRL (for moving on and off the screen)
        turn off chase flag if non-exist
```

STRPC

```
set HL points to NEWPOS
set DE points to PCSTR (police car store)
store position and 3 byte of information starting from
    R/W variable
store according to SKIP/FILL format
    all screen memory first
    then attribute file
```

Save the whole module as "police"

You need to edit the testing program again to test the storing and restoring

Although we included the STRPC routine we are not sure that it

the screen

change the testing program to the following

MOVE	CALL	TFCTRI
LD5074	CALL	RFSPC
	ALD	MIVTKP
	ALL	POLICE
	TD	A,?FH
	TN	A,(OFFS)
	AND	1
ZOEC	JR	NZ,MOVE

Adjust the relative jump before you test run the module with 'Frog'

them off

The RFSPC's logic is as below

RFSPC

```
return if police car nonexistent
restore the position and 5 bytes into variables
    starting from ROW
restore screen memory then attribute according to
    SKIP FILL format
return
```

of the "police" module.

Load "frog" module then re-load the new "police" module.

Edit the testing program in memory 320000's following

```

        A L      N T
        S      TRCTR
        R      RSPN
        M      MOVTRP
        P      POLICE
        S      SEREN
        A      A,(OFFH)
        I
        F      MOVE

```

This program will find the whole frog
 This is a constant search for the
 spot + downcount delay loop

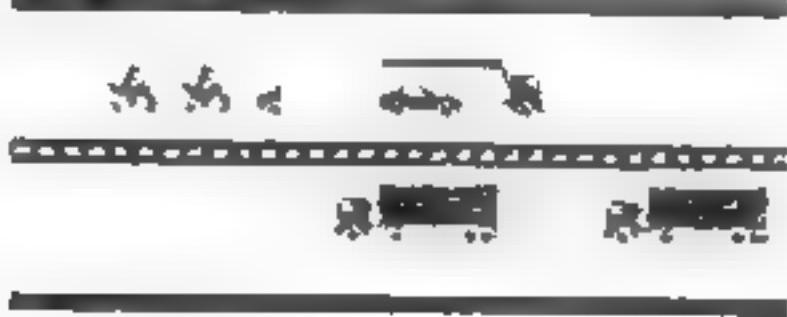
If you want to press

go to DELAY

IF no police

D DE, H
 11 01B

with "frog" in the memory and in
Score @ HIGH SCORE @



Stage 5-The Frog

In this stage we develop the frog routines.

We need to regenerate the frog when a frog dies.

We also need to handle the frog movement, save what is underneath it and restore what was stored back when the frog moves.

We also need to handle the frog position, home run and score etc.

We build up this module in three parts as below.

Regenerate and move frog
Save and restore what is underneath
Handle crashing, homerun and score

All routines within this module is as follows.

name	line#	from(H)	to(H)	from(D)	to(D)	checksum
FROG	10280	74C2H	74E2H	29890	29922	3818
REFRG	10320	74E3H	750FH	29923	29967	4079
MOVFRG	10770	7510H	75D5H	29968	30163	19943
RESFRG	11870	75D6H	7627H	30166	30247	8492
STRFRG	12440	7628H	7690H	30248	30352	10136
CRASH	13160	7691H	76A6H	30353	30374	2767
FRGDIE	13280	76A7H	7707H	30375	30471	9965
FRGTON	13890	7708H	771CH	30472	30492	2435
CALSCR	14040	771DH	776EH	30493	30574	8106

Module from 29890 to 30574, 685 bytes

Suggested name "frgrtn" (frog routine).

B1 R1 T1 P1 Z1 M1 L1 H1 R1 F1 C1 STAPS 4R 8A 1L 1R 1T

Edit the testing program as following.

```
        CALL    INIT
MOVE    CALL    TFCIR1
        CALL    RESPC
        AL     MI    TRF
        A_L    FP    A
        Z     R_A
```

```
LD      A,7FH
N      A,0FFEH
AND    I
JR    NZ,MovI
```

in line 13190 of the assembly listing by
30, 00, 00

Run the testing program and you should be able to move the frog.
The controls are "l"-up, "a"-down, "h"-left, "p"-right.

put in yet.

The description of these few routines are as following

FROI,

the control routine for the whole FROG module

```
IF frog crashed
    goto CRH

    if score-flag to no score (0)
    call RF_FRO
    decrement cycle count
    if count non zero
        return
    else
        reset cycle count
        call MKVFR
        if not crash
            return
    CRH call CHASH routine
    return
```

RF_FRO

```
If frog does not exist
    load frog initial database to working database
    update frog station to three position left
    initialise OLDFRG and NEWFRG to FR_PMS
    initialise frog storage area to 0
```

MGrFRC

```
initialise registers
    C - absolute movement
    B - frog direction
```

```
DE = frog shape
test Frog movement
    ) up, a - down, i - left, p - right
store shape, direction
if absolute move is zero
    return
else
    restore old frog position
    calculate new frog position and store
    test up screen position, right screen, bottom screen,
        frog station
    if valid
        store position into NEWFRG
        set score flag.
    restore OLDFRG
    if OLDFRG equal NEWFRG
        return

    call R_dN
    set OLDFRG equal NEWFRG,
    move stored direction, shape pointer to frog
        database
    call STRFR
    return
```

RESTFR

```
restore underneath frog based on OLDFRG position
memory first then attributes
```

STRFR()

```
store underneath frog based on NEWFRG position
draw frog while drawing
```

FRCDIE

```
reset crash flag
set frog nonexistent
call FRCDIE routine (dying procedures)
call RESTFRG routine (plate back what was underneath)
decrement number of Frog
```

and FRCTOM routine.

the following instruction

1698 CDA?76 CALL FRCDIE (30)60

Edit the testing program to the following

M. P. T. D.
together with the "frog" module

When the frog crosses, it will flash red and vanish.

FR-EIE

```
test frog reaches home or die
    if die tone, red colour attribute
        reaches home
            add one to third digit of score
                (bonus 100 points)
        call DESSCR (display score)
        set home tone, yellow colour attribute
        draw frog based on OLDFRC, FROUSH, and attribute just
            set by call DRWFRC
            h frog with the attribute five times
```

FR-TON

```
call TONE1 (tone code from SIREN routine)
move up or down the tone database depends upon attribute
    used to flash the frog ( if yellow then down db )
        ( if red then up db )
```

CASER

```
if frog non-exist
    return
else
    if score flag not set
        return
    else
        if go up
            add one to tenth digit of score
                (10 points)
        else
            if not within the highway
                return
            else
                add one to tenth digit
```

PlayStation
Sony Computer Entertainment America

Score 0 HIGH SCORE 0



Score 100 HIGH SCORE 0



Stage 6 - Control

When you get to the control line, type `Q` to quit.

After you quit with a key stroke, the score is updated and the game is restarted automatically.

To quit an active game, type `Q` to space key.

You will find that line 180 to 440 of the listing looks very similar to the testing program we have been developing.

Routines left for program are as follows:

name	line#	from H) to(H)	from(D) to(D)	checksum
START	180	6978H	69A9H	27000
OVFR	13200	77D9H	77FDH	30686

Now you can copy the module to memory and then link it with the "frog" module and save the whole module as "frog".

For now, just copy the module to memory and then link it with the "frog" module and save the whole module as "frog".

```
OVER
    compare 4 digits of HISCR and SCORE+1
    for the first nonequal digit
        if H CR digit is lower
            update HISCR to SCORE+1
    else
```

Score 1140 High Score 1140



Wednesday February 19th **1997** **10:00 AM**

HL	PRESERVE THE HL REGISTER DATA
IN I	INITIALISATION
REFLNL	TRAFFIC CONTROL POINTING TO TRAFFIC INPUT
END	END PROGRAM
POLICE	POWDER CAR ROUTINE
PIG	FIND PIGGY
AIHER	CALCULATE AND DISPLAY SCORE
N	MAIN PROGRAM
END	END OF PROGRAM

4. LỚN HƠN

LỊCH SỬ SCORE MANAGEMENT

THỜI KỲ MỚI CỦA TỔNG QUẢN TRỊ

THIẾT KẾ VÀ THỰC HIỆN

如上所述，本研究的实验结果表明，通过增加水的摄入量，可以有效降低尿酸水平。

450 J. Child Psychol. Psychiatr. 2000; 41: 447-459

END OF INTERRUPTS
RETURN TO BASIC BY SEPARATE

REFERENCES

44 77 5 44 11 11
177 6 44 11 11
77 5 44 11 11
11

DH 386, 340, 348, 749; 38, 22, 31, 52

• 10 • 91, 1, 75, 72, 111, 79, 223, 257.

• β $b_1 \neq 0$ b_0 a

卷之三

31 31, 31, 22, 22, 193, 113, 56

$$q = \lambda \cdot \delta | \quad q = \lambda_1 \quad x = q - \mu = \lambda_1$$

ਪ੍ਰਾਤਿ ਦੇਣਾ

⁷²⁸ 56, 113, 143, 79, 2, 127, 21, 31, 34

6017 00	01460	DB	O ₂ , N ₂ , H ₂ O, H ₂ , Ar, N ₂ O, He
6018 00	01470	DB	O ₂ , N ₂ , O ₃ , O ₂ , O ₂ , O ₃ , O

6022 00	01480	DB	O, C, O ₂ , N ₂ , H ₂ O, H ₂ , N ₂ O
6024 00	01490	DB	O, O, N ₂ , O ₂ , N ₂ O, H ₂ O, N ₂
6042 00	01500	DB	O ₂ , O, N ₂ , O ₃ , N ₂ O, H ₂ O, N ₂
6044 00	01510	DB	O, C, O ₂ , C ₂ H ₆ , H ₂ O, N ₂
6052 00	01520	DB	O, O, O ₂ , O, O ₃ , O ₂
6058 00	01530	DB	O, C, N, H ₂ O, O ₂ , O ₃
6062 00	01540	DB	O, O, O, O, O, O, O ₂ , O ₃

6066 00	01570 81477	DB	H ₂ S, S, S, O, S, S, S, C
6073 00	01580	DB	S, S, S, S, S, S, S, S, S
6076 00	01590	DB	S, S, S, S, S, S, O, O, S

6105 00	01620 28098	DB	C, O, O, C
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6026	01640	DB		
6028	01650 FR00274	DB	36	401844
	01660 F002818	DB	120	411104

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46 7 00	01730	DEFB	0	
46 10 00	01740	DEFB	0	
46 29 00 00	01750	DEFN	0	
46 29 00 00	01760	DEFM	0	
46 10 00 00	01770	DEFN	0	
46 29 00	01780	DEFR	0	
46 29 00	01790	DEFB	0	
46 29 00	01800 DEFB 00	DB	0, C, O, O	
46 29 00 00	01810	DB	0	170 4105 60 00 0000 F 20

6E39 0000	01830	DEFW	0
6E3B 00	01840	DB	0,0
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6E3D 00	01850 087EYT	DB	0,0,0,0
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6E42 0000	01870	DEFW	
6E43 0000	01880	DEFW	
6E47 00	01890	DB	
6F49 00	01900 089EYT	DB	
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6E4F 0000	01920	DEFW	
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6E59 0000	01960	DEFW	
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6E61 00	02000 089EYT	DB	
00 00 00			
6F65 0000	02110	DEFW	0
6F67 0000	02120	DEFW	0
6F69 0000	02130	DEFW	0
6E6B 00	02140	DB	0,0
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6F6E		0	1	0	
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6F77		0	0	0	
6E75		0	0	0	
6F77	"	0	0	0	
6E78	"	0	0	0	
6E79	"	0	0	0	(PADS Database)
6E7A		0	0	0	
6E7B	"	0	0	0	
6E7C		0	0	0	
6E7E		0	0	0	
6E80	"	0	0	0	
6E81 00	02150 089EYT	DB	0,0		(PRINT AL MPH T ON DE FREQ)
00 00 00					
6E84 0000	02160	DEFW	0,0,0,0		
6E85 0000	02170	DEFW	0,0,0,0		

• [View Details](#) | [Edit](#) | [Delete](#) | [Share](#)

LEFT BIKE DR
LEFT CAR DR
RIGHT TRUCK DR

한국어	영어	한국어	영어
한국어	Korean	영어	English
한국어	Korean	영어	English
한국어	Korean	영어	English
한국어	Korean	영어	English

• $\eta_1(T) = \delta_1(\mu_1)$

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• 21

STUDY + SCREEN

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R(H₂)₃, 4, 5-H₂O

REFUGIUS BYTES HOME REVERSE

RECEIVED FROM THE UNIVERSITY OF TORONTO LIBRARIES
BY CLAUDIO RODRIGUEZ

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JUL 19 1988
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SERIALS SECTION
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	DEFW	-	MF	+	IF	+	SC	-
	DEFW	-	MF	+	IF	+	SC	-
	DEFH	-	MF	+	IF	+	SC	-

EQU	01	10, 3B, 0A, 3F
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ENDL	4	10, 12B, 0A, 104
ENDU	4	
ENDU	4, 10	0, FF
ENDU	3	0

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For the first time, we have shown that the H_2O_2 produced by the $\text{Fe}^{2+}/\text{Fe}^{3+}$ redox couple can be reduced by the $\text{Fe}^{2+}/\text{Fe}^{3+}$ redox couple.

CLARIFICATION REQUESTED
PRESERVE ALL EVIDENCE
DO NOT DESTROY EVIDENCE
DO NOT DISTURB
DO NOT DISTURB

LEFT ALL CPU MAPPING

SET NO FOR THE CAR CHASE

SAT 614EN ON
INIT AL+SE SRNGE TD
ASC115 SPAN to 344

P₄ 0.75 0.00

$\mathcal{E} = \mathbb{L}(\mathbb{M}^{\mathbb{N}})$, $\mathcal{F} = \mathbb{L}(\mathbb{M}^{\mathbb{N}})$, $\mathcal{G} = \mathbb{L}(\mathbb{M}^{\mathbb{N}})$

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		10:00	AM

上 华南理工大学图书馆 国家图书馆数字资源

02761-1 drew all frogs left on the screen

中華書局影印
明刻本
周易

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	BD4	02890	AMP	DH	MOVE PTR
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	BD4	02910	BRA	A	
	CD3F	02920	BRA	P	
	CD3B	02930	ADL	A, BH	
	BD4	02940	LD	P, A11H	INIT FROG SHAPE ATTR
	T7	02950		HL, P	
	23	02970	INC	HL	TEST CHARACTER
	77	02980		D	
	ED42	02990	INC	HL, BC	DONE LINE UP
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10

DELETING THIS AS
THE MOVES OFF SCREEN

OFFICE CYCLE COUNT

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1. $\frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$
 2. $\frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$
 3. $\frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$

11. **What is the best way to handle a difficult customer?**

ADD A.A. 15 MAY 1947 BY R. H. L.

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7A0114F		E 0	08	
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2 IP ATHENS

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47. In the following sentence, the verb is underlined. The subject is italicized.

display the value from character set
M₁ - where DE₁ the message pointer
is the same after display
register as well.

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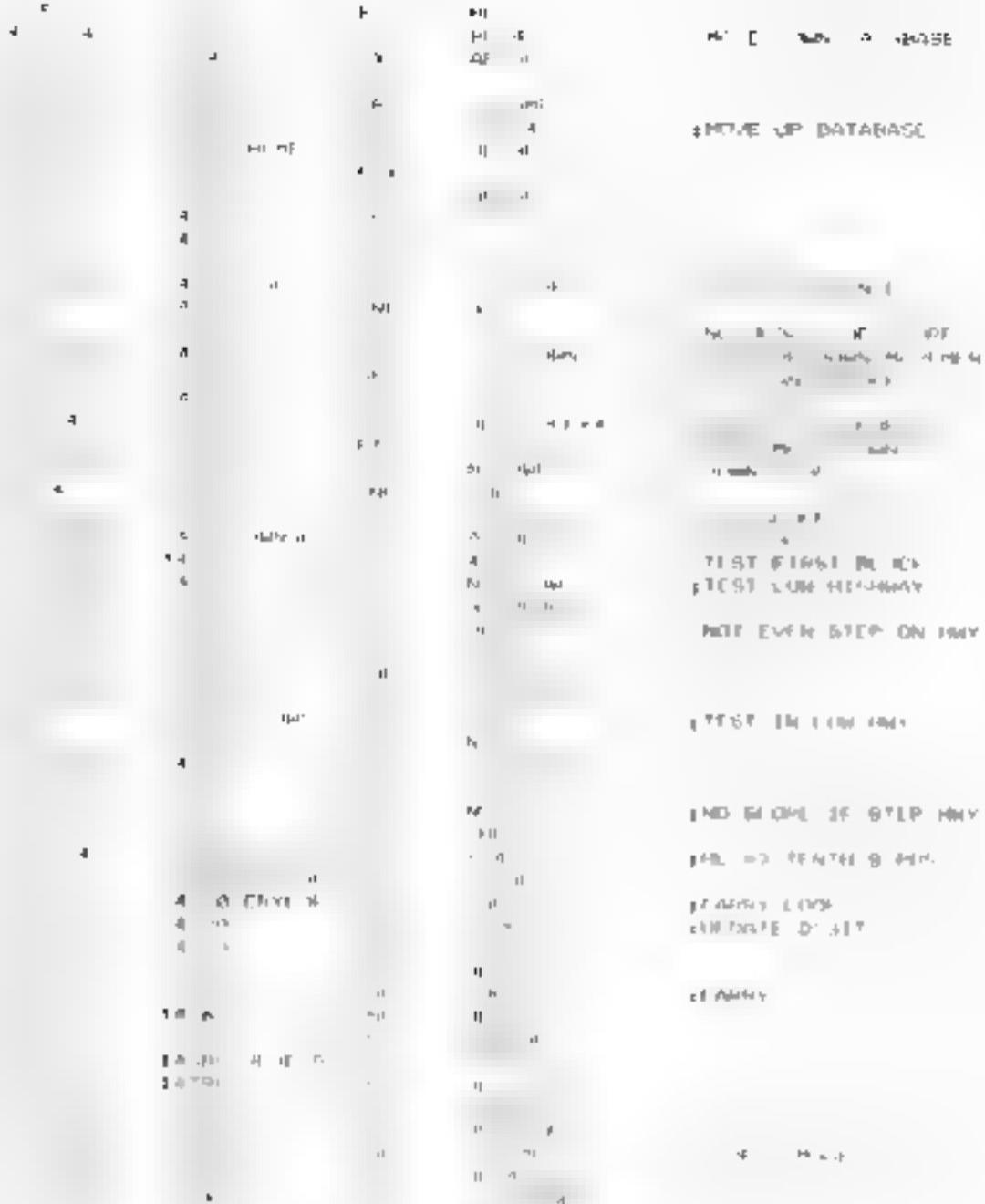
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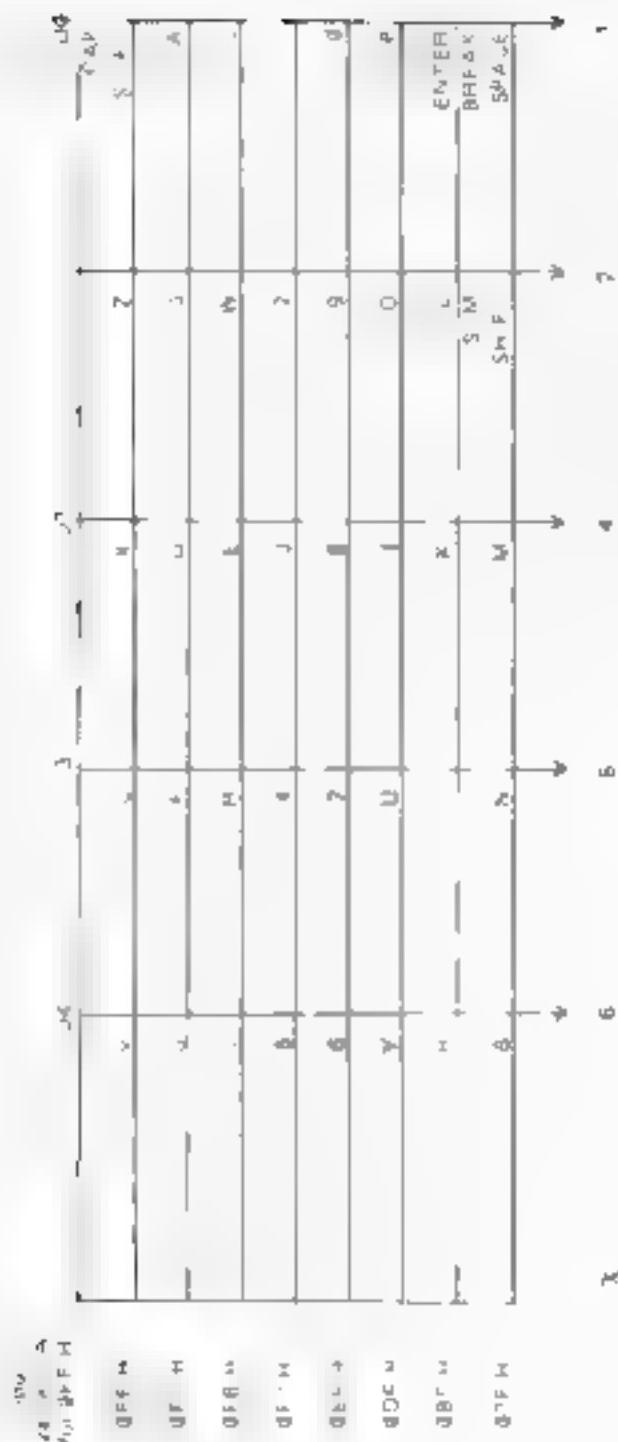
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313	314	315	316	317	318	319	320
321	322	323	324	325	326	327	328
329	330	331	332	333	334	335	336
337	338	339	340	341	342	343	344
345	346	347	348	349	350	351	352
353	354	355	356	357	358	359	360
361	362	363	364	365	366	367	368
369	370	371	372	373	374	375	376
377	378	379	380	381	382	383	384
385	386	387	388	389	390	391	392
393	394	395	396	397	398	399	400
401	402	403	404	405	406	407	408
409	410	411	412	413	414	415	416
417	418	419	420	421	422	423	424
425	426	427	428	429	430	431	432
433	434	435	436	437	438	439	440
441	442	443	444	445	446	447	448
449	450	451	452	453	454	455	456
457	458	459	460	461	462	463	464
465	466	467	468	469	470	471	472
473	474	475	476	477	478	479	480
481	482	483	484	485	486	487	488
489	490	491	492	493	494	495	496
497	498	499	500	501	502	503	504
505	506	507	508	509	510	511	512
513	514	515	516	517	518	519	520
521	522	523	524	525	526	527	528
529	530	531	532	533	534	535	536
537	538	539	540	541	542	543	544
545	546	547	548	549	550	551	552
553	554	555	556	557	558	559	560
561	562	563	564	565	566	567	568
569	570	571	572	573	574	575	576
577	578	579	580	581	582	583	584
585	586	587	588	589	590	591	592
593	594	595	596	597	598	599	600
601	602	603	604	605	606	607	608
609	610	611	612	613	614	615	616
617	618	619	620	621	622	623	624
625	626	627	628	629	630	631	632
633	634	635	636	637	638	639	640
641	642	643	644	645	646	647	648
649	650	651	652	653	654	655	656
657	658	659	660	661	662	663	664
665	666	667	668	669	670	671	672
673	674	675	676	677	678	679	680
681	682	683	684	685	686	687	688
689	690	691	692	693	694	695	696
697	698	699	700	701	702	703	704
705	706	707	708	709	710	711	712
713	714	715	716	717	718	719	720
721	722	723	724	725	726	727	728
729	730	731	732	733	734	735	736
737	738	739	740	741	742	743	744
745	746	747	748	749	750	751	752
753	754	755	756	757	758	759	760
761	762	763	764	765	766	767	768
769	770	771	772	773	774	775	776
777	778	779	780	781	782	783	784
785	786	787	788	789	790	791	792
793	794	795	796	797	798	799	800
801	802	803	804	805	806	807	808
809	810	811	812	813	814	815	816
817	818	819	820	821	822	823	824
825	826	827	828	829	830	831	832
833	834	835	836	837	838	839	840
841	842	843	844	845	846	847	848
849	850	851	852	853	854	855	856
857	858	859	860	861	862	863	864
865	866	867	868	869	870	871	872
873	874	875	876	877	878	879	880
881	882	883	884	885	886	887	888
889	890	891	892	893	894	895	896
897	898	899	900	901	902	903	904
905	906	907	908	909	910	911	912
913	914	915	916	917	918	919	920
921	922	923	924	925	926	927	928
929	930	931	932	933	934	935	936
937	938	939	940	941	942	943	944
945	946	947	948	949	950	951	952
953	954	955	956	957	958	959	960
961	962	963	964	965	966	967	968
969	970	971	972	973	974	975	976
977	978	979	980	981	982	983	984
985	986	987	988	989	990	991	992
993	994	995	996	997	998	999	1000



APPENDIX A

SPECTRUM KEY INPUT TABLE



VB : To trap a key

1. Load A with INPUT values of the corresponding row.
Bottom Row

VB : BREAK SPACE key
Normal state is always
JR Z, keyset

APPENDIX B

MEMORY ATTRIBUTE LINE

MEMORY ATTRIBUTE IN HEX	MEMORY ATTRIBUTE LINE IN HEX	MEMORY ATTRIBUTE IN HEX	MEMORY ATTRIBUTE IN HEX
4000	5800	0	0
4020	5820	1	0
4040	5840	2	0
4060	5860	3	0
4080	5880	4	0
40A0	58A0	5	0
40C0	58C0	6	0
40E0	58E0	7	0
40F0	58F0	8	0
4800	5900	9	0
4820	5920	9	0
4840	5940	10	0
4860	5960	11	0
4880	5980	12	0
48A0	59A0	13	0
48C0	59C0	14	0
48E0	59E0	15	0
5000	5A00	16	0
5020	5A20	17	0
5040	5A40	18	0
5060	5A60	19	0
5080	5A80	20	0
50A0	5AA0	21	0
50C0	5AC0	22	0
50E0	5AE0	23	0

APPENDIX C

SPECTRUM CHARACTER SET TABLE

Hex	Hex	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0H	0T3	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
1	0GUL	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
2	0D0	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
3	001	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
4	0'00	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
5	010	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
6	0'10	PRINT	EDIT	TAB ctrl	AT ctrl	~	6	6	6	6	6	6	6	6	6	6	6
7	0111	t000	t001	cursor left	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
8		t000	t001	cursor right	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
9				cursor down	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
A				cursor up	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
B				DELETE	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
C				ENTER	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
D				number	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
E					NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU	NU
F																	

NB NU Not Used

NON PRINTABLE ← → PRINTABLE

APPENDIX D

DECIMAL/MERADecimal CONVERSION TABLES

L	K	J	I	H	G	F	E	D	C
0	0	0	0	0	0	0	0	0	0
1	1	1	1	1	1	1	1	1	1
2	2	2	2	2	2	2	2	2	2
3	3	3	3	3	3	3	3	3	3
4	4	4	4	4	4	4	4	4	4
5	5	5	5	5	5	5	5	5	5
6	6	6	6	6	6	6	6	6	6
7	7	7	7	7	7	7	7	7	7
8	8	8	8	8	8	8	8	8	8
9	9	9	9	9	9	9	9	9	9
A	A	A	A	A	A	A	A	A	A
B	B	B	B	B	B	B	B	B	B
C	C	C	C	C	C	C	C	C	C
D	D	D	D	D	D	D	D	D	D
E	E	E	E	E	E	E	E	E	E
F	F	F	F	F	F	F	F	F	F

Appendix D

	w	x	y	z	
c	a	n	b	r	
0	0	0	0	0	0
1	0	0	0	0	1
2	0	0	0	1	0
3	0	0	0	1	1
4	0	0	1	0	0
5	0	0	1	0	1
6	0	0	1	1	0
7	0	0	1	1	1
8	0	1	0	0	0
9	0	1	0	0	1
10	0	1	0	1	0
11	0	1	0	1	1
12	0	1	1	0	0
13	0	1	1	0	1
14	0	1	1	1	0
15	0	1	1	1	1
16	1	0	0	0	0
17	1	0	0	0	1
18	1	0	0	1	0
19	1	0	0	1	1
20	1	0	1	0	0
21	1	0	1	0	1
22	1	0	1	1	0
23	1	0	1	1	1
24	1	1	0	0	0
25	1	1	0	0	1
26	1	1	0	1	0
27	1	1	0	1	1
28	1	1	1	0	0
29	1	1	1	0	1
30	1	1	1	1	0
31	1	1	1	1	1

We have to determine the 16-bit binary number,

i.e. bbbbbbbb bbbbbbbb
HOB LOB

	w	x	y	z	
c	a	n	b	r	
0	0	0	0	0	0
1	0	0	0	0	1
2	0	0	0	1	0
3	0	0	0	1	1
4	0	0	1	0	0
5	0	0	1	0	1
6	0	0	1	1	0
7	0	0	1	1	1
8	0	1	0	0	0
9	0	1	0	0	1
10	0	1	0	1	0
11	0	1	0	1	1
12	0	1	1	0	0
13	0	1	1	0	1
14	0	1	1	1	0
15	0	1	1	1	1
16	1	0	0	0	0
17	1	0	0	0	1
18	1	0	0	1	0
19	1	0	0	1	1
20	1	0	1	0	0
21	1	0	1	0	1
22	1	0	1	1	0
23	1	0	1	1	1
24	1	1	0	0	0
25	1	1	0	0	1
26	1	1	0	1	0
27	1	1	0	1	1
28	1	1	1	0	0
29	1	1	1	0	1
30	1	1	1	1	0
31	1	1	1	1	1

Four bits of the HOB (High Order Byte) to be 1 i.e. 01.

0001bbbb bbbbbbbb
HOB LOB

Since the difference is still greater than 255, we refer to the under the column heading 00x4 and find that 2104 is between 2048 and 2304. Again we take the lower value 2048 and arrive from the row value that the less significant four bytes of HOB is 8 i.e. 1000.

00011000 bbbbbbbb
HOB LOB

And then for the under the column heading 00x4 and find that 2104 is between 2048 and 2304. From the large middle big sub-table we find that 56 is at the intersection of row 3 and column 8. So we take the LOB as 38H.

10011000 0011100H
HOB LOB

So the HEX-value of the number 6200 is 1E18H.

卷之三

A COMPARISON OF CIVIL AND MILITARY DESIGN STANDARDS

<i>N</i>	<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>	<i>g</i>	<i>h</i>	<i>i</i>	<i>j</i>	<i>k</i>	<i>l</i>	<i>m</i>	<i>n</i>	<i>o</i>	<i>p</i>	<i>q</i>	<i>r</i>	<i>s</i>	<i>t</i>	<i>u</i>	<i>v</i>	<i>w</i>	<i>x</i>	<i>y</i>	<i>z</i>
H	1/21	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
A	1/6	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X
B	1/4	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
C	6/4	6/2	6/3	6/4	6/5	6/6	6/7	6/8	6/9	6/10	6/11	6/12	6/13	6/14	6/15	6/16	6/17	6/18	6/19	6/20	6/21	6/22	6/23	6/24	6/25	6/26	6/27
D	4/3	4/6	4/7	4/8	4/9	4/10	4/11	4/12	4/13	4/14	4/15	4/16	4/17	4/18	4/19	4/20	4/21	4/22	4/23	4/24	4/25	4/26	4/27	4/28	4/29	4/30	4/31
E	1/7	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26
F	6	5	4	3	2	1	0	-1	-2	-3	-4	-5	-6	-7	-8	-9	-10	-11	-12	-13	-14	-15	-16	-17	-18	-19	-20

HEXADECIMAL ADDITION TABLE

0		1		2		3		4		5		6		7		8		9		A		B		C		D		E		F		
0	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
1	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F	
2	2	3	4	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F		
3	3	4	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F			
4	4	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F				
5	5	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F					
6	6	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F						
7	7	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F							
8	8	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F								
9	9	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F									
A	A	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F										
B	B	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F											
C	C	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F												
D	D	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F													
E	E	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F														
F	F	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F															

APPENDIX G

FLAG OPERATION SUMMARY TABLE

INSTRUCTION	C	Z	P/V	S	N	H	COMMENTS
ADC HL SS	*	*	V	*	0	X	16-bit add with carry
ADX s AxDs	*	*	V	*	0	*	8-bit add or add with carry
ADD DD SS	*	-	-	-	0	X	16-bit add
AND s	*	*	P	*	0	1	Logical operations
BIT b,s	-	*	X	X	0	1	State of bit b of location s is copied into the Z flag
CCF	*	-	-	-	0	X	Complement carry
CPD CPDR CP CPIR	*	*	X	1	X		Block search instruction Z-1 if A=(HL) else Z 0 P/V+1 if BC>0, otherwise P/V=0
CP s	1	*	V	*	1	*	Compare accumulator
CPL	-	-	-	-	1	1	Complement accumulator
DAA	*	*	P	*	1	*	Decimal adjust accumulator
DECA	-	*	V	*	1	*	8-bit decrement
INC (C)	-	*	P	*	0	*	Input register indirect
INC s	-	*	V	*	0	*	8-bit increment
IND (INI)	-	*	X	X	1	X	Block input Z 0 if B=0 else Z=1
INDR INR	-	1	X	X	1	X	Block input Z 0 if B=0 else Z=1
LDA LDAR	*	FF	*	0	0	0	Content of interupt enable Flip Flop is copied into the P/V flag
LDD LDH	-	X	*	X	0	0	Block transfer instructions
LDHR LDH	-	X	0	X	0	0	P/V 1 if BC>0, otherwise P/V 0
NEG	*	*	V	*	1	*	Negate accumulator
OR s	0	*	P	*	0	0	Logical OR accumulator
OTDR OTIR	-	1	X	X	1	X	Block output, Z 0 if B=0 otherwise Z=1
OUTD OUTI	*	X	X	1	X		Block output Z 0 if B=0 otherwise Z=1
RLA RLCA RRA RRCA	*	-	-	-	0	0	Rotate accumulator
RLD RRD	*	-	P	*	0		Rotated quartet and eight
RLS RLCS RRH RRCH s	*	*	P	*	0	0	Rotate and shift locations
SLA s, SRA s, SRAL s	*	*	P	*	0		
SBC HL SS	*	*	V	*	1	X	16-bit subtract with carry
SCF	1	-	-	-	0	0	Set carry
SBC s, SUB s	-	-	V	-	1		8-bit subtract with carry
XOR x	0	-	P	*	0	0	Exclusive OR accumulator

Appendix G

INFORMATION

Carry flag. C=1 if the operation produced a carry from the most significant bit of the operand or result.

Zero flag. Z=1 if the result of the operation is

Sign flag. S=1 if the most significant bit of the result is one, i.e. a negative number

P/V Parity or overflow flag. Parity (P) and overflow (V) share the same flag. Logical operations affect this flag with the parity of the result while arithmetic operations affect this flag with the overflow of the result.

If P/V holds parity, P V 1 if the result of the operation is even, P V 0 if result is odd.

If P/V holds overflow, P V 1 If the result of the operation produced an overflow.

Carry flag. H 1 if the add or subtract operation produced a carry into or borrow from bit 8 of the accumulator.

Add Subtract flag. N 1 if the previous operations was a subtract.

H and N flags are used in conjunction with the decimal adjust instruction (DAA) to properly correct the result into packed BCD format following addition or subtraction using operands with packed BCD format.

The flag is affected according to the result of the operation.

The flag is unchanged by the operation.

The flag is reset (0) by the operation.

The flag is set (-1) by the operation.

The flag result is unknown.

The P/V flag is affected according to the overflow result of the operation.

P/V flag is affected according to the parity result of the operation.

Any one of the CPU registers A,B,C,D,E,H,L.

Any 8-bit location for all the addressing modes allowed for the particular instructions.

Any 16-bit location for all the addressing modes allowed for that instruction.

Refresh register

8-bit value in range 0-255.

16-bit value in range 0-65535

APPENDIX H
Z80 CPU INSTRUCTIONS SORTED BY OP-CODE

HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC
00	NOP	40	LD C,C	80	SUB D
01 XXXX	LD BC NN	4A	LD C,D	83	SUB E
02	LD (BC) A	4B	LD C,E	84	SUB H
03	INC BC	4C	LD C,H	85	SUB L
04	INC B	4D	LD C,L	86	SUB HL
05	DEC B	4E	LD C,HL	87	SUBA
06		4F	A	88	SH A,B
07 XXXX	RECA	50	LD D,B	89	SBC A,C
08	EXAF AF	51	LD D,C	8A	SBC A,D
09		52	L	8B	SH A
0A		53	F	8C	SBC A,H
0B		54	H	8D	SBC A,I
0C	INC C	55	LD D,L	8E	SBC A,HL
0D	DEC C	56	LD D,HL	8F	SBC A,A
0E XXXX		57	A	90	AND H
0F		58	D	91	AND
10 XXXX		59	C	92	ANI
11 XXXX		5A	A	93	ANL A
12 XXXX		5B	H	94	ANL H
13 XXXX		5C	L	95	ANL
14 XXXX		5D	B	96	ANL =
15 XXXX		5E	C	97	XOR H
16 XXXX		5F	D	98	XOR C
17 XXXX		60	E	99	XOR D
18 XXXX		61	F	9A	XORE
19 XXXX		62	G	9B	XORN
1A XXXX		63	H	9C	XORN A
1B XXXX		64	I	9D	DR B
1C XXXX		65	J	9E	DR C
1D XXXX		66	K	9F	DR D
1E XXXX		67	L	A0	DR E
1F XXXX		68	M	A1	DR H
20 XXXX		69	N	A2	DR L
21 XXXX		6A	O	A3	DR (HL)
22 XXXX		6B	P	A4	DR A
23 XXXX		6C	Q	A5	DR B
24 XXXX		6D	R	A6	DR C
25 XXXX		6E	S	A7	DR D
26 XXXX		6F	T	A8	DR E
27 XXXX		70	U	A9	DR H
28 XXXX	CPL	71	V	AA	DR L
29 XXXX	JR NC DIS	72	W	BA	CP (HL)
2A XXXX	LD SP NN	73	X	CA	N.
2B XXXX	LD (NN) A	74	Y	CB	POP BC
2C XXXX	INC SP	75	Z	CC	JP NZ NN
2D XXXX	INC (HL)	76		CD	JP VN
2E XXXX		77		CE	CALL NZ NN
2F XXXX		78		CF	PUSH BC
30 XXXX		79		CG	ADD A,N
31 XXXX		7A		CH	RET Z
32 XXXX		7B		CI	RET
33 XXXX		7C		CJ	JP Z NN
34 XXXX		7D		CK	CALL Z NN
35 XXXX		7E		CL	ADC A,N
36 XXXX		7F		CM	ADC A,N
37 XXXX		80		CO	ADC A,B
38 XXXX		81		CP	ADC A,C
39 XXXX		82		CA	ADC A,D
3A XXXX		83		CC	ADC A,E
3B XXXX		84		CD	ADC A,H
3C XXXX		85		CE	ADC A,L
3D XXXX		86		CF	ADC A,I
3E XXXX		87		CG	ADC A,HL
3F XXXX		88		CH	ADC A,A
40 XXXX		89		CI	ADC A,B
41 XXXX		8A		CC	ADC A,C
42 XXXX		8B		CD	ADC A,D
43 XXXX		8C		CE	ADC A,E
44 XXXX		8D		CF	ADC A,H
45 XXXX		8E		CG	ADC A,L
46 XXXX		8F		CH	ADC A,I
47 XXXX		90		CI	ADC A,HL
48 XXXX		91		CC	ADC A,A
			SUB C	DBXX	JL NN
				DBXX	IN A/N

HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC
00 00 00	ADD A,B	CB 00	SRA B	C 4	SHL
00 00 00	SBC A,M	CB 01	SRA C	C 5	SHL
00 00	PLST 10H	CB 02	SRA D	C 6	SHL
00	POP HL	CB 03	SRA E	C 7	SHL
00 00 00	RR	CB 04	SRA H	C 8	SHL
00 00 00	EX ISP HL	CB 05	SRA L	C 9	SHL
00 00 00	CALL PO,NW	CB 06	SRA (HL)	C 0A	SHL
00	L	CB 07	SRA A	C 0B	SHL
00 00 00	H	CB 08	A-H	C 0C	SHL
00 00 00	I	CB 09		C 0D	SHL
00 00 00	J	CB 0A		C 0E	SHL
00 00 00	K	CB 0B		C 0F	SHL
00 00 00	L	CB 0C		C 10	SHL
00 00 00	M	CB 0D		C 11	SHL
00 00 00	N	CB 0E		C 12	SHL
00 00 00	O	CB 0F		C 13	SHL
00 00 00	P	CB 10		C 14	SHL
00 00 00	Q	CB 11		C 15	SHL
00 00 00	R	CB 12		C 16	SHL
00 00 00	S	CB 13		C 17	SHL
00 00 00	T	CB 14		C 18	SHL
00 00 00	U	CB 15		C 19	SHL
00 00 00	V	CB 16		C 1A	SHL
00 00 00	W	CB 17		C 1B	SHL
00 00 00	X	CB 18		C 1C	SHL
00 00 00	Y	CB 19		C 1D	SHL
00 00 00	Z	CB 1A		C 1E	SHL
00 00 00	AA	CB 1B		C 1F	SHL
00 00 00	BB	CB 20		C 20	SHL
00 00 00	CC	CB 21		C 21	SHL
00 00 00	DD	CB 22		C 22	SHL
00 00 00	EE	CB 23		C 23	SHL
00 00 00	FF	CB 24		C 24	SHL
00 00 00	GG	CB 25		C 25	SHL
00 00 00	HH	CB 26		C 26	SHL
00 00 00	II	CB 27		C 27	SHL
00 00 00	JJ	CB 28		C 28	SHL
00 00 00	KK	CB 29		C 29	SHL
00 00 00	LL	CB 2A		C 2A	SHL
00 00 00	MM	CB 2B		C 2B	SHL
00 00 00	NN	CB 2C		C 2C	SHL
00 00 00	OO	CB 2D		C 2D	SHL
00 00 00	PP	CB 2E		C 2E	SHL
00 00 00	QQ	CB 2F		C 2F	SHL
00 00 00	RR	CB 30		C 20	SHL
00 00 00	SS	CB 31		C 21	SHL
00 00 00	TT	CB 32		C 22	SHL
00 00 00	UU	CB 33		C 23	SHL
00 00 00	VV	CB 34		C 24	SHL
00 00 00	WW	CB 35		C 25	SHL
00 00 00	XX	CB 36		C 26	SHL
00 00 00	YY	CB 37		C 27	SHL
00 00 00	ZZ	CB 38		C 28	SHL
00 00 00	AA	CB 39		C 29	SHL
00 00 00	BB	CB 3A		C 2A	SHL
00 00 00	CC	CB 3B		C 2B	SHL
00 00 00	DD	CB 3C		C 2C	SHL
00 00 00	EE	CB 3D		C 2D	SHL
00 00 00	FF	CB 3E		C 2E	SHL
00 00 00	GG	CB 3F		C 2F	SHL
00 00 00	HH	CB 40		C 20	SHL
00 00 00	II	CB 41		C 21	SHL
00 00 00	JJ	CB 42		C 22	SHL
00 00 00	KK	CB 43		C 23	SHL
00 00 00	LL	CB 44		C 24	SHL
00 00 00	MM	CB 45		C 25	SHL
00 00 00	NN	CB 46		C 26	SHL
00 00 00	OO	CB 47		C 27	SHL
00 00 00	PP	CB 48		C 28	SHL
00 00 00	QQ	CB 49		C 29	SHL
00 00 00	RR	CB 4A		C 2A	SHL
00 00 00	SS	CB 4B		C 2B	SHL
00 00 00	TT	CB 4C		C 2C	SHL
00 00 00	UU	CB 4D		C 2D	SHL
00 00 00	VV	CB 4E		C 2E	SHL
00 00 00	WW	CB 4F		C 2F	SHL
00 00 00	XX	CB 50		C 20	SHL
00 00 00	YY	CB 51		C 21	SHL
00 00 00	ZZ	CB 52		C 22	SHL
00 00 00	AA	CB 53		C 23	SHL
00 00 00	BB	CB 54		C 24	SHL
00 00 00	CC	CB 55		C 25	SHL
00 00 00	DD	CB 56		C 26	SHL
00 00 00	EE	CB 57		C 27	SHL
00 00 00	FF	CB 58		C 28	SHL
00 00 00	GG	CB 59		C 29	SHL
00 00 00	HH	CB 5A		C 2A	SHL
00 00 00	II	CB 5B		C 2B	SHL
00 00 00	JJ	CB 5C		C 2C	SHL
00 00 00	KK	CB 5D		C 2D	SHL
00 00 00	LL	CB 5E		C 2E	SHL
00 00 00	MM	CB 5F		C 2F	SHL
00 00 00	NN	CB 60		C 20	SHL
00 00 00	OO	CB 61		C 21	SHL
00 00 00	PP	CB 62		C 22	SHL
00 00 00	QQ	CB 63		C 23	SHL
00 00 00	RR	CB 64		C 24	SHL
00 00 00	SS	CB 65		C 25	SHL
00 00 00	TT	CB 66		C 26	SHL
00 00 00	UU	CB 67		C 27	SHL
00 00 00	VV	CB 68		C 28	SHL
00 00 00	WW	CB 69		C 29	SHL
00 00 00	XX	CB 6A		C 2A	SHL
00 00 00	YY	CB 6B		C 2B	SHL
00 00 00	ZZ	CB 6C		C 2C	SHL
00 00 00	AA	CB 6D		C 2D	SHL
00 00 00	BB	CB 6E		C 2E	SHL
00 00 00	CC	CB 6F		C 2F	SHL
00 00 00	DD	CB 70		C 20	SHL
00 00 00	EE	CB 71		C 21	SHL
00 00 00	FF	CB 72		C 22	SHL
00 00 00	GG	CB 73		C 23	SHL
00 00 00	HH	CB 74		C 24	SHL
00 00 00	II	CB 75		C 25	SHL
00 00 00	JJ	CB 76		C 26	SHL
00 00 00	KK	CB 77		C 27	SHL
00 00 00	LL	CB 78		C 28	SHL

HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC
A-7		A Y X	X A	4 X X X	S N E
F-4	H	4 X X	L X	5	
F-5	L	C X X	L L X	6	
CBC6	4 L	6 X X X	4 X	7	
CBC7	D	6 X X X	X X	8	
CBC8	E	L X X	X B	9	
CBC9		X X X	X C	10	
CBCA		X X X	X D	11	
CBCB		X X X	X E	12	
BCC	F	4 X X	X F	13	
		4 X X	X G	14	
		X X X	X H	15	
		X X X	X I	16	
		X X X	X J	17	
		X X X	X K	18	
		X X X	X L	19	
		X X X	X M	20	
		X X X	X N	21	
		X X X	X O	22	
		X X X	X P	23	
		X X X	X Q	24	
		X X X	X R	25	
		X X X	X S	26	
		X X X	X T	27	
		X X X	X U	28	
		X X X	X V	29	
		X X X	X W	30	
		X X X	X X	31	
		X X X	X Y	32	
		X X X	X Z	33	
		X X X	X A	34	
		X X X	X B	35	
		X X X	X C	36	
		X X X	X D	37	
		X X X	X E	38	
		X X X	X F	39	
		X X X	X G	40	
		X X X	X H	41	
		X X X	X I	42	
		X X X	X J	43	
		X X X	X K	44	
		X X X	X L	45	
		X X X	X M	46	
		X X X	X N	47	
		X X X	X O	48	
		X X X	X P	49	
		X X X	X Q	50	
		X X X	X R	51	
		X X X	X S	52	
		X X X	X T	53	
		X X X	X U	54	
		X X X	X V	55	
		X X X	X W	56	
		X X X	X X	57	
		X X X	X Y	58	
		X X X	X Z	59	
		X X X	X A	60	
		X X X	X B	61	
		X X X	X C	62	
		X X X	X D	63	
		X X X	X E	64	
		X X X	X F	65	
		X X X	X G	66	
		X X X	X H	67	
		X X X	X I	68	
		X X X	X J	69	
		X X X	X K	70	
		X X X	X L	71	
		X X X	X M	72	
		X X X	X N	73	
		X X X	X O	74	
		X X X	X P	75	
		X X X	X Q	76	
		X X X	X R	77	
		X X X	X S	78	
		X X X	X T	79	
		X X X	X U	80	
		X X X	X V	81	
		X X X	X W	82	
		X X X	X X	83	
		X X X	X Y	84	
		X X X	X Z	85	
		X X X	X A	86	
		X X X	X B	87	
		X X X	X C	88	
		X X X	X D	89	
		X X X	X E	90	
		X X X	X F	91	
		X X X	X G	92	
		X X X	X H	93	
		X X X	X I	94	
		X X X	X J	95	
		X X X	X K	96	
		X X X	X L	97	
		X X X	X M	98	
		X X X	X N	99	
		X X X	X O	100	
		X X X	X P	101	
		X X X	X Q	102	
		X X X	X R	103	
		X X X	X S	104	
		X X X	X T	105	
		X X X	X U	106	
		X X X	X V	107	
		X X X	X W	108	
		X X X	X X	109	
		X X X	X Y	110	
		X X X	X Z	111	
		X X X	X A	112	
		X X X	X B	113	
		X X X	X C	114	
		X X X	X D	115	
		X X X	X E	116	
		X X X	X F	117	
		X X X	X G	118	
		X X X	X H	119	
		X X X	X I	120	
		X X X	X J	121	
		X X X	X K	122	
		X X X	X L	123	
		X X X	X M	124	
		X X X	X N	125	
		X X X	X O	126	
		X X X	X P	127	
		X X X	X Q	128	
		X X X	X R	129	
		X X X	X S	130	
		X X X	X T	131	
		X X X	X U	132	
		X X X	X V	133	
		X X X	X W	134	
		X X X	X X	135	
		X X X	X Y	136	
		X X X	X Z	137	
		X X X	X A	138	
		X X X	X B	139	
		X X X	X C	140	
		X X X	X D	141	
		X X X	X E	142	
		X X X	X F	143	
		X X X	X G	144	
		X X X	X H	145	
		X X X	X I	146	
		X X X	X J	147	
		X X X	X K	148	
		X X X	X L	149	
		X X X	X M	150	
		X X X	X N	151	
		X X X	X O	152	
		X X X	X P	153	
		X X X	X Q	154	
		X X X	X R	155	
		X X X	X S	156	
		X X X	X T	157	
		X X X	X U	158	
		X X X	X V	159	
		X X X	X W	160	
		X X X	X X	161	
		X X X	X Y	162	
		X X X	X Z	163	
		X X X	X A	164	
		X X X	X B	165	
		X X X	X C	166	
		X X X	X D	167	
		X X X	X E	168	
		X X X	X F	169	
		X X X	X G	170	
		X X X	X H	171	
		X X X	X I	172	
		X X X	X J	173	
		X X X	X K	174	
		X X X	X L	175	
		X X X	X M	176	
		X X X	X N	177	
		X X X	X O	178	
		X X X	X P	179	
		X X X	X Q	180	
		X X X	X R	181	
		X X X	X S	182	
		X X X	X T	183	
		X X X	X U	184	
		X X X	X V	185	
		X X X	X W	186	
		X X X	X X	187	
		X X X	X Y	188	
		X X X	X Z	189	
		X X X	X A	190	
		X X X	X B	191	
		X X X	X C	192	
		X X X	X D	193	
		X X X	X E	194	
		X X X	X F	195	
		X X X	X G	196	
		X X X	X H	197	
		X X X	X I	198	
		X X X	X J	199	
		X X X	X K	200	
		X X X	X L	201	
		X X X	X M	202	
		X X X	X N	203	
		X X X	X O	204	
		X X X	X P	205	
		X X X	X Q	206	
		X X X	X R	207	
		X X X	X S	208	
		X X X	X T	209	
		X X X	X U	210	
		X X X	X V	211	
		X X X	X W	212	
		X X X	X X	213	
		X X X	X Y	214	
		X X X	X Z	215	
		X X X	X A	216	
		X X X	X B	217	
		X X X	X C	218	
		X X X	X D	219	
		X X X	X E	220	
		X X X	X F	221	
		X X X	X G	222	
		X X X	X H	223	
		X X X	X I	224	
		X X X	X J	225	
		X X X	X K	226	
		X X X	X L	227	
		X X X	X M	228	
		X X X	X N	229	
		X X X	X O	230	
		X X X	X P	231	
		X X X	X Q	232	
		X X X	X R	233	
		X X X	X S	234	
		X X X	X T	235	
		X X X	X U	236	
		X X X	X V	237	
		X X X	X W	238	
		X X X	X X	239	
		X X X	X Y	240	
		X X X	X Z	241	
		X X X	X A	242	
		X X X	X B	243	
		X X X	X C	244	
		X X X	X D	245	
		X X X	X E	246	
		X X X	X F	247	
		X X X	X G	248	
		X X X	X H	249	
		X X X	X I	250	
		X X X	X J	251	
		X X X	X K	252	
		X X X	X L	253	
		X X X	X M	254	
		X X X	X N	255	
		X X X	X O	256	
		X X X	X P	257	
		X X X	X Q	258	
		X X X	X R	259	
		X X X	X S	260	
		X X X	X T	261	
		X X X	X U	262	
		X X X	X V	263	
		X X X	X W	264	
		X X X	X X	265	
		X X X	X Y	266	
		X X X	X Z	267	
		X X X	X A	268	
		X X X	X B	269	
		X X X	X C	270	
		X X X	X D	271	
		X X X	X E	272	
		X X X	X F	273	
		X X X	X G	274	
		X X X	X H	275	
		X X X	X I	276	
		X X X	X J	277	
		X X X	X K	278	
		X X X	X L	279	
		X X X	X M	280	
		X X X	X N	281	
		X X X	X O	282	
		X X X	X P	283	
		X X X	X Q	284	
		X X X	X R	285	
		X X X	X S	286	
		X X X	X T	287	
		X X X	X U	288	
		X X X	X V	289	
		X X X	X W	290	
		X X X	X X	291	
		X X X	X Y	292	
		X X X	X Z	293	
		X X X	X A	294	
		X X X	X B	295	
		X X X	X C	296	
		X X X	X D	297	
		X X X	X E		

APPENDIX I

Z80-CPU INSTRUCTIONS SORTED BY MNEMONIC

MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL
LD A P	ED XX XX	LD B HL	0 XXXX	LD DR	ED B8
LD A R	ED XX XX	LD C HL	4E	LD IR	ED 90
LD ADDR DA	ED XX XX	LD D HL	DD 4E XX	NEG	ED 44
LD ADDR EA	ED XX XX	LD E HL	FD 4E XX	NOP	00
LD ADDR RA	ED XX XX	LD F HL	4F	OR H	B2
LD ADDR IB	ED XX XX	LD G HL	48	OR L	DD B6 XX
LD ADDR IC	ED XX XX	LD H HL	49	OR M	FD B6 XX
LD ADDR ID	ED XX XX	LD I HL	4A	OR A	B7
LD ADDR IA	ED XX XX	LD J HL	4B	OR B	BD
LD ADDR IB	ED XX XX	LD K HL	4C	OR C	B8
LD ADDR IC	ED XX XX	LD L HL	4D	OR D	BB
LD ADDR ID	ED XX XX	LD M HL	4E	OUT (C) A	ED 70
LD ADDR IA	ED XX XX	LD N HL	4F	OUT (C) C	ED 41
LD ADDR IB	ED XX XX	LD O HL	50	OUT (C) D	ED 51
LD ADDR IC	ED XX XX	LD P HL	51	OUT (C) E	ED 50
LD ADDR ID	ED XX XX	LD Q HL	52	OUT (C) L	ED 69
LD ADDR SP	ED 73 XXXX	LD D D	53	OUT part,A	ED 60
LD BC A	02	LD D E	54		
LD DE A	12	LD D M	55		
LD HL A	77	LD D L	56		
LD HL B	70	LD E M	57		
LD HL C	71	LD E L	58		
LD HL D	72	LD D MM	59		
LD HL E	56 XX	LD E (HL)	5A		
LD HL F	-	LD D MM	5B		
LD IXde A	ED 77 XX	LD E MM	5C		
LD IXde B	77	LD E E	5D		
LD IXde C	78	LD E M	5E		
LD IXde D	79	LD E L	5F		
LD IXde E	7A	LD E MM	60		
LD IXde F	7B	LD E E	61		
LD IXde G	7C	LD E M	62		
LD IXde H	7D	LD E L	63		
LD IXde I	7E	LD E MM	64		
LD IXde J	7F	LD E E	65		
LD IXde K	80	LD E M	66		
LD IXde L	81	LD H, HL	67		
LD IXde M	82	LD H, IXdeH	68		
LD IXde N	83	LD H, IXdeL	69		
LD IXde O	84	LD H (IXdeH)	6A		
LD IXde P	85	LD H (IXdeL)	6B		
LD IXde Q	86	LD H, HL	6C		
LD IXde R	87	LD H, IXdeH	6D		
LD IXde S	88	LD H, IXdeL	6E		
LD IXde T	89	LD H (IXdeH)	6F		
LD IXde U	8A	LD H (IXdeL)	70		
LD IXde V	8B	LD H, HL	71		
LD IXde W	8C	LD H, IXdeH	72		
LD IXde X	8D	LD H, IXdeL	73		
LD IXde Y	8E	LD H (IXdeH)	74		
LD IXde Z	8F	LD H (IXdeL)	75		
LD IXde A	8A	LD H, HL	76		
LD IXde B	8B	LD H, IXdeH	77		
LD IXde C	8C	LD H, IXdeL	78		
LD IXde D	8D	LD H (IXdeH)	79		
LD IXde E	8E	LD H (IXdeL)	7A		
LD IXde F	8F	LD H, HL	7B		
LD IXde G	90	LD H, IXdeH	7C		
LD IXde H	91	LD H, IXdeL	7D		
LD IXde I	92	LD H (IXdeH)	7E		
LD IXde J	93	LD H (IXdeL)	7F		
LD IXde K	94	LD H, HL	80		
LD IXde L	95	LD H, IXdeH	81		
LD IXde M	96	LD H, IXdeL	82		
LD IXde N	97	LD H (IXdeH)	83		
LD IXde O	98	LD H (IXdeL)	84		
LD IXde P	99	LD H, HL	85		
LD IXde Q	9A	LD H, IXdeH	86		
LD IXde R	9B	LD H, IXdeL	87		
LD IXde S	9C	LD H (IXdeH)	88		
LD IXde T	9D	LD H (IXdeL)	89		
LD IXde U	9E	LD H, HL	8A		
LD IXde V	9F	LD H, IXdeH	8B		
LD IXde W	90	LD H, IXdeL	8C		
LD IXde X	91	LD H (IXdeH)	8D		
LD IXde Y	92	LD H (IXdeL)	8E		
LD IXde Z	93	LD H, HL	8F		
LD IXde A	94	LD H, IXdeH	90		
LD IXde B	95	LD H, IXdeL	91		
LD IXde C	96	LD H (IXdeH)	92		
LD IXde D	97	LD H (IXdeL)	93		
LD IXde E	98	LD H, HL	94		
LD IXde F	99	LD H, IXdeH	95		
LD IXde G	9A	LD H, IXdeL	96		
LD IXde H	9B	LD H (IXdeH)	97		
LD IXde I	9C	LD H (IXdeL)	98		
LD IXde J	9D	LD H, HL	99		
LD IXde K	9E	LD H, IXdeH	9A		
LD IXde L	9F	LD H, IXdeL	9B		
LD IXde M	90	LD H (IXdeH)	9C		
LD IXde N	91	LD H (IXdeL)	9D		
LD IXde O	92	LD H, HL	9E		
LD IXde P	93	LD H, IXdeH	9F		
LD IXde Q	94	LD H, IXdeL	90		
LD IXde R	95	LD H (IXdeH)	91		
LD IXde S	96	LD H (IXdeL)	92		
LD IXde T	97	LD H, HL	93		
LD IXde U	98	LD H, IXdeH	94		
LD IXde V	99	LD H, IXdeL	95		
LD IXde W	9A	LD H (IXdeH)	96		
LD IXde X	9B	LD H (IXdeL)	97		
LD IXde Y	9C	LD H, HL	98		
LD IXde Z	9D	LD H, IXdeH	99		
LD IXde A	9E	LD H, IXdeL	90		
LD IXde B	9F	LD H (IXdeH)	91		
LD IXde C	90	LD H (IXdeL)	92		
LD IXde D	91	LD H, HL	93		
LD IXde E	92	LD H, IXdeH	94		
LD IXde F	93	LD H, IXdeL	95		
LD IXde G	94	LD H (IXdeH)	96		
LD IXde H	95	LD H (IXdeL)	97		
LD IXde I	96	LD H, HL	98		
LD IXde J	97	LD H, IXdeH	99		
LD IXde K	98	LD H, IXdeL	90		
LD IXde L	99	LD H (IXdeH)	91		
LD IXde M	9A	LD H (IXdeL)	92		
LD IXde N	9B	LD H, HL	93		
LD IXde O	9C	LD H, IXdeH	94		
LD IXde P	9D	LD H, IXdeL	95		
LD IXde Q	9E	LD H (IXdeH)	96		
LD IXde R	9F	LD H (IXdeL)	97		
LD IXde S	90	LD H, HL	98		
LD IXde T	91	LD H, IXdeH	99		
LD IXde U	92	LD H, IXdeL	90		
LD IXde V	93	LD H (IXdeH)	91		
LD IXde W	94	LD H (IXdeL)	92		
LD IXde X	95	LD H, HL	93		
LD IXde Y	96	LD H, IXdeH	97		
LD IXde Z	97	LD H, IXdeL	98		
LD IXde A	98	LD H (IXdeH)	99		
LD IXde B	99	LD H, HL	90		
LD IXde C	9A	LD H, IXdeH	91		
LD IXde D	9B	LD H, IXdeL	92		
LD IXde E	9C	LD H, HL	93		
LD IXde F	9D	LD H, IXdeH	94		
LD IXde G	9E	LD H, IXdeL	95		
LD IXde H	9F	LD H, HL	96		
LD IXde I	90	LD H, IXdeH	97		
LD IXde J	91	LD H, IXdeL	98		
LD IXde K	92	LD H, HL	99		
LD IXde L	93	LD H, IXdeH	90		
LD IXde M	94	LD H, IXdeL	91		
LD IXde N	95	LD H, HL	92		
LD IXde O	96	LD H, IXdeH	93		
LD IXde P	97	LD H, IXdeL	94		
LD IXde Q	98	LD H, HL	95		
LD IXde R	99	LD H, IXdeH	96		
LD IXde S	9A	LD H, IXdeL	97		
LD IXde T	9B	LD H, HL	98		
LD IXde U	9C	LD H, IXdeH	99		
LD IXde V	9D	LD H, IXdeL	90		
LD IXde W	9E	LD H, HL	91		
LD IXde X	9F	LD H, IXdeH	92		
LD IXde Y	90	LD H, IXdeL	93		
LD IXde Z	91	LD H, HL	94		
LD IXde A	92	LD H, IXdeH	95		
LD IXde B	93	LD H, IXdeL	96		
LD IXde C	94	LD H, HL	97		
LD IXde D	95	LD H, IXdeH	98		
LD IXde E	96	LD H, IXdeL	99		
LD IXde F	97	LD H, HL	90		
LD IXde G	98	LD H, IXdeH	91		
LD IXde H	99	LD H, IXdeL	92		
LD IXde I	9A	LD H, HL	93		
LD IXde J	9B	LD H, IXdeH	94		
LD IXde K	9C	LD H, IXdeL	95		
LD IXde L	9D	LD H, HL	96		
LD IXde M	9E	LD H, IXdeH	97		
LD IXde N	9F	LD H, IXdeL	98		
LD IXde O	90	LD H, HL	99		
LD IXde P	91	LD H, IXdeH	90		
LD IXde Q	92	LD H, IXdeL	91		
LD IXde R	93	LD H, HL	94		
LD IXde S	94	LD H, IXdeH	95		
LD IXde T	95	LD H, IXdeL	96		
LD IXde U	96	LD H, HL	97		
LD IXde V	97	LD H, IXdeH	98		
LD IXde W	98	LD H, IXdeL	99		
LD IXde X	99	LD H, HL	90		
LD IXde Y	9A	LD H, IXdeH	91		
LD IXde Z	9B	LD H, IXdeL	92		
LD IXde A	9C	LD H, HL	93		
LD IXde B	9D	LD H, IXdeH	94		
LD IXde C	9E	LD H, IXdeL	95		
LD IXde D	9F	LD H, HL	90		
LD IXde E	90	LD H, IXdeH	91		
LD IXde F	91	LD H, IXdeL	92		
LD IXde G	92	LD H, HL	93		
LD IXde H	93	LD H, IXdeH	94		
LD IXde I	94	LD H, IXdeL	95		
LD IXde J	95	LD H, HL	96		
LD IXde K	96	LD H, IXdeH	97		
LD IXde L	97	LD H, IXdeL	98		
LD IXde M	98	LD H, HL	99		
LD IXde N	99	LD H, IXdeH	90		
LD IXde O	9A	LD H, IXdeL	91		
LD IXde P	9B	LD H, HL	92		
LD IXde Q	9C	LD H, IXdeH	93		
LD IXde R	9D	LD H, IXdeL	94		
LD IXde S	9E	LD H, HL	95		
LD IXde T	9F	LD H, IXdeH	96		
LD IXde U	90	LD H, IXdeL	97		
LD IXde V	91	LD H, HL	98		
LD IXde W	92	LD H, IXdeH	99		
LD IXde X	93	LD H, IXdeL	90		
LD IXde Y	94	LD H, HL	91		
LD IXde Z	95	LD H, IXdeH	92		
LD IXde A	96	LD H, IXdeL	93		
LD IXde B	97	LD H, HL	94		
LD IXde C	98	LD H, IXdeH	95		
LD IXde D	99	LD H, IXdeL	90		
LD IXde E	9A	LD H, HL	91		
LD IXde F	9B	LD H, IXdeH	92		
LD IXde G	9C	LD H, IXdeL	93		
LD IXde H	9D	LD H, HL	94		
LD IXde I	9E	LD H, IXdeH	95		
LD IXde J	9F	LD H, IXdeL	96		
LD IXde K	90	LD H, HL	97		
LD IXde L	91	LD H, IXdeH	98		
LD IXde M	92	LD H, IXdeL	99		
LD IXde N	93	LD H, HL	90		
LD IXde O	94	LD H, IXdeH	91		
LD IXde P	95	LD H, IXdeL	92		
LD IXde Q	96	LD H, HL	93		
LD IXde R	97	LD H, IXdeH	94		
LD IXde S	98	LD H, IXdeL	95		
LD IXde T	99	LD H, HL	96		
LD IXde U	9A	LD H, IXdeH	97		
LD IXde V	9B	LD H, IXdeL	98		
LD IXde W	9C	LD H, HL	99		
LD IXde X	9D	LD H, IXdeH	90		
LD IXde Y	9E	LD H, IXdeL	91		
LD IXde Z	9F	LD H, HL	92		
LD IXde A	90	LD H, IXdeH	93		
LD IXde B	91	LD H, IXdeL	94		
LD IXde C	92	LD H, HL	95		
LD IXde D	93	LD H, IXdeH	96		
LD IXde E	94				

MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL
SET B	F35	RLC C	CB01	SET 1 L	CB CD
SET C	F34	RLC O	CB02	SET 2 (HL)	CB D6
SET E	F32	RLC E	CB03	SET 2 (IX+dis)	DD CB XX D6
SET H	F31	RLC H	CB04	SET 2 (IV+dis)	FD CB XX D6
SET L	F30	RLC L	CB05	SET 2 A	CB D7
SET	F2F	RLCA	CB07	SET 2,B	CB D0
RLD	ED8F	RLD	CB1E	SET 2,C	CB D1
DD-CB XX AB	FF1H	RR X	CB1F	SET 2,D	CB D2
CB AB	CB1E	RR Y	CB1G	SET 2,H	CB D4
DD-CB XX AE	FF1H	RR X+dis	CB1H	SET 2,L	CB D6
CB AE	CB1I	RR Y+dis	CB1J	SET 3, (HL)	CB DE
CB A7	CB1A	RRA	CB1K	SET 3, /	L MAX LL
CB A0	CB1L	RRA	CB1M	SET 3, (IV+dis)	FD CB XX DE
CB A1	CB1N	RRA	CB1P	SET 3,A	CB DF
CB A4	CB1Q	RRA	CB1R	SET 3,B	CB D6
CB A6	CB1S	RRA	CB1T	SET 3,C	CB D9
CB A8	CB1U	RRA	CB1V	SET 3,D	CB DA
CB A9	CB1W	RRC	CB1X	SET 3,E	CB DB
CB AA	CB1Y	RRC	CB1Z	SET 3,H	CB D4
CB AB	CB1A	RRC	CB2B	SET 3,L	CB DL
CB AC	CB2C	RRC	CB2D	SET 4, (HL)	CB E6
CB AD	CB2E	RRC	CB2F	SET 4, (IV+dis)	FD CB XX E6
CB AE	CB2G	RRC	CB2H	SET 4,A	CB E7
CB AF	CB2I	RRC	CB2J	SET 4,B	CB E0
CB AG	CB2K	RRC	CB2L	SET 4,C	CB E1
CB AH	CB2M	RRC	CB2N		
CB AI	CB2O	RRC	CB2P		
CB AJ	CB2Q	RRC	CB2R		
CB AK	CB2S	RRC	CB2T		
CB AL	CB2U	RRC	CB2V		
CB AM	CB2W	RRC	CB2X		
CB AN	CB2Y	RRC	CB2Z		
CB AO	CB2A	RRC	CB2B		
CB AP	CB2C	RRC	CB2D		
CB AR	CB2E	RRC	CB2F		
CB AS	CB2G	RRC	CB2H		
CB AT	CB2I	RRC	CB2J		
CB AU	CB2K	RRC	CB2L		
CB AV	CB2M	RRC	CB2N		
CB AW	CB2O	RRC	CB2P		
CB AX	CB2Q	RRC	CB2R		
CB AZ	CB2S	RRC	CB2T		
CB BA	CB2U	RRC	CB2V		
CB BB	CB2W	RRC	CB2X		
CB BC	CB2Y	RRC	CB2Z		
CB BD	CB2A	RRC	CB2B		
CB BE	CB2C	RRC	CB2D		
CB BF	CB2E	RRC	CB2F		
CB BG	CB2G	RRC	CB2H		
CB BH	CB2I	RRC	CB2J		
CB BI	CB2K	RRC	CB2L		
CB BJ	CB2M	RRC	CB2N		
CB BK	CB2O	RRC	CB2P		
CB BL	CB2Q	RRC	CB2R		
CB BM	CB2S	RRC	CB2T		
CB BN	CB2U	RRC	CB2V		
CB BO	CB2W	RRC	CB2X		
CB BP	CB2Y	RRC	CB2Z		
CB BR	CB2A	RRC	CB2B		
CB BS	CB2C	RRC	CB2D		
CB BT	CB2E	RRC	CB2F		
CB BU	CB2G	RRC	CB2H		
CB BV	CB2I	RRC	CB2J		
CB BW	CB2K	RRC	CB2L		
CB BX	CB2M	RRC	CB2N		
CB BY	CB2O	RRC	CB2P		
CB BZ	CB2Q	RRC	CB2R		
SET 0	(HL)	SET 0	(HL)	SET 7 A	CB FF
SET 0	(IX+dis)	SET 0	(IX+dis)	SET 7 B	CB FB
SET 0	(IV+dis)	SET 0	(IV+dis)	SET 7 C	CB FA
SET 0A		SET 0A		SET 7 D	CB FD
SET 0B		SET 0B		SET 7 E	CB FE
SET 0C		SET 0C		SET 7 F	DD CB XX FE
SET 0D		SET 0D		SET 7 G	FD CB XX FE
SET 0E		SET 0E		SET 7 H	
SET 0F		SET 0F		SET 7 I	
SET 0G		SET 0G		SET 7 J	
SET 0H		SET 0H		SET 7 K	
SET 0I		SET 0I		SET 7 L	
SET 0J		SET 0J		SET 7 M	
SET 0K		SET 0K		SET 7 N	
SET 0L		SET 0L		SET 7 O	
SET 0M		SET 0M		SET 7 P	
SET 0N		SET 0N		SET 7 Q	
SET 0O		SET 0O		SET 7 R	
SET 0P		SET 0P		SET 7 S	
SET 0Q		SET 0Q		SET 7 T	
SET 0R		SET 0R		SET 7 U	
SET 0S		SET 0S		SET 7 V	
SET 0T		SET 0T		SET 7 W	
SET 0U		SET 0U		SET 7 X	
SET 0V		SET 0V		SET 7 Y	
SET 0W		SET 0W		SET 7 Z	
SET 0X		SET 0X		SET 7 A	
SET 0Y		SET 0Y		SET 7 B	
SET 0Z		SET 0Z		SET 7 C	
SET 1A		SET 1A		SET 7 D	
SET 1B		SET 1B		SET 7 E	
SET 1C		SET 1C		SET 7 F	
SET 1D		SET 1D		SET 7 G	
SET 1E		SET 1E		SET 7 H	
SET 1F		SET 1F		SET 7 I	
SET 1G		SET 1G		SET 7 J	
SET 1H		SET 1H		SET 7 K	
SET 1I		SET 1I		SET 7 L	
SET 1J		SET 1J		SET 7 M	
SET 1K		SET 1K		SET 7 N	
SET 1L		SET 1L		SET 7 O	
SET 1M		SET 1M		SET 7 P	
SET 1N		SET 1N		SET 7 Q	
SET 1O		SET 1O		SET 7 R	
SET 1P		SET 1P		SET 7 S	
SET 1Q		SET 1Q		SET 7 T	
SET 1R		SET 1R		SET 7 U	
SET 1S		SET 1S		SET 7 V	
SET 1T		SET 1T		SET 7 W	
SET 1U		SET 1U		SET 7 X	
SET 1V		SET 1V		SET 7 Y	
SET 1W		SET 1W		SET 7 Z	
SET 1X		SET 1X		SET 8 A	
SET 1Y		SET 1Y		SET 8 B	
SET 1Z		SET 1Z		SET 8 C	
SET 2A		SET 2A		SET 8 D	
SET 2B		SET 2B		SET 8 E	
SET 2C		SET 2C		SET 8 F	
SET 2D		SET 2D		SET 8 G	
SET 2E		SET 2E		SET 8 H	
SET 2F		SET 2F		SET 8 I	
SET 2G		SET 2G		SET 8 J	
SET 2H		SET 2H		SET 8 K	
SET 2I		SET 2I		SET 8 L	
SET 2J		SET 2J		SET 8 M	
SET 2K		SET 2K		SET 8 N	
SET 2L		SET 2L		SET 8 O	
SET 2M		SET 2M		SET 8 P	
SET 2N		SET 2N		SET 8 Q	
SET 2O		SET 2O		SET 8 R	
SET 2P		SET 2P		SET 8 S	
SET 2Q		SET 2Q		SET 8 T	
SET 2R		SET 2R		SET 8 U	
SET 2S		SET 2S		SET 8 V	
SET 2T		SET 2T		SET 8 W	
SET 2U		SET 2U		SET 8 X	
SET 2V		SET 2V		SET 8 Y	
SET 2W		SET 2W		SET 8 Z	
SET 2X		SET 2X		SET 9 A	
SET 2Y		SET 2Y		SET 9 B	
SET 2Z		SET 2Z		SET 9 C	
SET 3A		SET 3A		SET 9 D	
SET 3B		SET 3B		SET 9 E	
SET 3C		SET 3C		SET 9 F	
SET 3D		SET 3D		SET 9 G	
SET 3E		SET 3E		SET 9 H	
SET 3F		SET 3F		SET 9 I	
SET 3G		SET 3G		SET 9 J	
SET 3H		SET 3H		SET 9 K	
SET 3I		SET 3I		SET 9 L	
SET 3J		SET 3J		SET 9 M	
SET 3K		SET 3K		SET 9 N	
SET 3L		SET 3L		SET 9 O	
SET 3M		SET 3M		SET 9 P	
SET 3N		SET 3N		SET 9 Q	
SET 3O		SET 3O		SET 9 R	
SET 3P		SET 3P		SET 9 S	
SET 3Q		SET 3Q		SET 9 T	
SET 3R		SET 3R		SET 9 U	
SET 3S		SET 3S		SET 9 V	
SET 3T		SET 3T		SET 9 W	
SET 3U		SET 3U		SET 9 X	
SET 3V		SET 3V		SET 9 Y	
SET 3W		SET 3W		SET 9 Z	
SET 4A		SET 4A		SET 10 A	
SET 4B		SET 4B		SET 10 B	
SET 4C		SET 4C		SET 10 C	
SET 4D		SET 4D		SET 10 D	
SET 4E		SET 4E		SET 10 E	
SET 4F		SET 4F		SET 10 F	
SET 4G		SET 4G		SET 10 G	
SET 4H		SET 4H		SET 10 H	
SET 4I		SET 4I		SET 10 I	
SET 4J		SET 4J		SET 10 J	
SET 4K		SET 4K		SET 10 K	
SET 4L		SET 4L		SET 10 L	
SET 4M		SET 4M		SET 10 M	
SET 4N		SET 4N		SET 10 N	
SET 4O		SET 4O		SET 10 O	
SET 4P		SET 4P		SET 10 P	
SET 4Q		SET 4Q		SET 10 Q	
SET 4R		SET 4R		SET 10 R	
SET 4S		SET 4S		SET 10 S	
SET 4T		SET 4T		SET 10 T	
SET 4U		SET 4U		SET 10 U	
SET 4V		SET 4V		SET 10 V	
SET 4W		SET 4W		SET 10 W	
SET 4X		SET 4X		SET 10 X	
SET 4Y		SET 4Y		SET 10 Y	
SET 4Z		SET 4Z		SET 10 Z	
SET 5A		SET 5A		SET 11 A	
SET 5B		SET 5B		SET 11 B	
SET 5C		SET 5C		SET 11 C	
SET 5D		SET 5D		SET 11 D	
SET 5E		SET 5E		SET 11 E	
SET 5F		SET 5F		SET 11 F	
SET 5G		SET 5G		SET 11 G	
SET 5H		SET 5H		SET 11 H	
SET 5I		SET 5I		SET 11 I	
SET 5J		SET 5J		SET 11 J	
SET 5K		SET 5K		SET 11 K	
SET 5L		SET 5L		SET 11 L	
SET 5M		SET 5M		SET 11 M	
SET 5N		SET 5N		SET 11 N	
SET 5O		SET 5O		SET 11 O	
SET 5P		SET 5P		SET 11 P	
SET 5Q		SET 5Q		SET 11 Q	
SET 5R		SET 5R		SET 11 R	
SET 5S		SET 5S		SET 11 S	
SET 5T		SET 5T		SET 11 T	
SET 5U		SET 5U		SET 11 U	
SET 5V		SET 5V		SET 11 V	
SET 5W		SET 5W		SET 11 W	
SET 5X		SET 5X		SET 11 X	
SET 5Y		SET 5Y		SET 11 Y	
SET 5Z		SET 5Z		SET 11 Z	
SET 6A		SET 6A		SET 12 A	
SET 6B		SET 6B		SET 12 B	
SET 6C		SET 6C		SET 12 C	
SET 6D		SET 6D		SET 12 D	
SET 6E		SET 6E		SET 12 E	
SET 6F		SET 6F		SET 12 F	
SET 6G		SET 6G		SET 12 G	
SET 6H		SET 6H		SET 12 H	
SET 6I		SET 6I		SET 12 I	
SET 6J		SET 6J		SET 12 J	
SET 6K		SET 6K		SET 12 K	
SET 6L		SET 6L		SET 12 L	
SET 6M		SET 6M		SET 12 M	
SET 6N		SET 6N		SET 12 N	
SET 6O		SET 6O		SET 12 O	
SET 6P		SET 6P		SET 12 P	
SET 6Q		SET 6Q		SET 12 Q	
SET 6R		SET 6R		SET 12 R	
SET 6S		SET 6S		SET 12 S	
SET 6T		SET 6T		SET 12 T	
SET 6U		SET 6U		SET 12 U	
SET 6V		SET 6V		SET 12 V	
SET 6W		SET 6W		SET 12 W	
SET 6X		SET 6X</td			

MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL	MNEMONIC	HEXADECIMAL
SRA Y+dn	F C9 XX ZE				
SRA A	0 0				
SRA B	0 0				
SRA C	0 0				
SRA D	0 0				
SRA E	0 0				
SRA H	C0 0				
SRA L	0 0				
SRL (HL)	0 0				
~P ~A	0 XX BE				
SAL (Y+dn)	F 5 XH JE				
H A	0 F				
H L	0 39				
A	0 19				
O	0 1A				
S P	0 0				
S L P	0 0				
S	0 0				
S RL H	96				
G H X Z	L 36 RX				
G H P	F 36 RX				
~B A	0 9				
B H	0 0				
~B H	0 0				
N H	0 0				
S H E	0 0				
~ H ~	0 0				
X P	0 A				
X Y	0 AE XX				
Y X	0 AE XX				
X A	0 A				
X D H	0 A				
X D	0 A				
X D	0 A				
X D H	0 XX				
X D H	0 A				
X D H	0 A				
X D H	0 A				

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DATAE0	11
DATAE8	11
DATAF0	11
DATAF8	11
DATAH0	11
DATAH8	11
DATAI0	11
DATAI8	11
DATAJ0	11
DATAJ8	11
DATAK0	11
DATAK8	11
DATAO0	11
DATAO8	11
DATAP0	11
DATAP8	11
DATAQ0	11
DATAQ8	11
DATAR0	11
DATAR8	11
DATAU0	11
DATAU8	11
DATAV0	11
DATAV8	11
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